

# Component-based Construction of Real-time Systems in BIP

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in collaboration with

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**Workshop on Foundations and Applications of  
Component-based Design**

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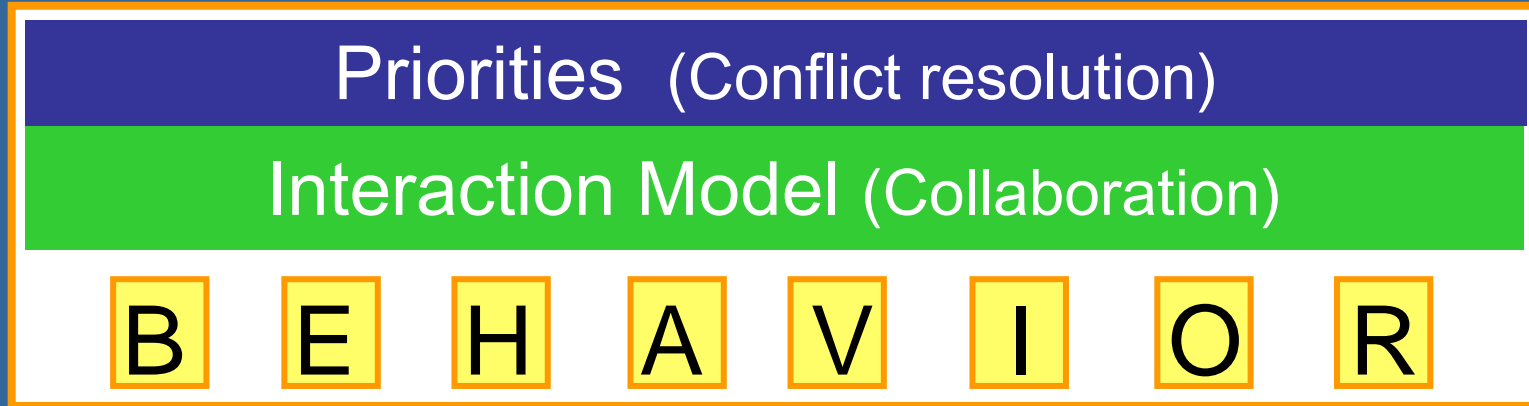
# Component-based construction – Objectives

Develop a rigorous and general basis for real-time system design and implementation:

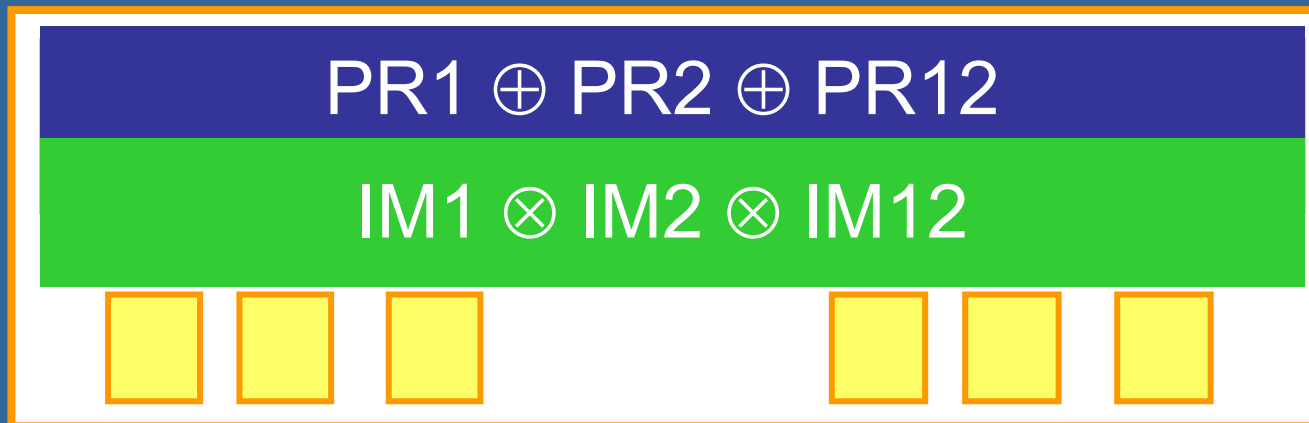
- Concept of component and associated composition operators for incremental description and correctness by construction
- Concept for real-time architecture encompassing heterogeneity, paradigms and styles of computation e.g.
  - Synchronous vs. asynchronous execution
  - Event driven vs. data driven computation
  - Distributed vs. centralized execution
- Automated support for component integration and generation of glue code meeting given requirements

# Component-based construction – The BIP framework

## Layered component model



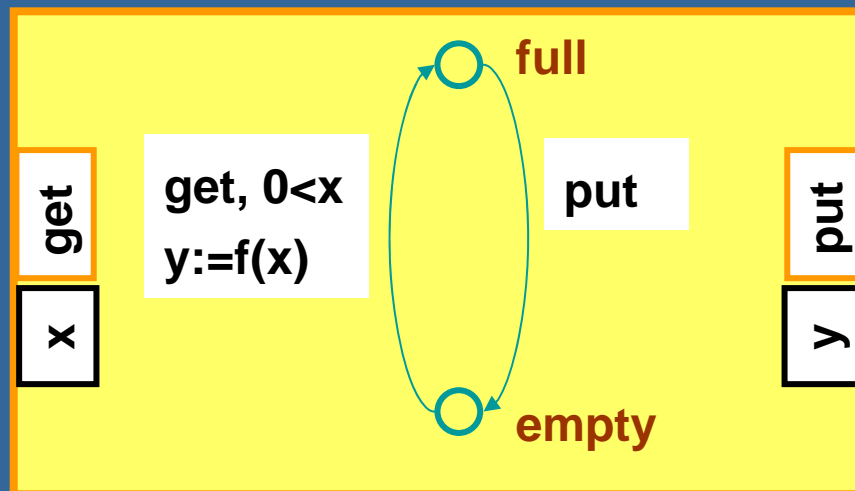
## Composition (incremental description)



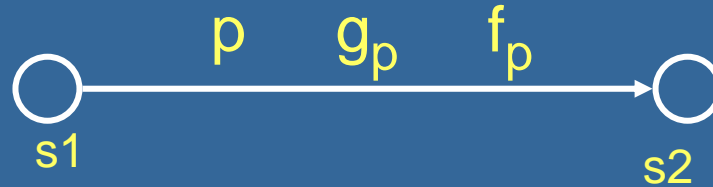
# Component-based construction – The BIP framework: Behavior

An atomic component has

- A set of ports  $P$ , for interaction with other components
- A set of control states  $S$
- A set of variables  $V$
- A set of transitions of the form
  - $p$  is a port
  - $g_p$  is a guard, boolean expression on  $V$
  - $f_p$  is a function on  $V$  (block of code)



# Component-based construction – The BIP framework: Behavior



$p$ : a port through which interaction is sought

$g_p$ : a pre-condition for interaction through  $p$

$f_p$ : a computation (local state transformation)

## Semantics

- **Enabledness:**  $g_p$  is true and some interaction involving  $p$  is possible
- **Execution:** interaction involving  $p$  followed by the execution of  $f_p$

# Overview



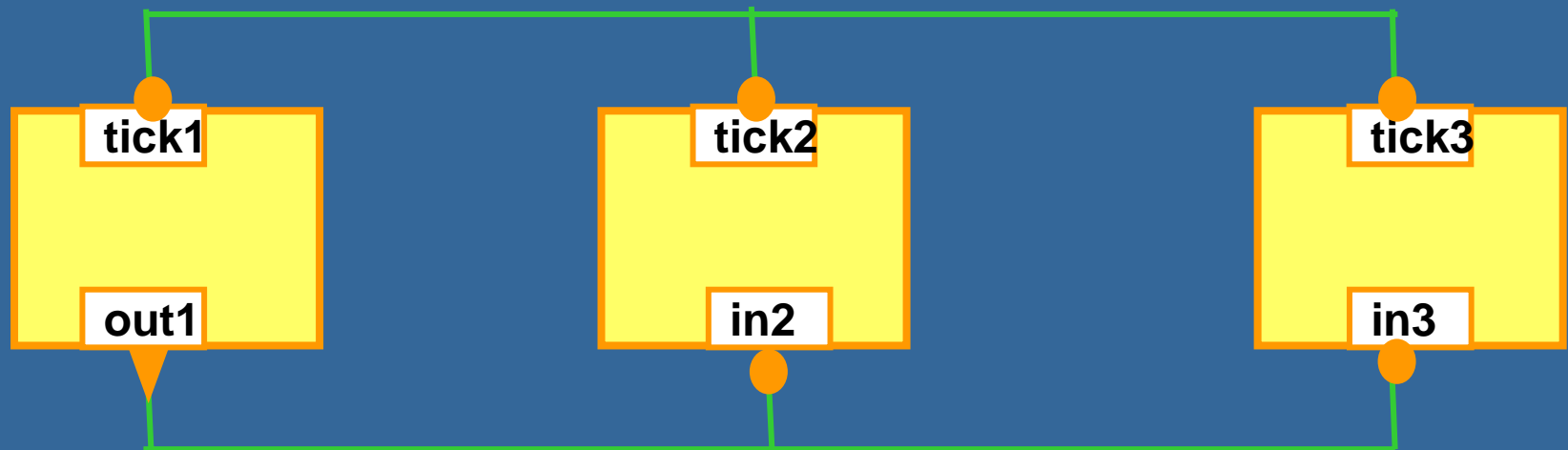
- Interaction modeling
- Priority modeling
- Implementation
- Modeling systems in BIP
- Discussion

# Interaction modeling

- A **connector** is a set of ports which can be involved in an interaction

- Port attributes (**complete** ▽, **incomplete** ●) are used to distinguish between rendezvous and broadcast.

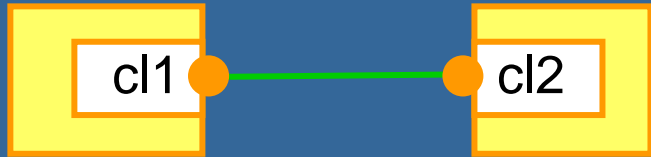
- An **interaction** of a connector is a set of ports such that: either it contains some complete port or it is maximal.



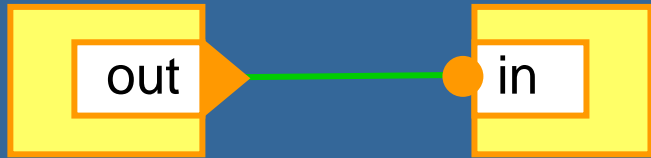
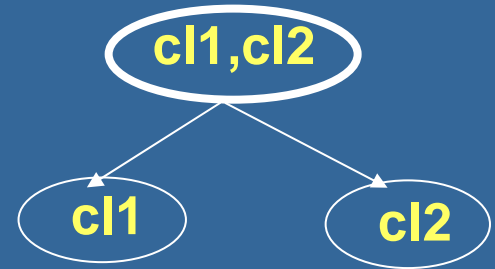
Interactions:

{tick1,tick2,tick3} {out1} {out1,in2} {out1,in3} {out1,in2, in3}

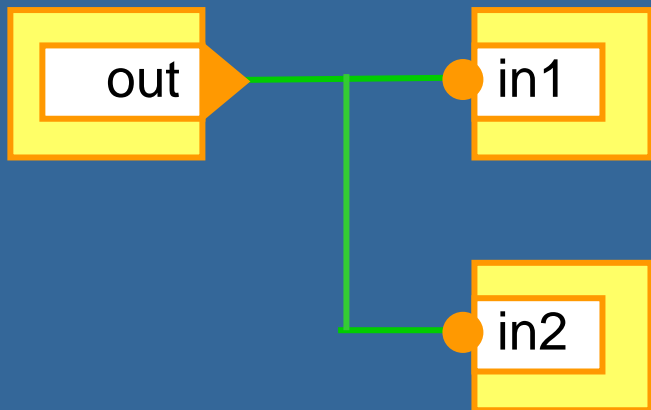
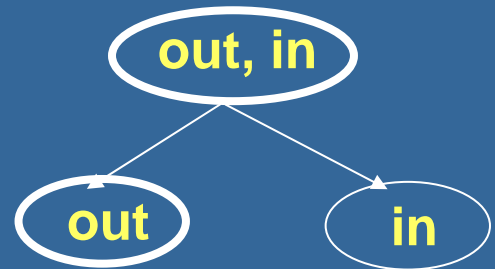
# Interaction modeling - Examples



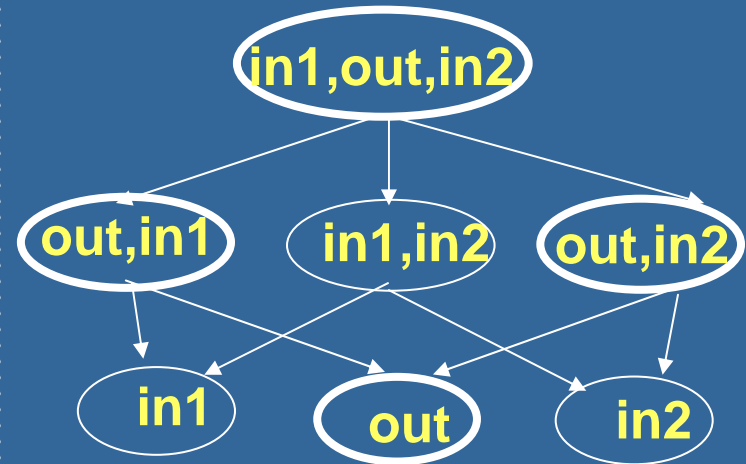
CN: {cl1, cl2}  
CP:  $\emptyset$



CN: {out, in}  
CP: {out}

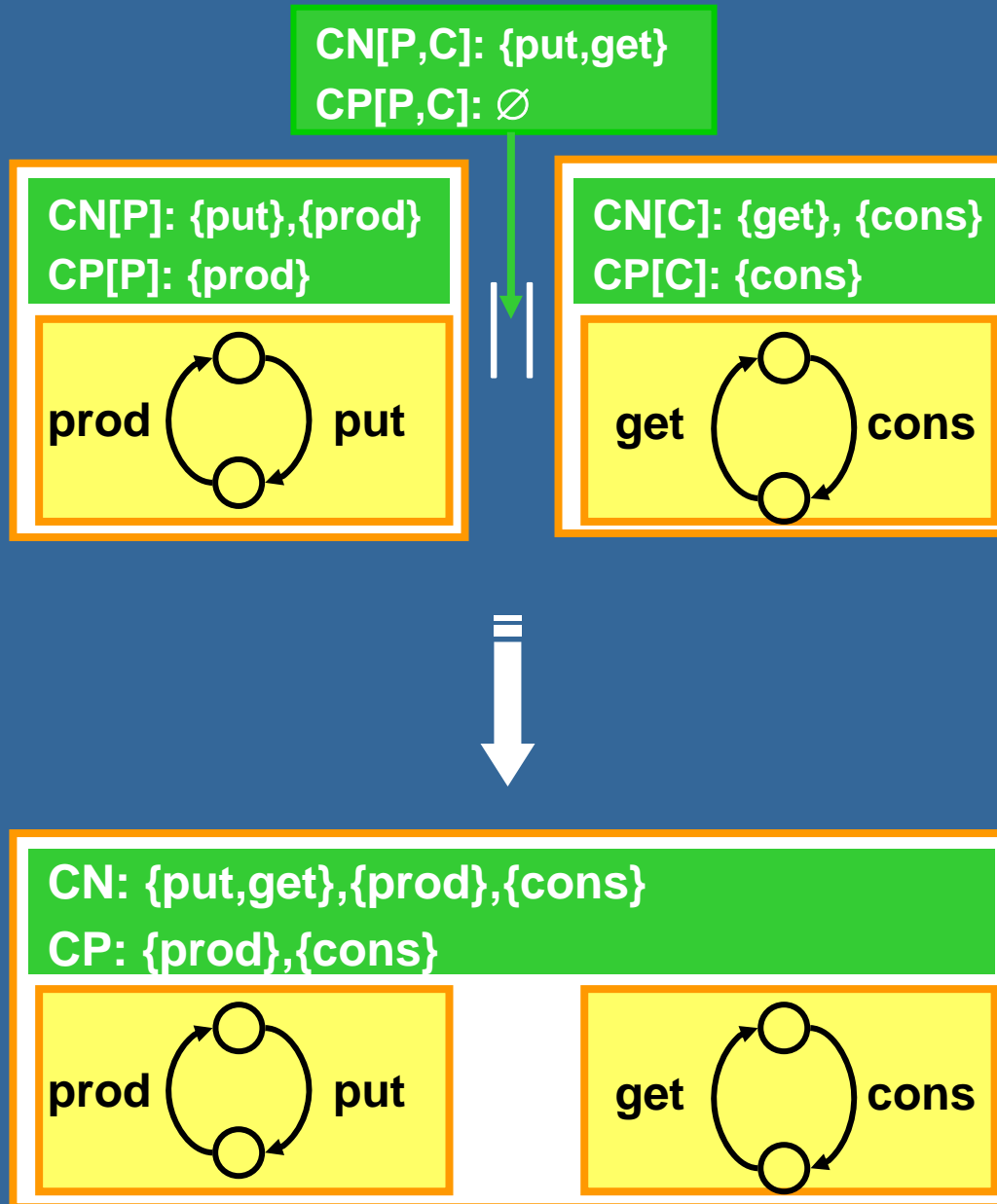


CN: {in1, out, in2}  
CP: {out}





# Interaction modeling – Composition



# Interaction modeling – Data transfer

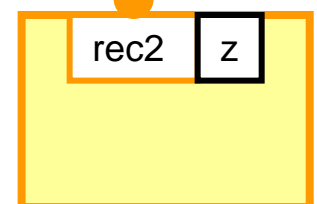
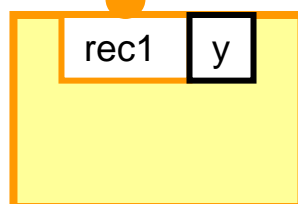
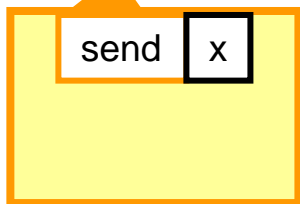
CN: BUS={send,rec1,rec2}

{send}: true  $\rightarrow$  skip

{send,rec1}:  $x < y \rightarrow x := y - x, y := y + x$

{send,rec2}:  $x < z \rightarrow x := z - x, z := z + x$

{send,rec1,rec2}:  $x < z + y \rightarrow x := y + z - x, y := y + x, z := z + x$



- Notice the difference between control flow and data flow (input, output)
- Maximal progress: execute a maximal enabled interaction

# Overview

- Interaction modeling
- Priority modeling
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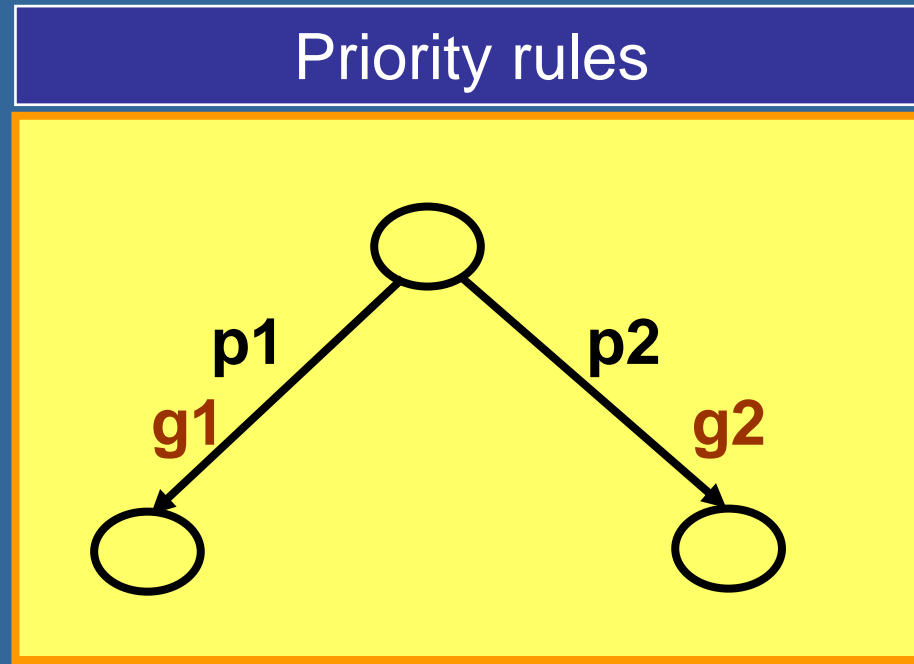


# Priorities

Priorities are a powerful tool for restricting non-determinism:

- they allow straightforward modeling of urgency and scheduling policies for real-time systems
- run to completion and synchronous execution can be modeled by assigning priorities to threads
- they can advantageously replace (static) restriction of process algebras

# Priorities - Definition

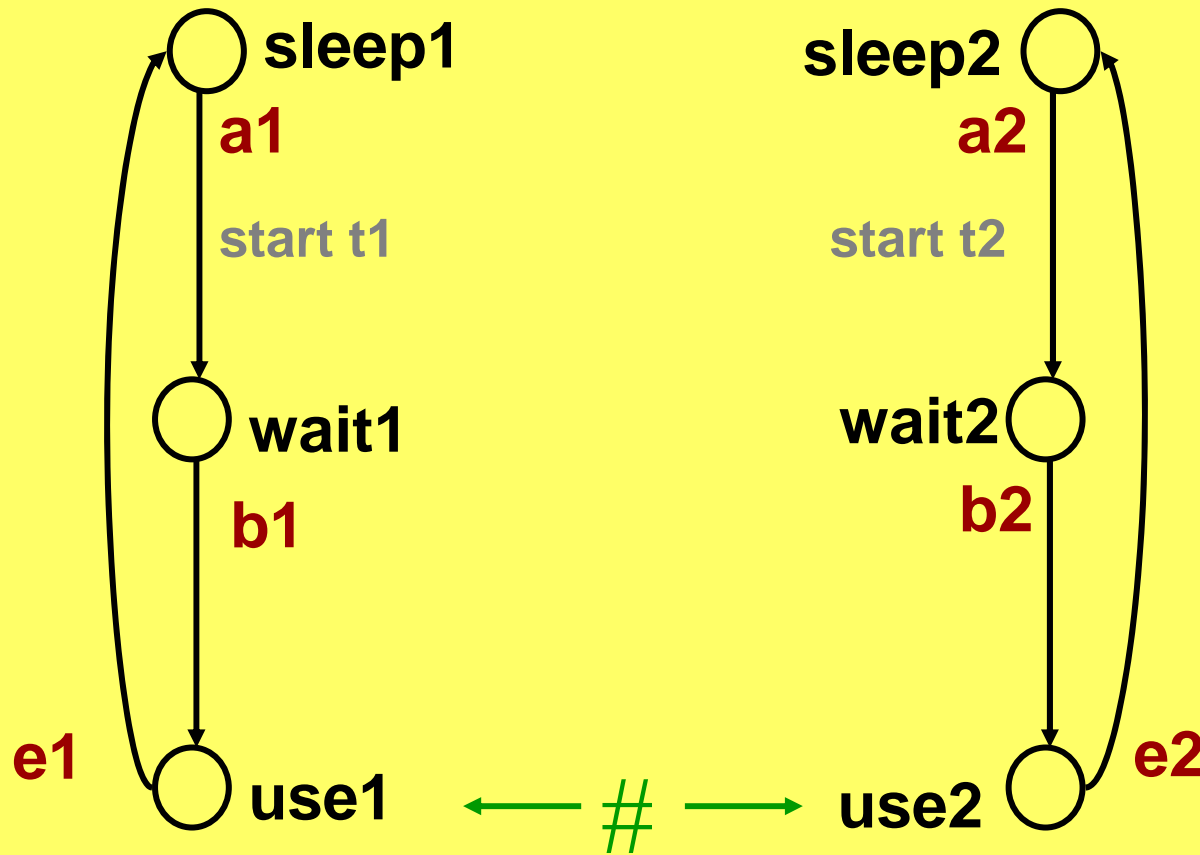


Priority rule	Restricted guard $g1'$
$\text{true} \rightarrow p1 \prec p2$	$g1' = g1 \wedge \neg g2$
$C \rightarrow p1 \prec p2$	$g1' = g1 \wedge \neg(C \wedge g2)$

# Priorities – Example: FIFO policy

$t1 \leq t2 \rightarrow b1 \prec b2$

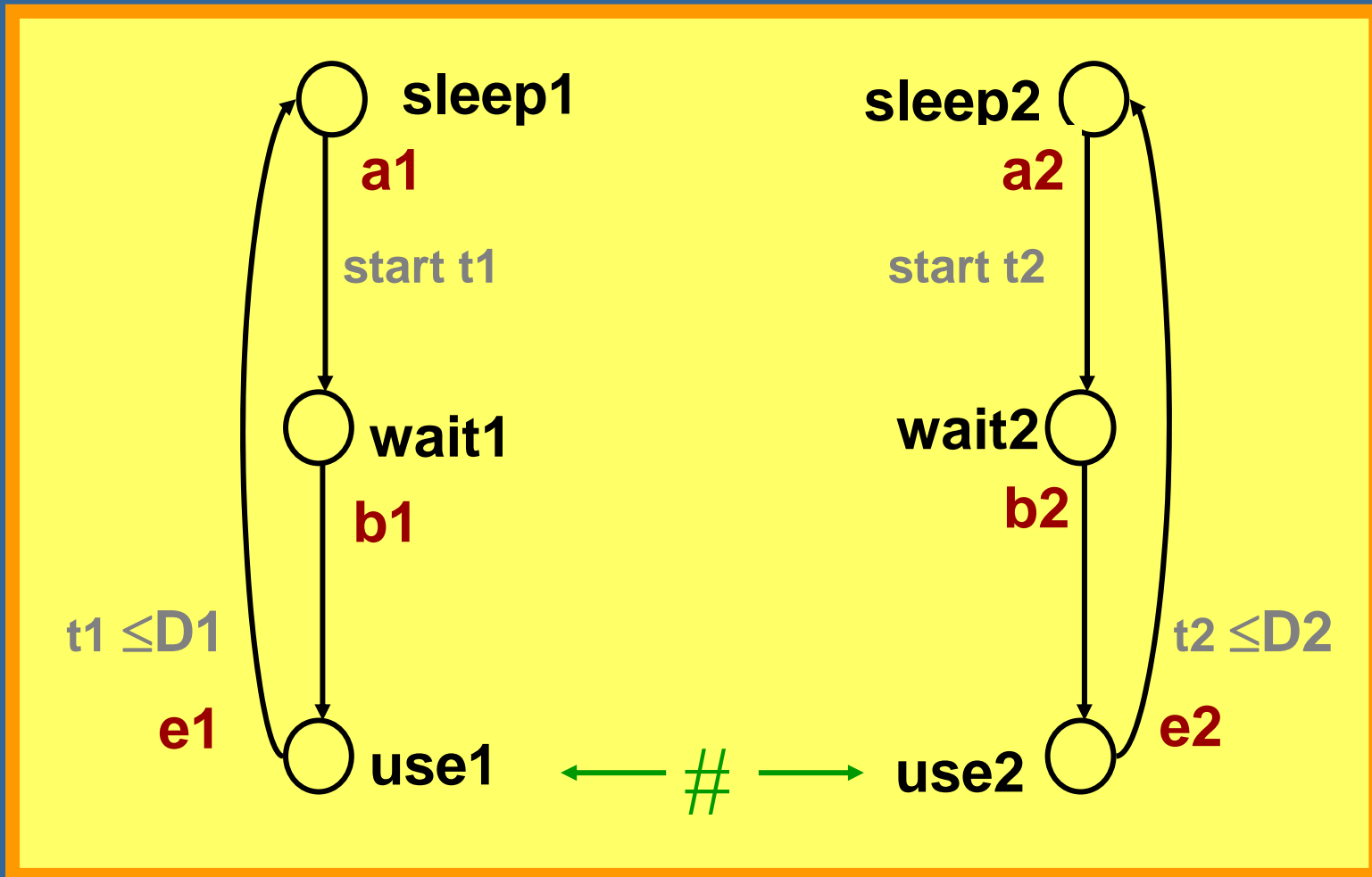
$t2 < t1 \rightarrow b2 \prec b1$



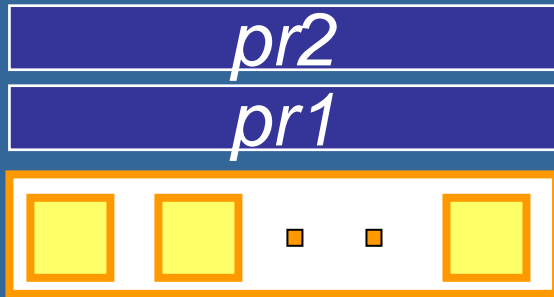
# Priorities – Example: EDF policy

$D1-t1 \leq D2-t2 \rightarrow b2 < b1$

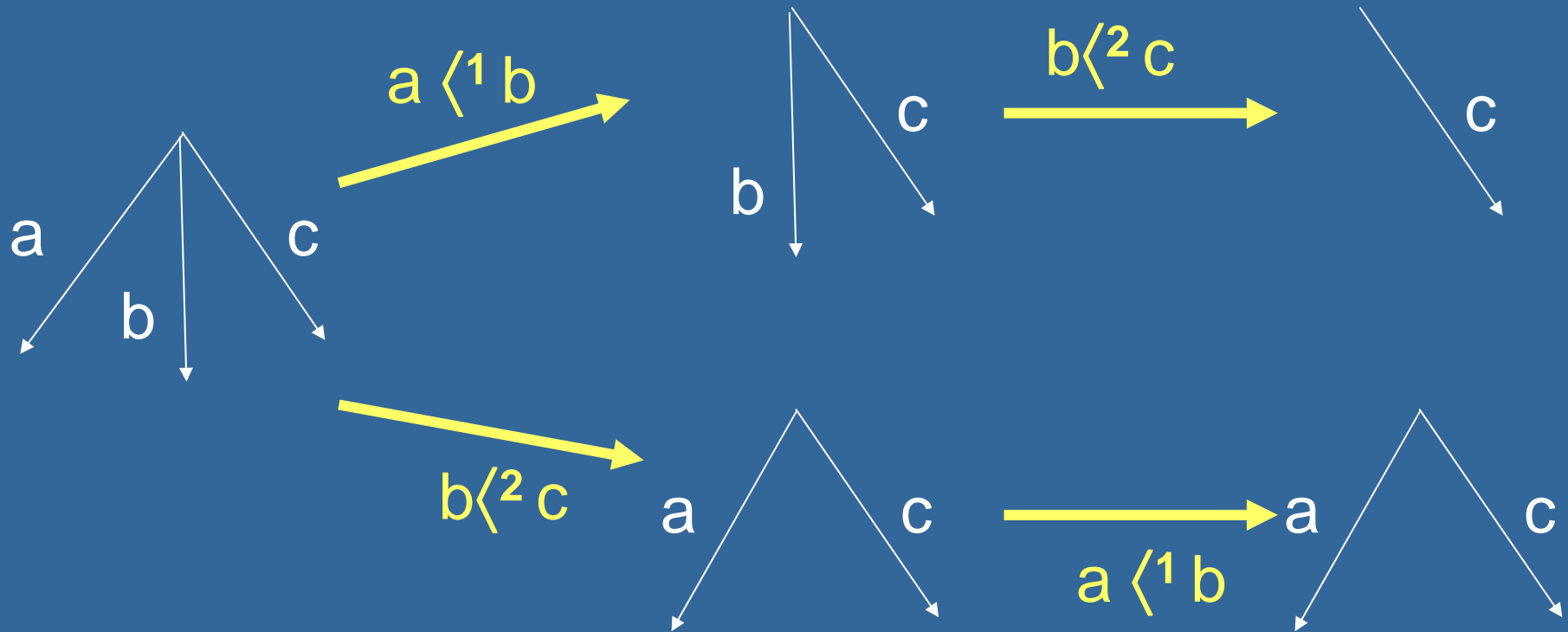
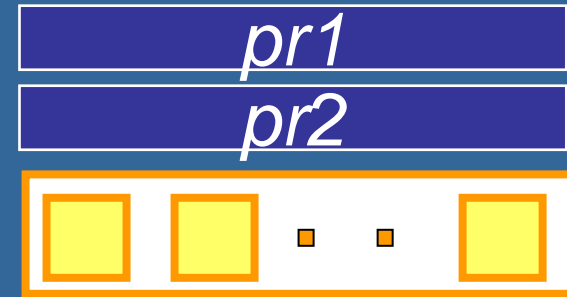
$D2-t2 < D1-t1 \rightarrow b1 < b2$



# Priorities – Composition



$\neq$





## Priorities – Composition (2)

Take:



$pr1 \oplus pr2$  is the least priority containing  $pr1 \cup pr2$

### Results :

- The operation  $\oplus$  is partial, associative and commutative
- $pr1(pr2(B)) \neq pr1(pr2(B))$
- $pr1 \oplus pr2(B)$  refines  $pr1 \cup pr2(B)$  refines  $pr1(pr2(B))$
- Priorities preserve deadlock-freedom

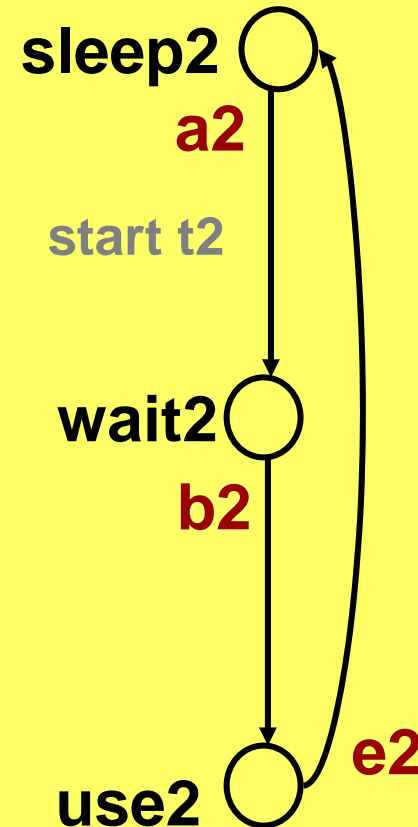
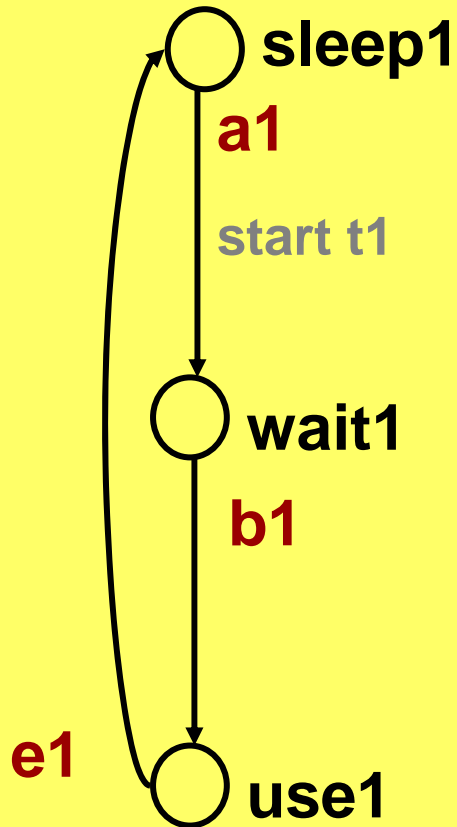
# Priorities – Example: Mutual exclusion + FIFO policy

$t1 \leq t2 \rightarrow b1 \prec b2$

$t2 < t1 \rightarrow b2 \prec b1$

$true \rightarrow b1 \prec e2$

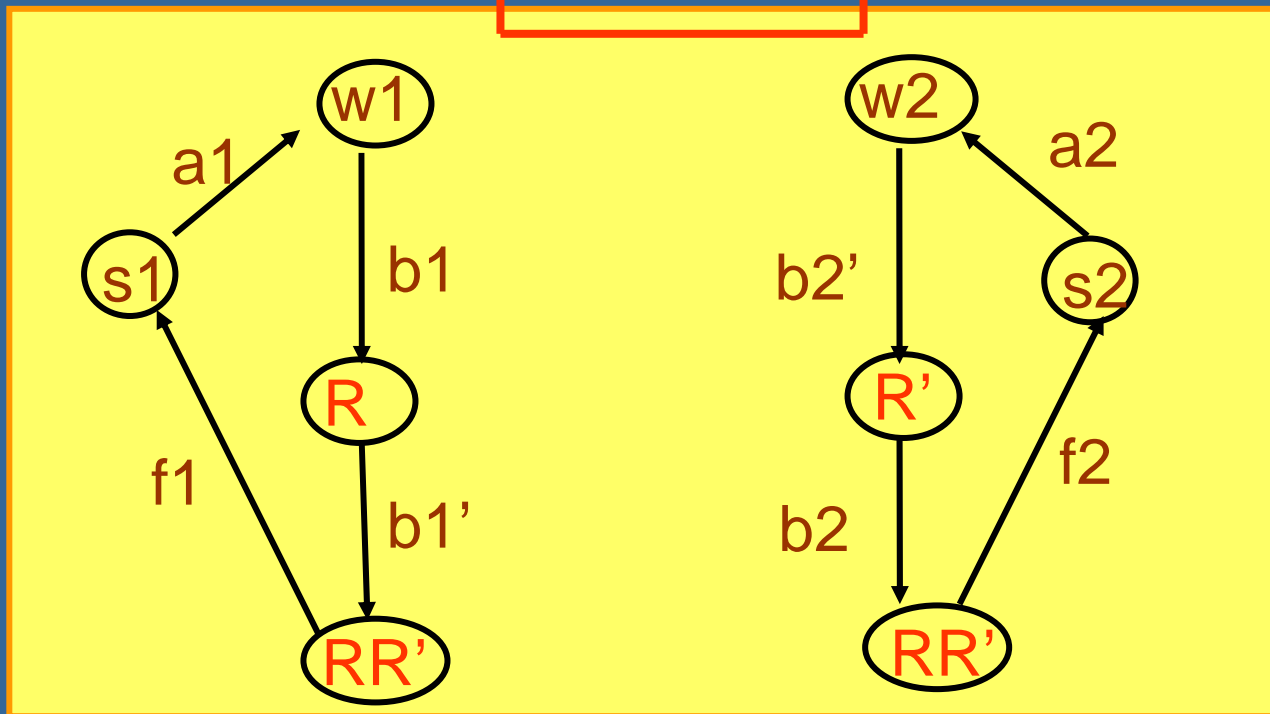
$true \rightarrow b2 \prec e1$



# Priorities – Checking for deadlock-freedom: Example


Mutex on  $R'$  :  $b1 \prec f2$     $b2 \prec \{f1, b1'\}$

Mutex on  $R$  :  $b1' \prec \{f2, b2\}$     $b2' \prec f1$

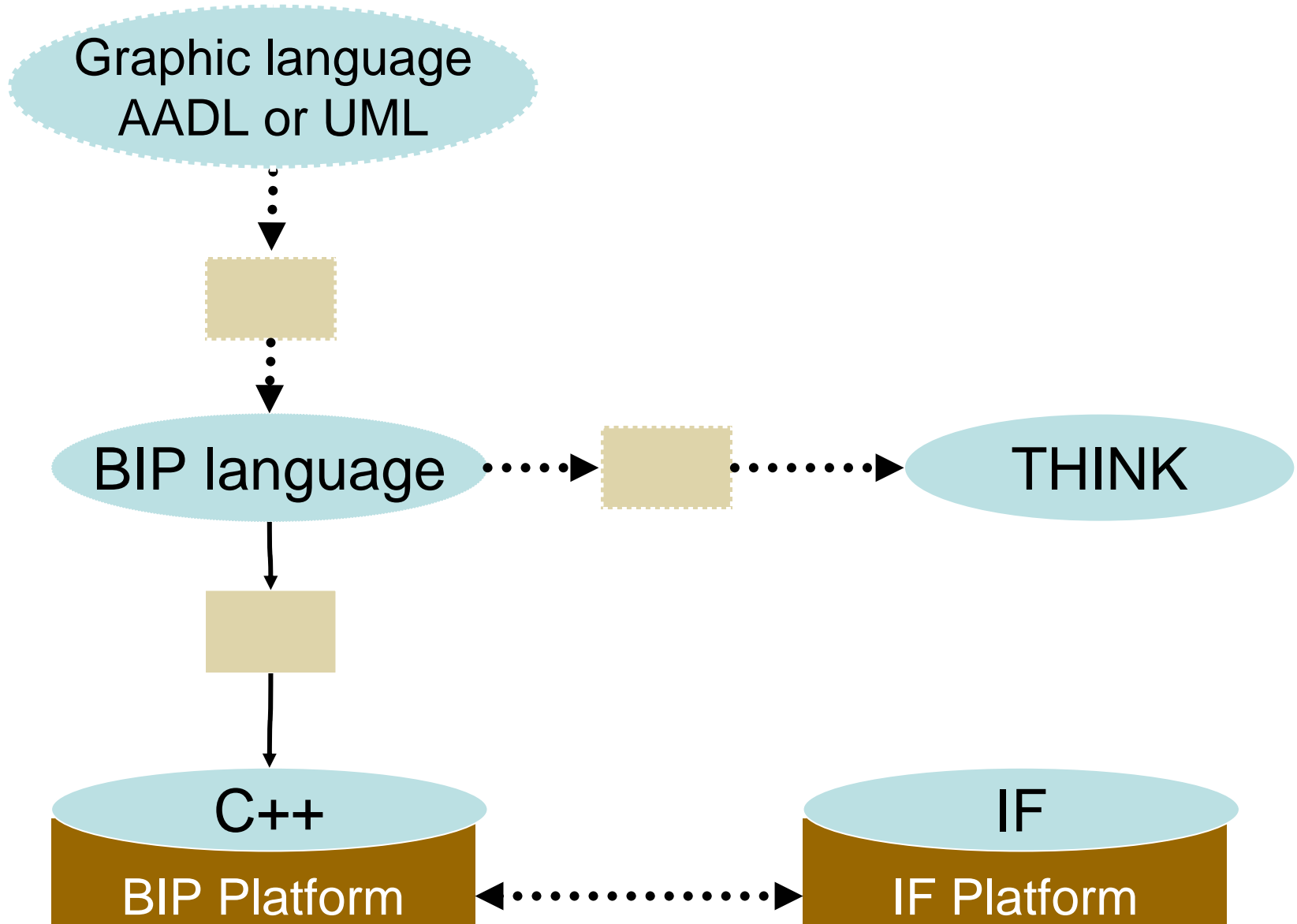


Risk of deadlock:  $b1' \prec b2$  and  $b2 \prec b1'$

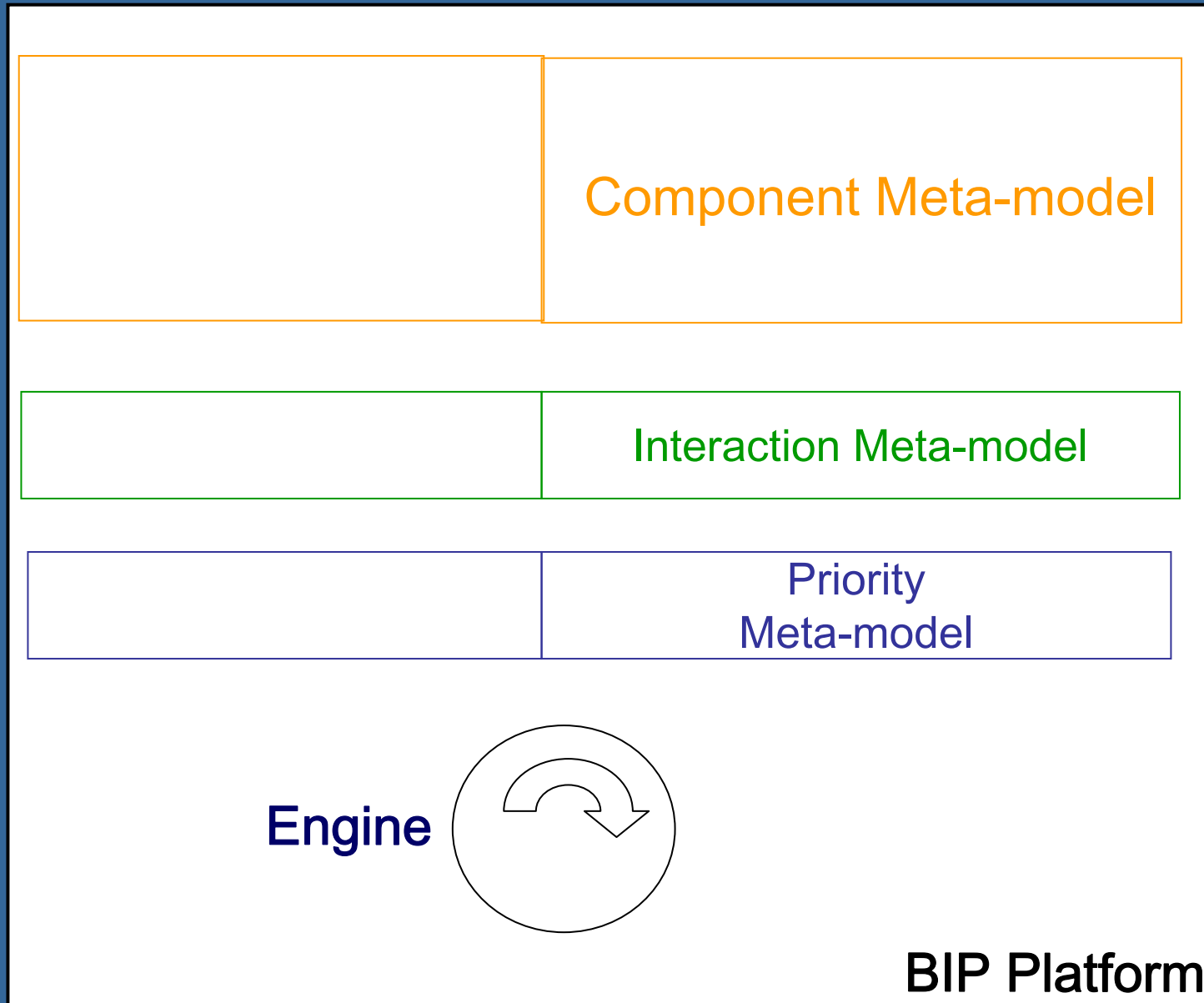
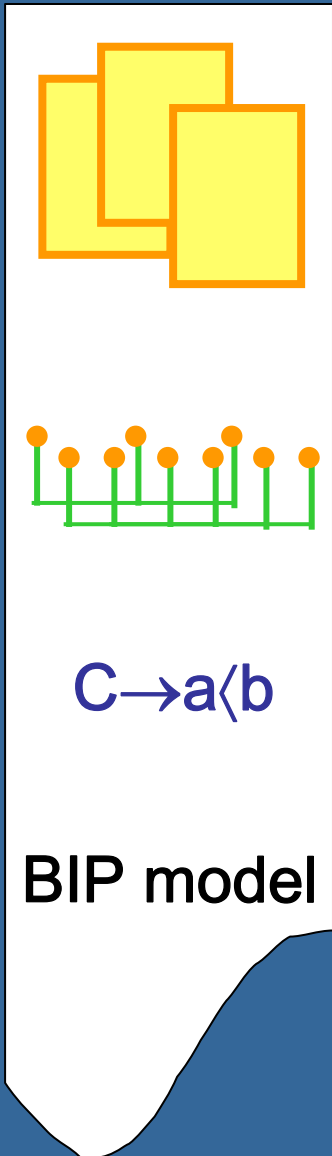
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-  • Implementation
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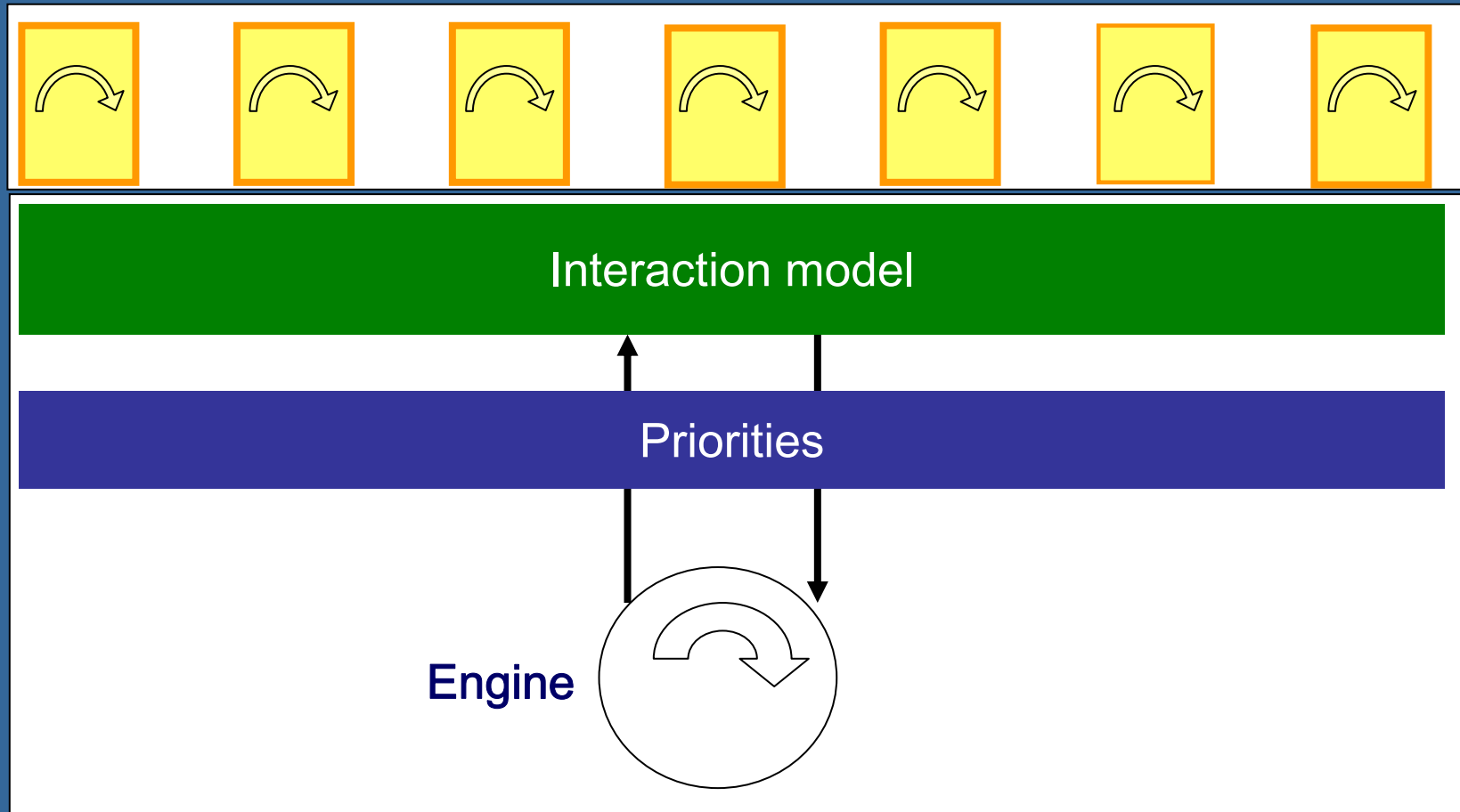
# Implementation – the BIP toolset



# Implementation – C++ code generation for the BIP platform

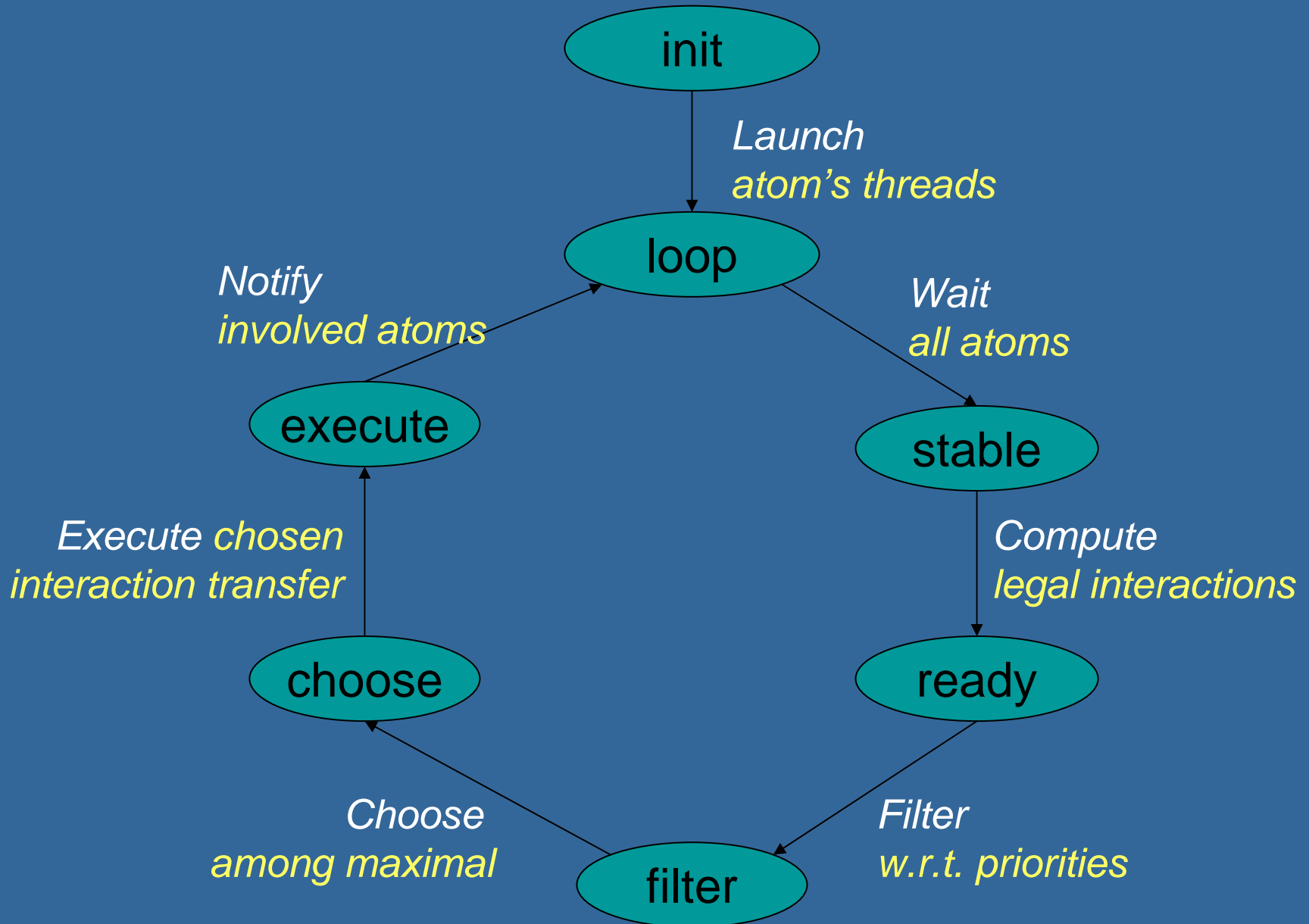


# Implementation – The BIP platform



- Code execution and state space exploration features
- Implementation in C++ on Linux using POSIX threads

# Implementation – The BIP platform: The engine





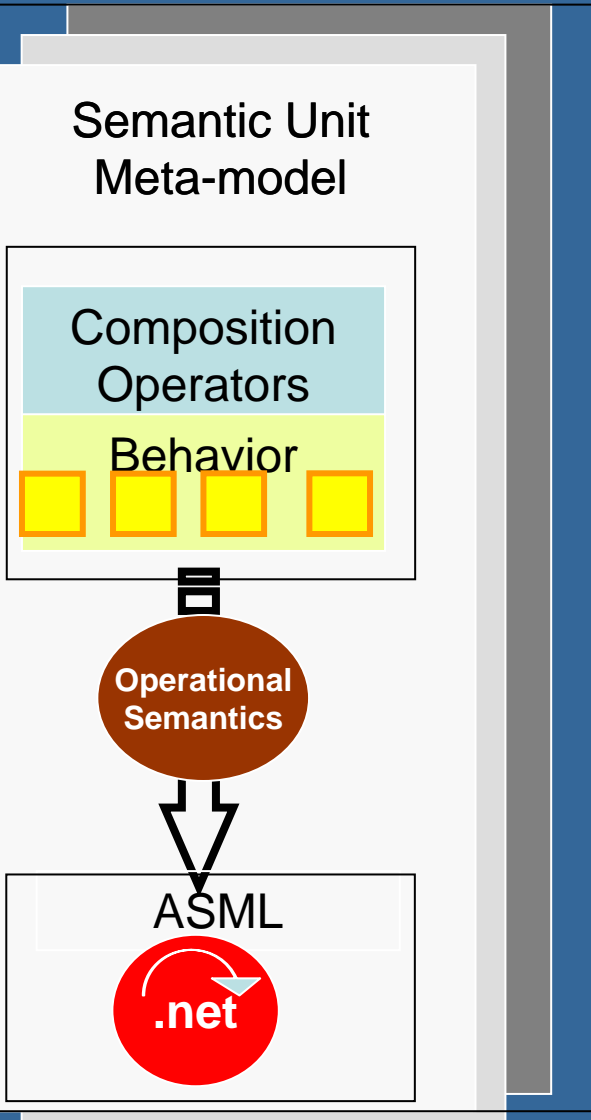
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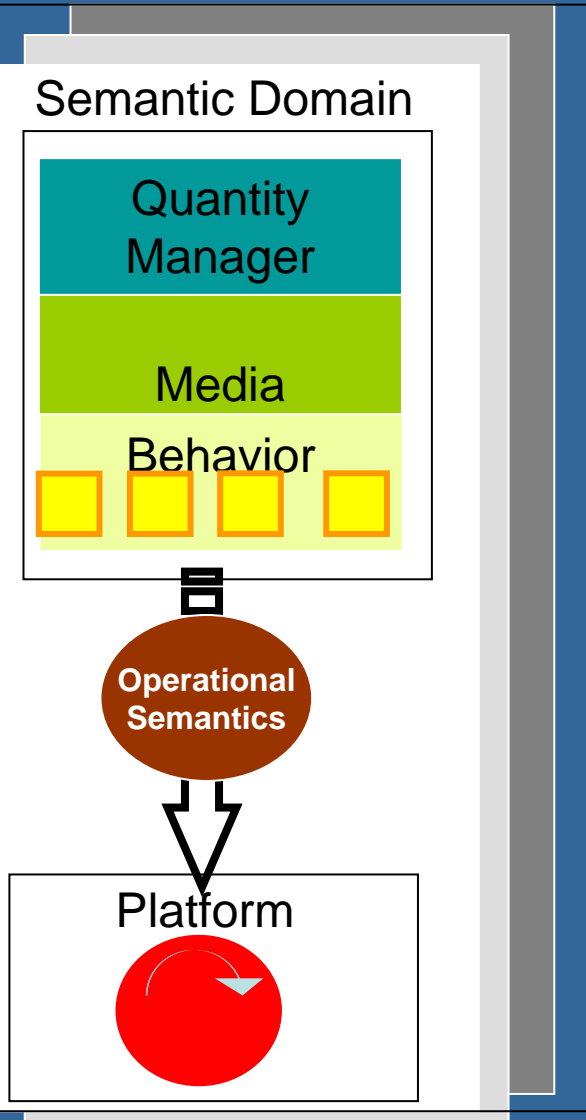


# Modeling in BIP– Other approaches encompassing heterogeneity

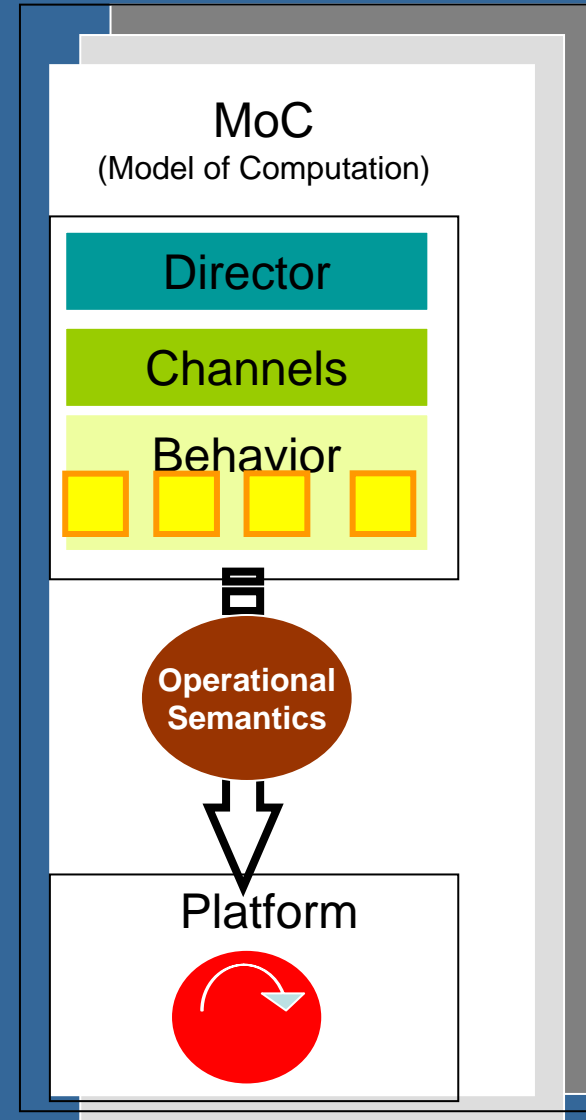
Vanderbilt's Approach



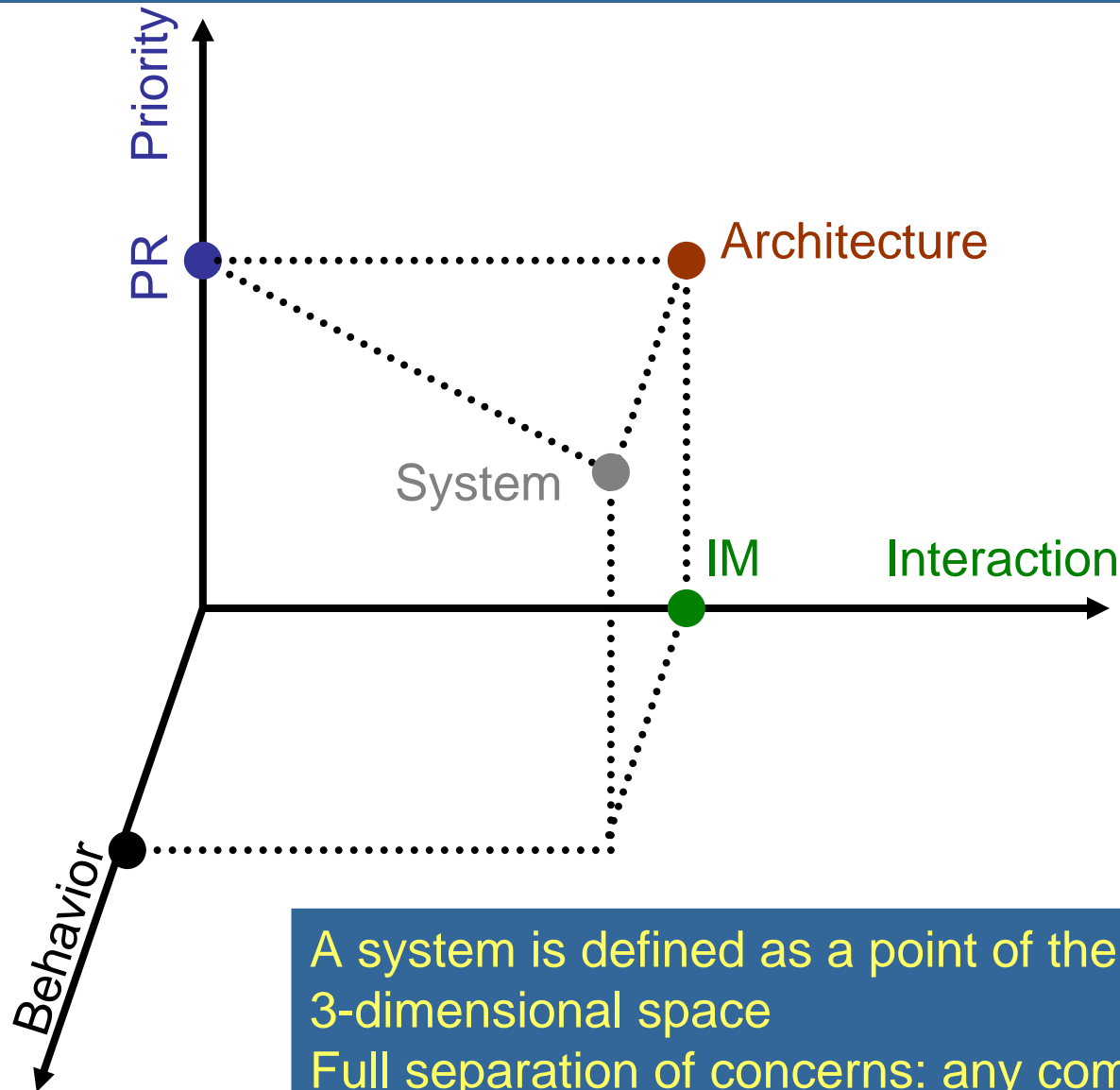
Metropolis



PTOLEMY

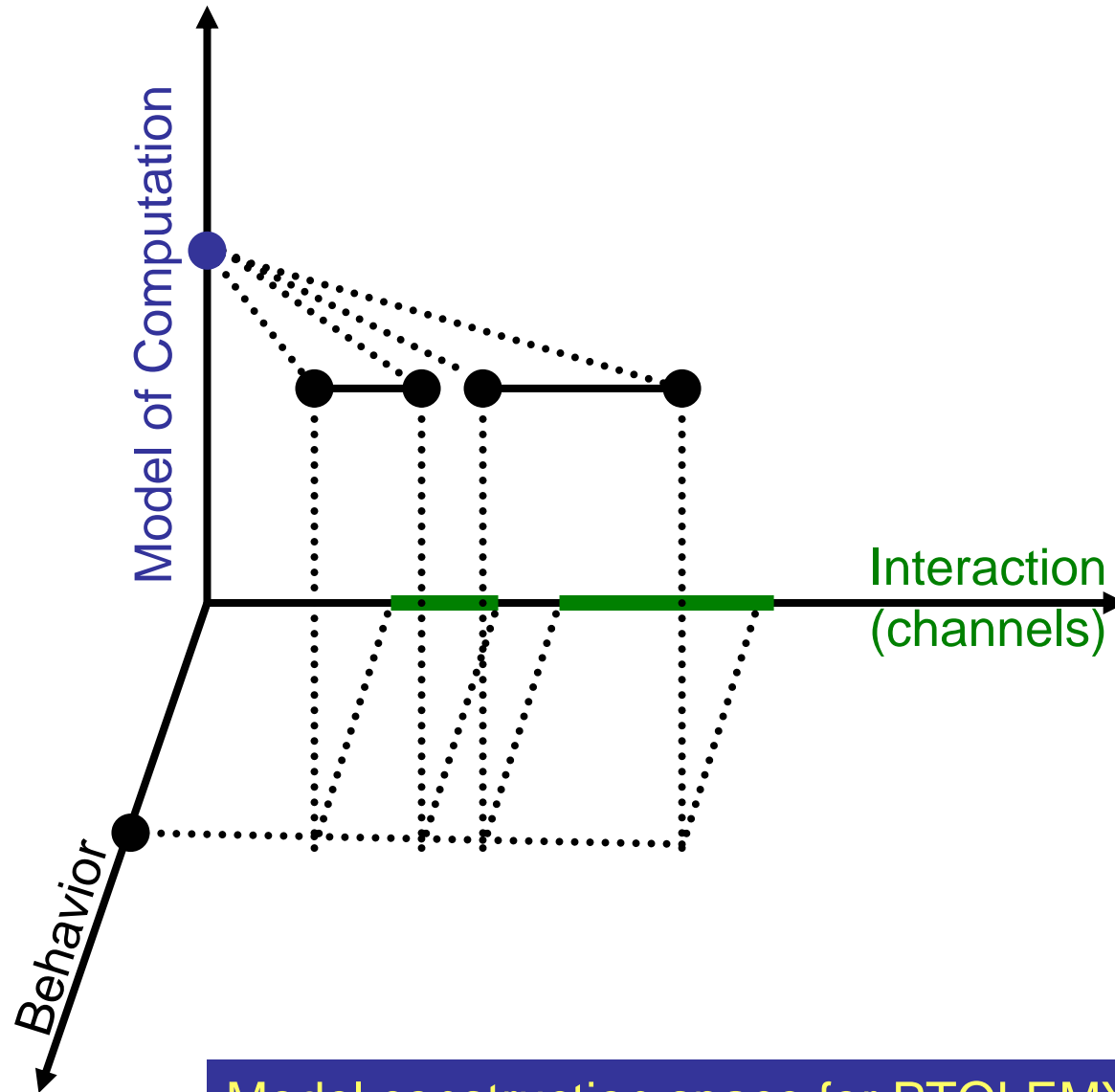


# Modeling in BIP– Model construction space



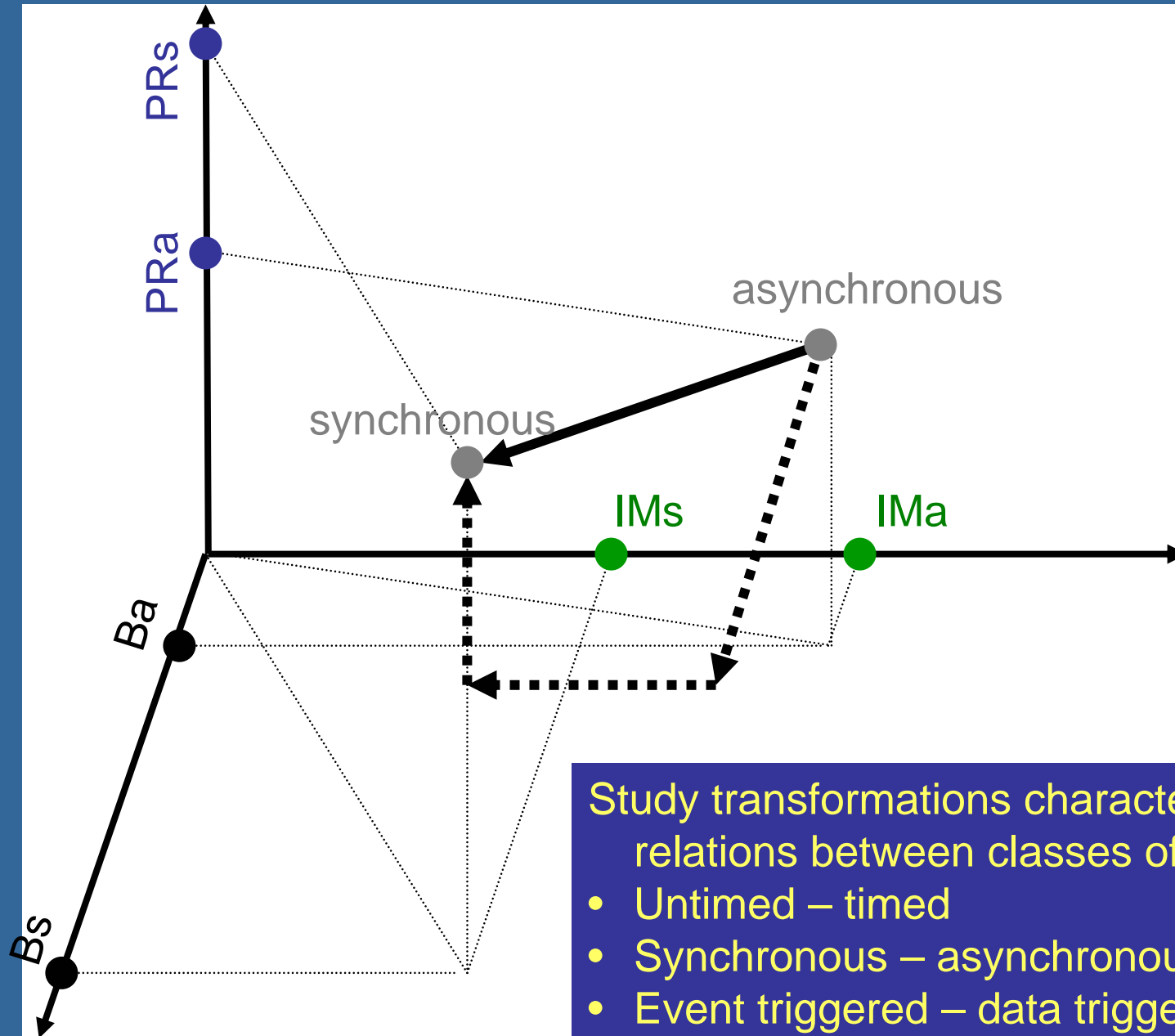
A system is defined as a point of the 3-dimensional space  
Full separation of concerns: any combination of coordinates defines a system

# Modeling in BIP – Model construction space (2)



Model construction space for PTOLEMY

# The BIP framework – Relating classes of components

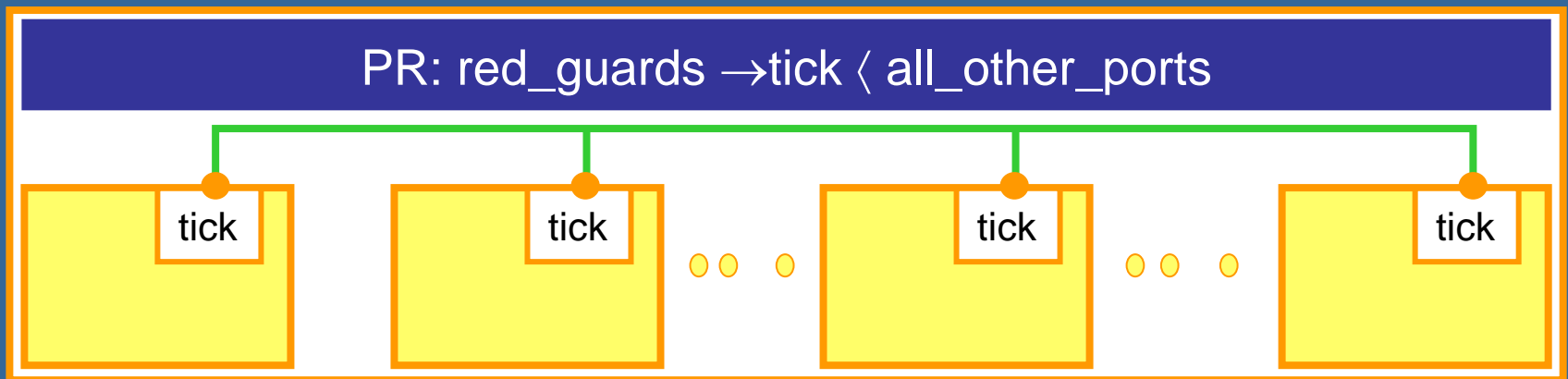
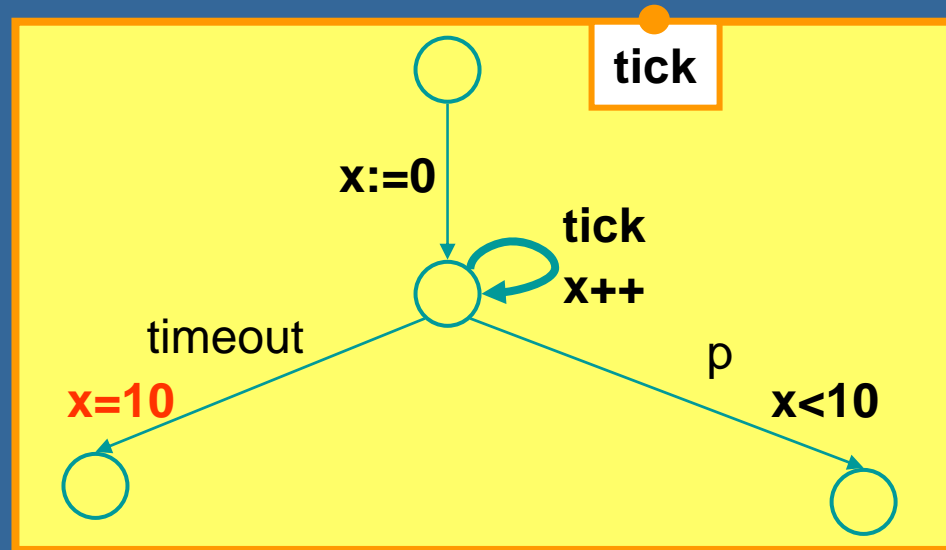


Study transformations characterizing relations between classes of systems:

- Untimed – timed
- Synchronous – asynchronous
- Event triggered – data triggered

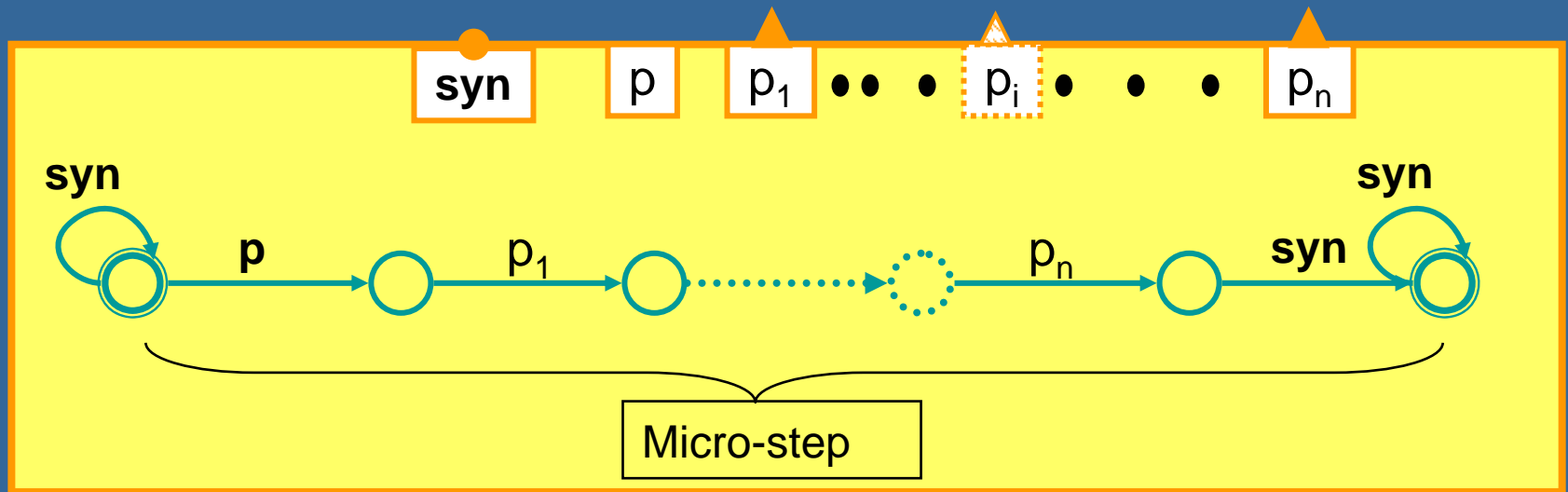
# Modeling in BIP – Timed systems

Timed Component

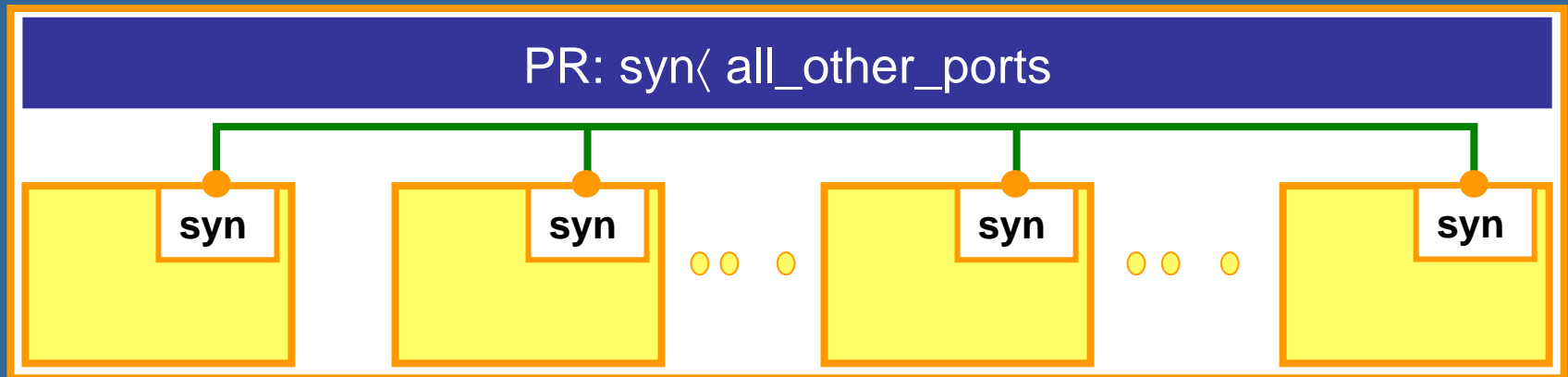


Timed architecture

# Modeling in BIP – Synchronous systems



Synchronous component



Synchronous architecture

# Modeling in BIP – MPEG4 Video encoder: Componentization

Transform a monolithic program into a componentized one

- ++ reconfigurability, schedulability
- overheads (memory, execution time)

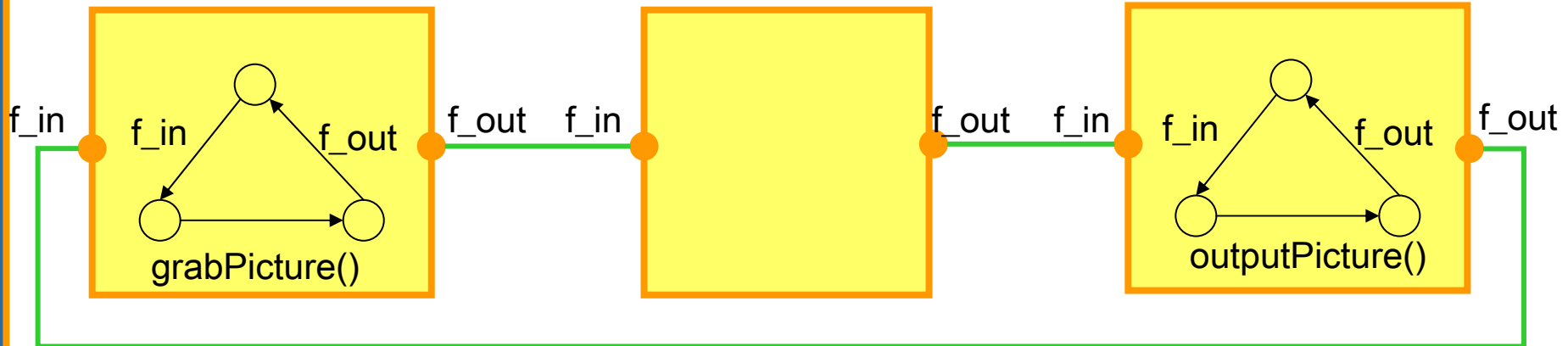
Video encoder characteristics:

- 12000 lines of C code
- Encodes one frame at a time:
  - grabPicture() : gets a frame
  - outputPicture() : produces an encoded frame

GrabPicture

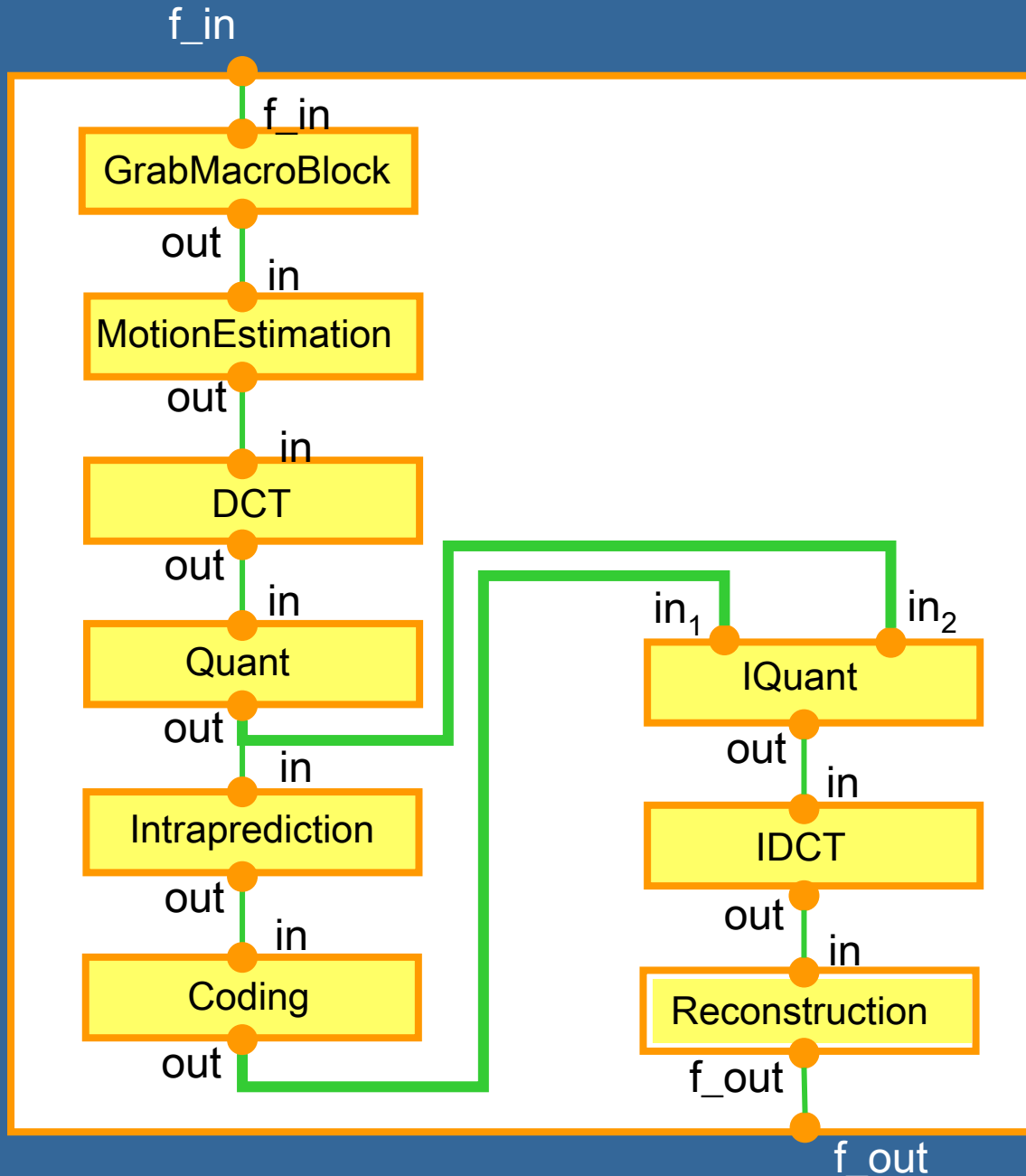
Encode

OutputPicture





# Modeling in BIP –Video encoder: The Encode component

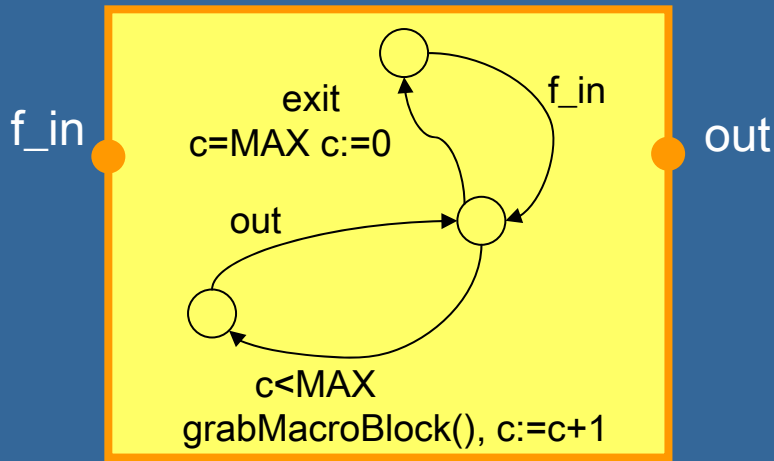


GrabMacroBlock:  
splits a frame in  
 $(W*H)/256$  macro  
blocks, outputs one  
at a time

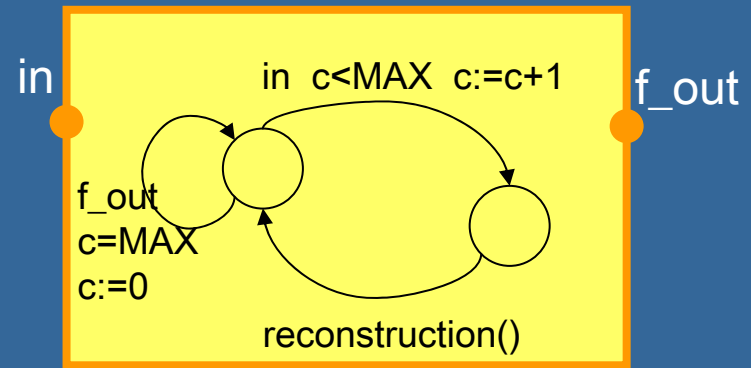
Reconstruction:  
regenerates the  
encoded frame from  
the encoded macro  
blocks.

— : buffered  
connections

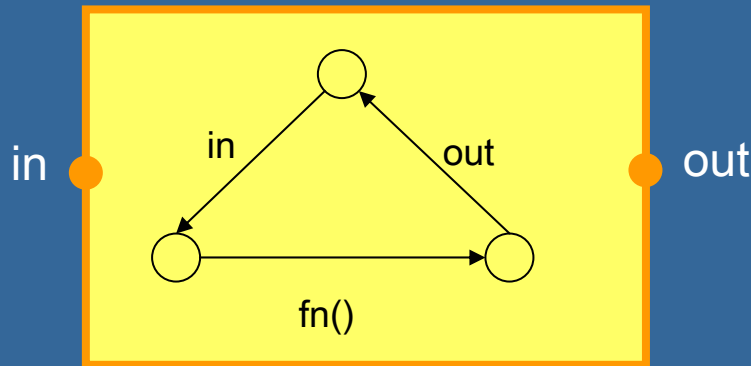
# Modeling in BIP – Video encoder : Atomic components



GrabMacroBlock



Reconstruction



Generic Functional component

$MAX = (W * H) / 256$   
 $W = \text{width of frame}$   
 $H = \text{height of frame}$

# Modeling in BIP – Video encoder: The BIP Encoder features

- BIP code describes a control skeleton for the encoder
  - Consists of 20 atomic components and 34 connectors
  - ~ 500 lines of BIP code
  - Functional components call routines from the encoder library
- The generated C++ code from BIP is ~ 2,000 lines
- The size of the BIP binary is 288 Kb compared to 172 Kb of monolithic binary.

# Modeling in BIP – Video encoder : Componentization overhead

## Overhead in execution time wrt monolithic code:

- ~66% due to communication (can be reduced by composing components at compile time)
  - function calls by atomic components to the execution engine for synchronization.
- ~34% due to resolution of non determinism (can be reduced by narrowing the search space at compile time)
  - time spent by engine to evaluate feasible interactions

**Problem: Reduce execution time overhead  
for componentized code**

# Overview

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## Discussion – The BIP framework: summary

Framework for component-based construction encompassing heterogeneity and relying on a **minimal set of constructs and principles**

Clear separation between structure (interaction +priority) and behavior

- Structure is a first class entity
- Layered description => separation of concerns => incrementality
- Correct-by-construction techniques for deadlock-freedom and liveness, based (mainly) on sufficient conditions on the structure

# Discussion - The BIP framework: Work directions (1)

## Methodology

- Modeling: BIP as a programming model, reference architectures in BIP
- Implementation techniques

## BIP toolset

- Generation of BIP models from system description languages such as SysML (IST/SPEEDS project), AADL and SystemC (ITEA/Spices project)
- Model transformation techniques in particular for code optimization
- Validation techniques
  - connection to Verimag's IF simulation/validation environment
  - specific techniques e.g. checking conditions for correctness by construction

# Discussion – The BIP framework: Work directions (2)

## Theory

- Study Component Algebras  $CA = (\mathbf{B}, GL, \oplus, \cong)$ , where
  - $(GL, \oplus)$  is a monoid and  $\oplus$  is idempotent
  - $\cong$  is a congruence compatible with operational semantics
- Study notions of **expressiveness** characterizing structure: Given two component algebras defined on the same set of atomic components,  
**CA1 is more expressive than CA2**  
if  $\forall P \exists gl2 \in GL2 \text{ } gl2(B1, \dots, Bn) \text{ sat } P \Rightarrow \exists gl1 \in GL1. \text{ } gl1(B1, \dots, Bn) \text{ sat } P$
- Model transformations
  - relating classes of systems
  - preserving properties
- Distributed implementations of BIP



## More about BIP:

- <http://www-verimag.imag.fr/index.php?page=tools>
- Email to [Joseph.Sifakis@imag.fr](mailto:Joseph.Sifakis@imag.fr)

THANK YOU

# Implementation – the BIP language: atomic component

```
component C  
port complete: p1, ... ; incomplete: p2, ...  
data {# int x, float y, bool z, .... #}  
init {# z=false; #}  
behavior  
    state s1  
        on p1 provided g1 do f1 to s1'  
        .....  
        on pn provided gn do fn to sn'  
  
    state s2  
        on .....  
        ....  
  
    state sn  
        on ....  
  
end  
end
```

# Implementation – the BIP language: connectors and priorities

```
connector BUS= {p, p', ... , }  
complete()  
  behavior  
    on  $\alpha_1$  provided  $g_{\alpha_1}$  do  $f_{\alpha_1}$   
    .....  
    on  $\alpha_n$  provided  $g_{\alpha_n}$  do  $f_{\alpha_n}$   
end
```

```
priority PR  
  if C1 ( $\alpha_1 < \alpha_2$ ), ( $\alpha_3 < \alpha_4$ ) , ...  
  if C2 ( $\alpha < \dots$ ), ( $\alpha < \dots$ ) , ...  
  ...  
  if Cn ( $\alpha < \dots$ ), ( $\alpha < \dots$ ) , ...
```

# Implementation – the BIP language: compound component

**component name**

**contains c\_name1 i\_name1(par\_list)**

.....

**contains c\_namen i\_namen(par\_list)**

**connector name1**

.....

**connector namem**

**priority name1**

.....

**priority namek**

**end**