

Foundations of Model Transformations

A "lambda calculus" for MDD ?

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Motivation

- * At the heart of model-driven engineering are activities like
 - maintaining consistency
 - evolution
 - translation
 - execution
- * A math. foundation is needed for studying
 - Expressiveness and complexity
 - Execution and optimisation
 - Well-definedness
 - Semantic correctness of transformations
- * These are examples of model transformations
- * This lecture is about graph transformations as one such foundation

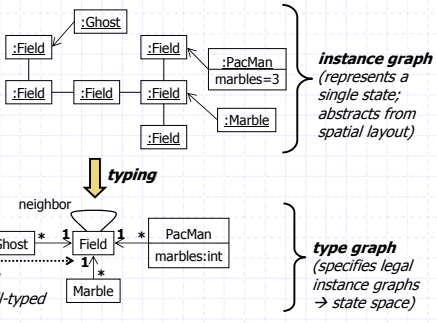
Outline

- * Graph transformation
 - why it is fun
 - how it works
- * Model transformation
- * Theory and Tools

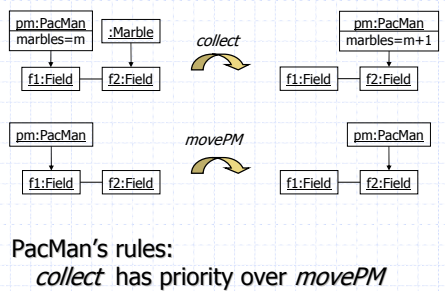
Why it is fun: Programming By Example

- StageCast (www.stagecast.com): a visual programming environment for kids (from 8 years on), based on
- behavioral rules associated to graphical objects
 - visual pattern matching
 - simple control structures (priorities, sequence, choice, ...)
 - external keyboard control
- intuitive rule-based behavior modelling
- Next: abstract from concrete visual presentation

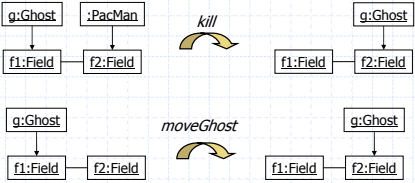
States of the PacMan Game: Graph-Based Presentation



Rules of the PacMan Game: Graph-Based Presentation, PacMan

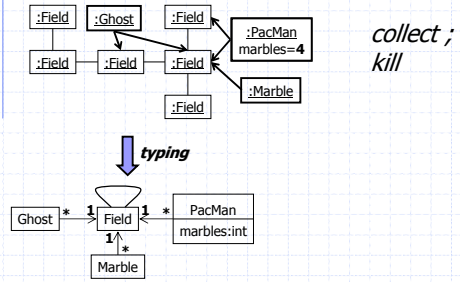


Rules of the PacMan Game: Graph-Based Presentation, Ghost



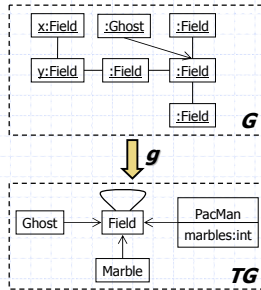
Ghost's rules:
kill has priority over *moveGhost*

Graph Transformation



How it works: Typed Graphs

Directed graphs as algebraic structures $G = (V, E, src, tar)$ with $src, tar: E \rightarrow V$
 Graph homomorphism as pair of mappings $h = (h_V, h_E): G_1 \rightarrow G_2$ with
 • $h_V: V_1 \rightarrow V_2$
 • $h_E: E_1 \rightarrow E_2$
 preserving src and tar
 Typed graphs given by
 • fixed type graph TG
 • instance graphs G typed over TG by homomorphism $g: G \rightarrow TG$



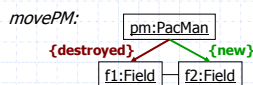
Rules

$p: L \rightarrow R$ with $L \cap R$ well-defined, in different presentations
 • like above (cf. PacMan example)
 • with $L \cap R$ explicit [DPO]: $L \leftarrow K \rightarrow R$

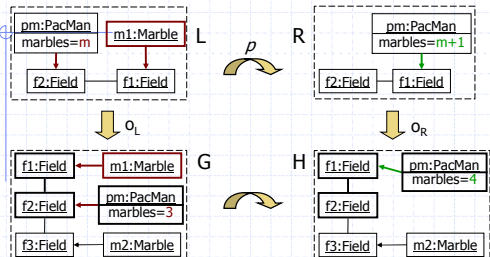


Rules

$p: L \rightarrow R$ with $L \cap R$ well-defined, in different presentations
 • like above (cf. PacMan example)
 • with $L \cap R$ explicit [DPO]: $L \leftarrow K \rightarrow R$
 • with L, R integrated [UML]: $L \cup R$ and marking
 • $L - R$ as {destroyed}
 • $R - L$ as {new}

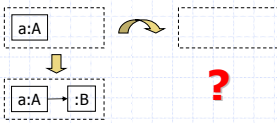


Transformation Step: Operational



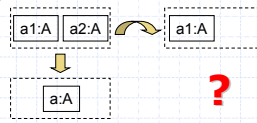
1. select rule $p: L \rightarrow R$; occurrence $o_L: L \rightarrow G$
2. remove from G the occurrence of $R \setminus L$
3. add to result a copy of $R \setminus L$

Semantic Questions: Dangling Edges



- ✳ conservative solution: application is forbidden
 - invertible transformations, no side-effects
- ✳ radical solution: *delete dangling edges*
 - more complex behavior, requires explicit control

Semantic Questions: Conflicts

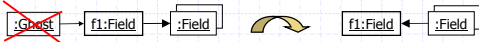


- ✳ conservative solution: application is forbidden
 - invertible transformations, no side-effects
- ✳ radical solution: *give priority to deletion*
 - more complex behavior, requires explicit control

Advanced Features

Dealing with unknown context

- set-nodes (multi-objects): match all nodes with the required connections
- explicit (negative) context conditions

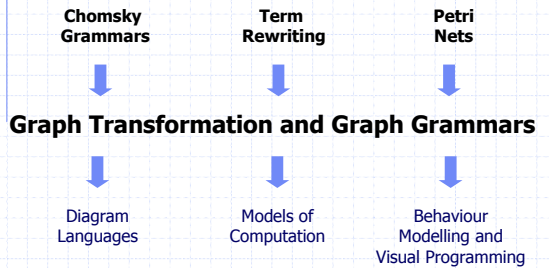


(turns f1 into a trap by reversing all outgoing edges to Field vertices, but only if there is no Ghost)

Control

- priorities: *movePM* only if *collect* is not possible
- programmed transformation: IF NOT *collect* THEN *movePM*;

Where it comes from and what it is good for

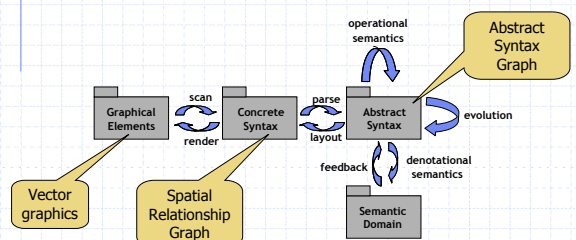


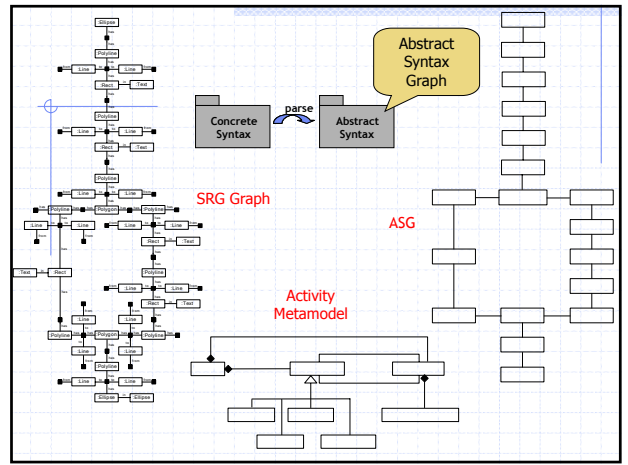
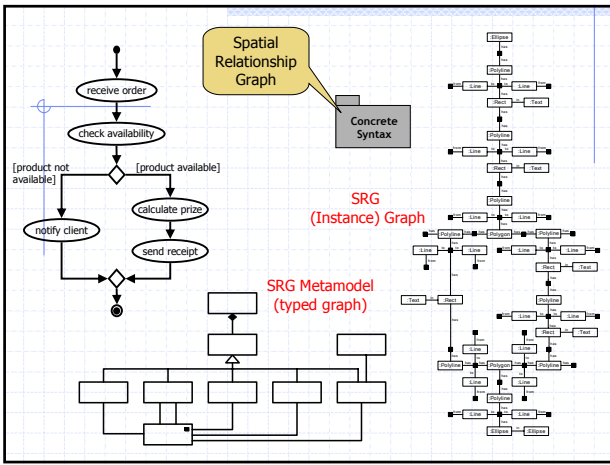
Outline

- ✓ Graph transformation
- ✳ Model transformation
 - diagram languages
 - execution
 - translation
- ✳ Theory and Tools

Diagram Languages

- ✳ theory and tools like for textual languages





Semantics

- * Operational semantics (execution)
 - concrete syntax: animation rules
 - abstract syntax: graph transformation rules
- * Denotational semantics (translation)

operational semantics

Abstract Syntax

denotational semantics

Semantic Domain

The diagram shows a bidirectional arrow between 'operational semantics' and 'Abstract Syntax', and another bidirectional arrow between 'denotational semantics' and 'Semantic Domain'.

Operational Semantics

- * diagram syntax plus *runtime state*
- * visual rules to model state transitions

operational semantics

Abstract Syntax

The diagram shows a state transition: a red oval (representing a state) transitions to a white oval labeled 'op(...)' (representing an operation). A curved arrow labeled 'op(...)' indicates the transition.

- * produce sequences of states / labels, visualized as animations

Extended Meta Model: Original plus *Runtime State*

operational semantics

Abstract Syntax

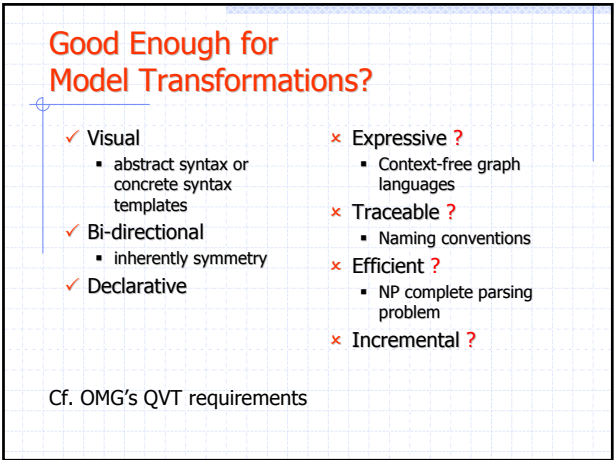
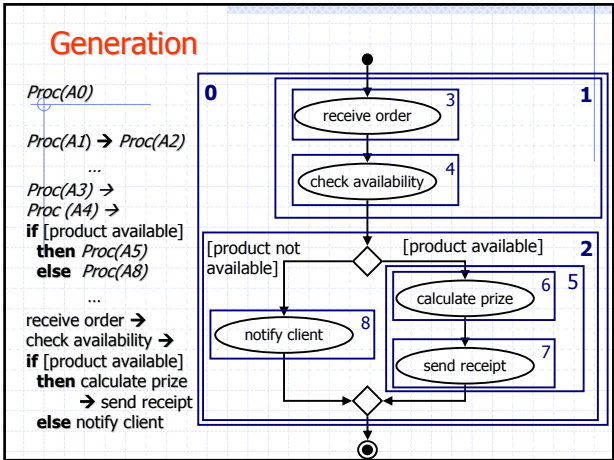
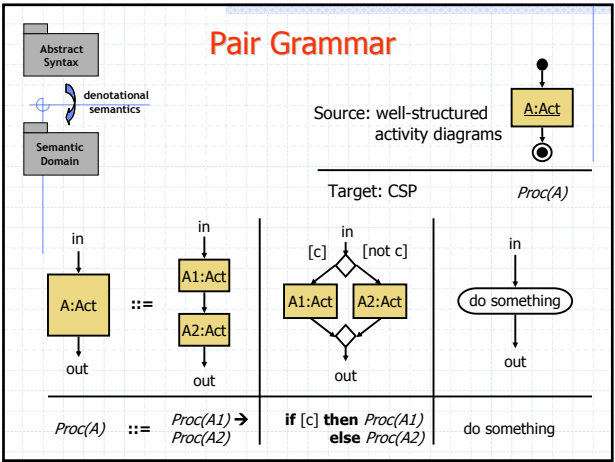
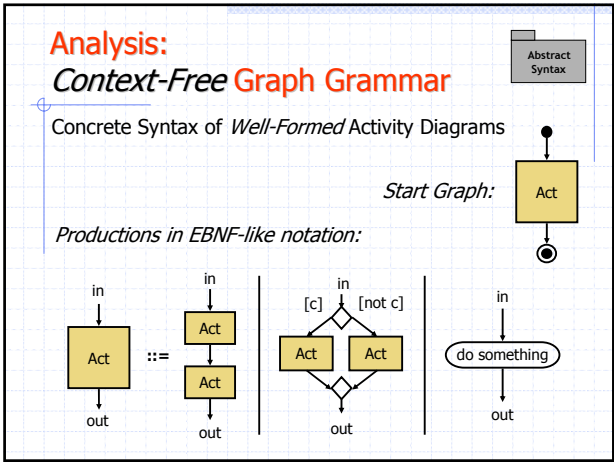
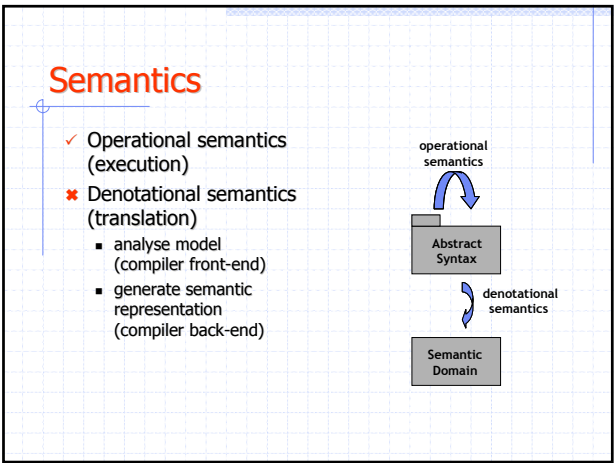
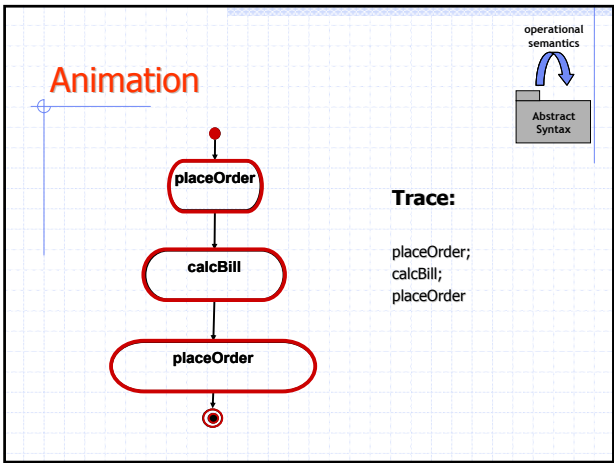
The diagram shows a class hierarchy with inheritance arrows. A red oval highlights a specific node in the hierarchy, representing a runtime state.

Operational Rules: GT on Meta Model Instances

operational semantics

Abstract Syntax

The diagram shows a graph transformation rule (GT) applied to meta-model instances. A curved arrow indicates the transformation from one graph state to another.



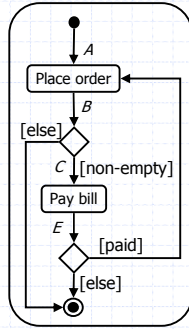
A Non-Well-Structured Example

Actions

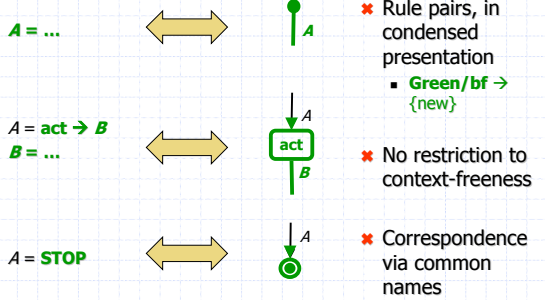
Place_order, Pay_bill

Processes

$A = \text{Place_order} \rightarrow B$
 $B = \text{if 'non-empty' then } C \text{ else STOP}$
 $C = \text{Pay_bill} \rightarrow E$
 $E = \text{if 'paid' then } A \text{ else STOP}$

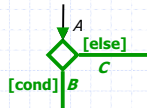


Correspondence Rules: Initial, Action, and Final Nodes

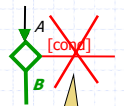


Correspondence Rules: Choice and Join

$A = \text{if cond then } B \text{ else } C$
 $B = \dots$
 $C = \dots$



$A = B$
 $B = \dots$

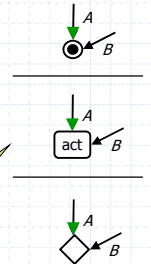


Negative application condition

Correspondence Rules: Connection to Existing Nodes

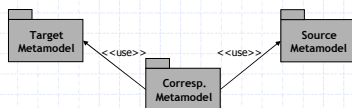
$A = B$

Either alternative is consistent with left side



Formally: Triple Graph Grammars

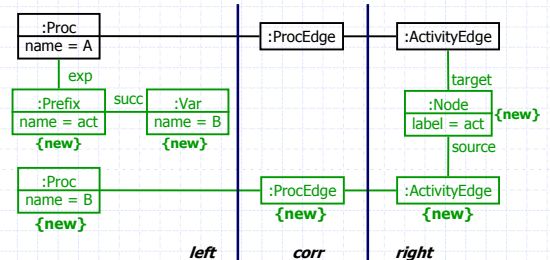
- Meta model for correspondence
 - traceability



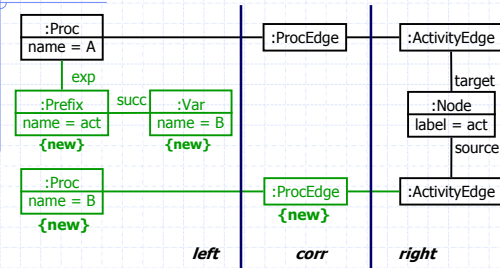
- Symmetric rule triplets (left, corr, right), generating directed rules
 - Declarative \rightarrow operational

Example TGG Rule

$A = \text{act} \rightarrow B$
 $B = \dots$



Derived Operational GT Rule: right → left



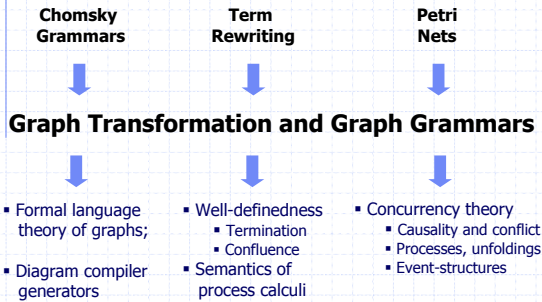
Alternatively:

- left → right
- (left, right) → corr

Outline

- ✓ Graph transformation
- ✓ Model transformation
- ✗ Theory and Tools

Theory: Sources of Inspirations

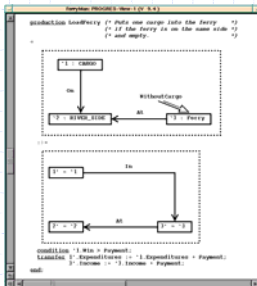


Outline

- ✓ Graph transformation
- ✓ Model transformation
- ✗ Theory and *Tools*
 - General purpose modeling environments
 - PROGRES, AGG, Fujaba, ...
 - Environments for specifying visual notations
 - DIAGEN, GENGE, ...
 - Analysis tools
 - Groove, ...

PROGRES

(PROgrammed Graph Rewriting Systems)

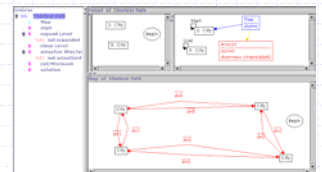


- ✗ Graphical/textual language to specify graph transformations
- ✗ Graph rewrite rules with complex and negative conditions
- ✗ Embedded in programming notation for OO data base transactions
- ✗ Cross compilation in Modula 2, C and Java

AGG

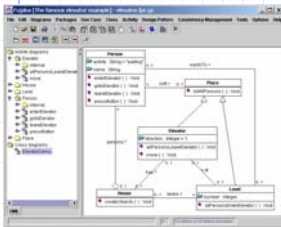
(The Attributed Graph Grammar System)

- ✗ Algebraic approach to graph transformation
- ✗ Attribute computations in Java
- ✗ Efficient graph parsing
- ✗ Analysis facilities for critical pairs, consistency of rules and invariants



Fujaba

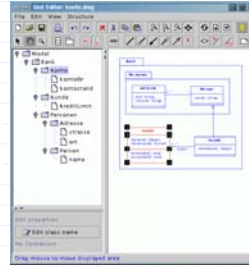
(From UML to Java and Back Again)



- * Round trip engineering with UML, Java, and design patterns
- * Class and activity diagrams; GT rules as collaborations
- * Generates standalone Java classes
- * Various plugins, including triple graph grammars

DiaGen

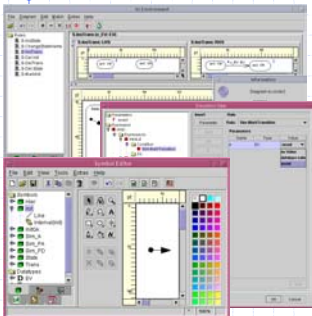
(The Diagram Editor Generator)



- * VL concrete / abstract syntax specified as CF hypergraph grammars
- * Generation of diagram editors, parsers, simulators

GenGED

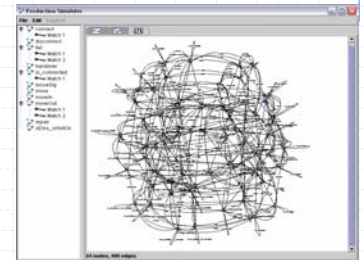
(Generation of Graphical Env.s for Design)



- * Roughly similar to the above, based on AGG

GRaphs for Object-Oriented VERification (GROOVE)

- * (bounded) generation of LTS from GT systems
 - edge-labelled graphs
 - application conditions
 - rule priorities



<http://wwwhome.cs.utwente.nl/~groove/groove-index>

Graph Transformation

A "lambda calculus" for MDD ?

- * Expressiveness and complexity
 - CF graph grammars: Rozenberg; Habel
 - Relation with MSO logic: Courcelle
- * Execution and optimisation
 - Various tools
 - Graph matching: Zuendorf 96; Varro et al 05
- * Well-definedness
 - Term graph rewriting: Plump
 - Confluence, termination of attributed GT: Ehrig et al 04
- * Semantic correctness
 - Requires formal semantics of the languages involved
 - operational: Hausmann et al
 - denotational: Kuester et al
 - Verification through
 - proofs: several authors
 - model checking: Varro; Rensink
 - testing: Koenig et al 04

Inspired?

Join us in Vienna!

ETAPS 2006

- * The event: 25 March - 2 April 2006

- * Abstract submission: Friday 7 October 2005

- * Paper submission: Friday 14 October 2005

Fundamental Approaches to Software Engineering



Vienna (Austria) March 27-29, 2006

GT-VMT* 2006

- * Paper submission: 12 December 2005

*: 5th Intl Workshop on Graph Transformation and Visual Modelling Techniques

Questions ?

