



Timing analysis and validation with UML: the case of the embedded MARS bus manager

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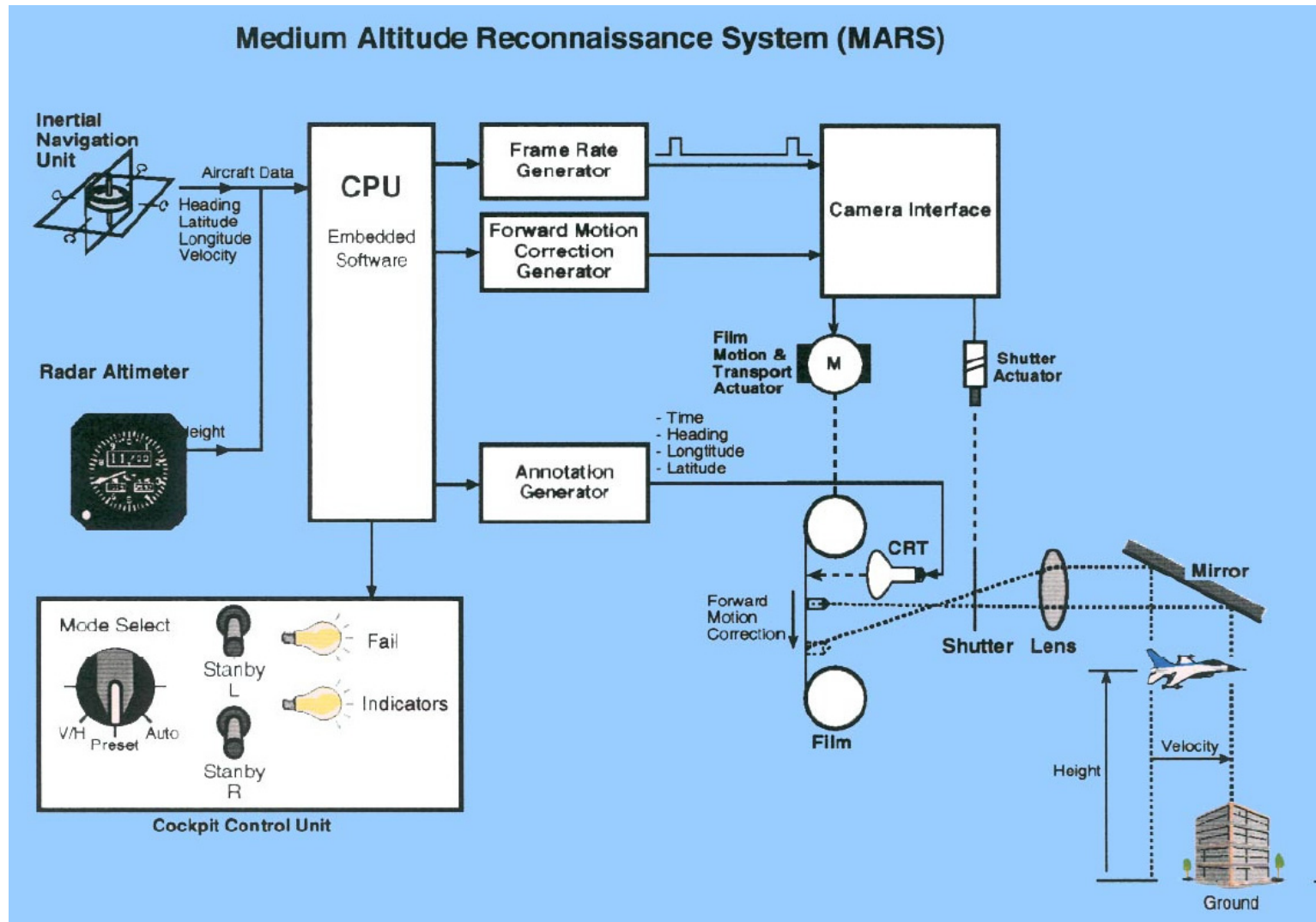




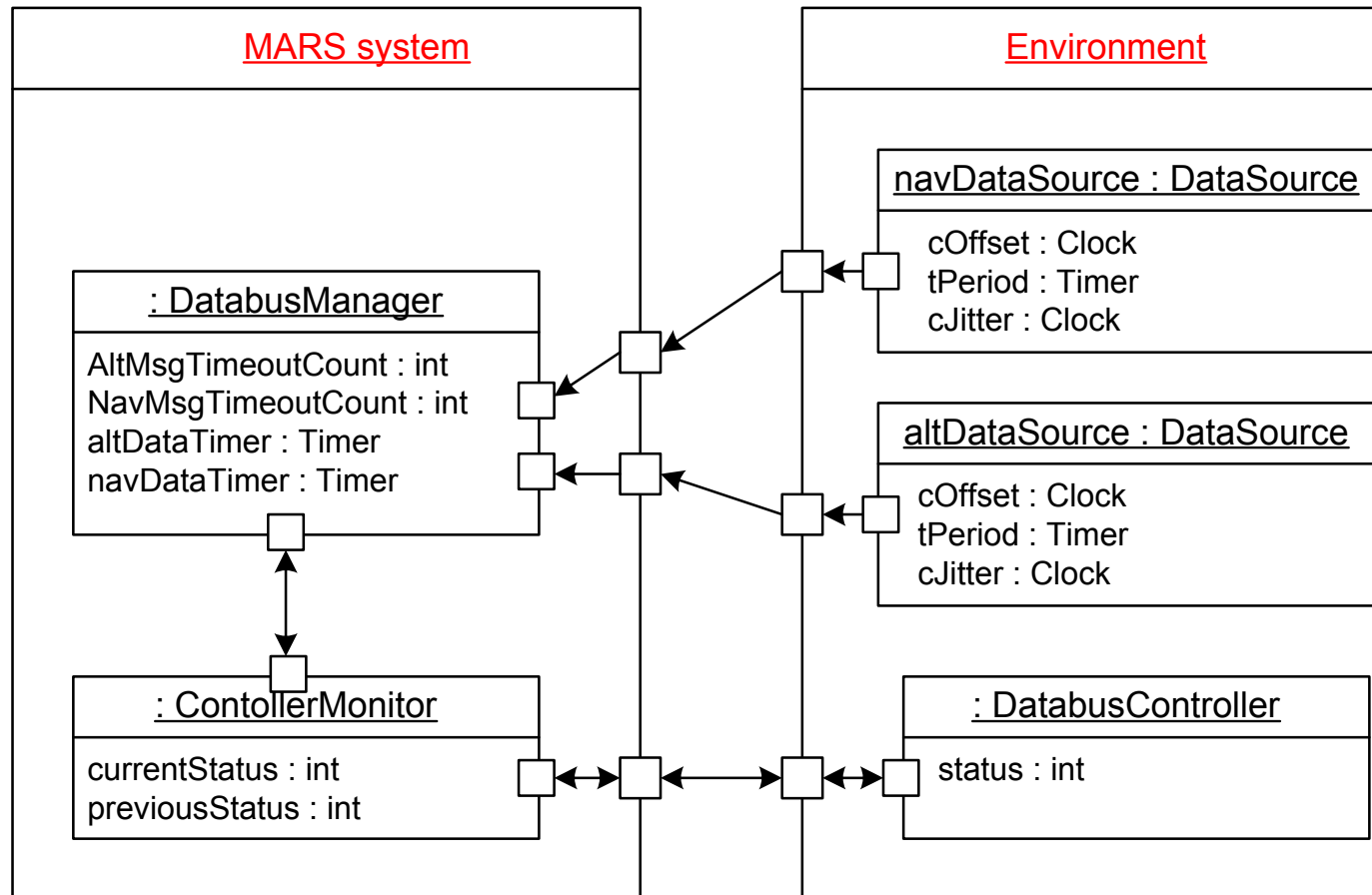
Overview

- MARS system :
 overview and requirements
- UML model and overview of tools
- Verification experiments
- Conclusions

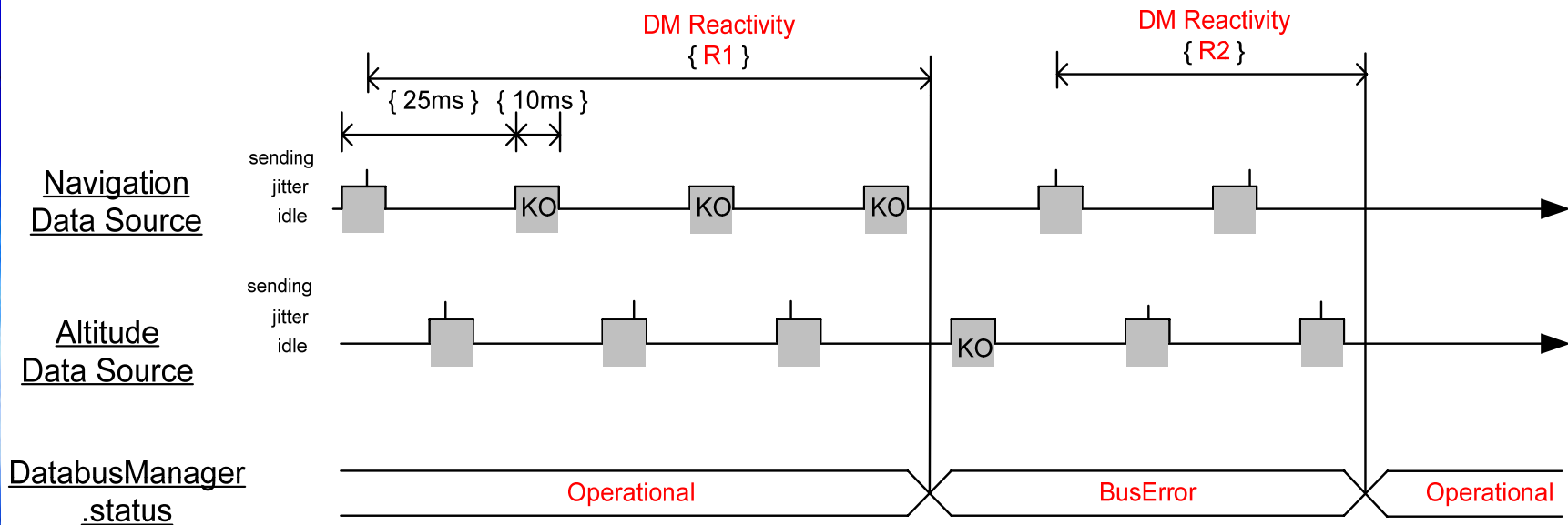
MARS overall architecture



Model architecture



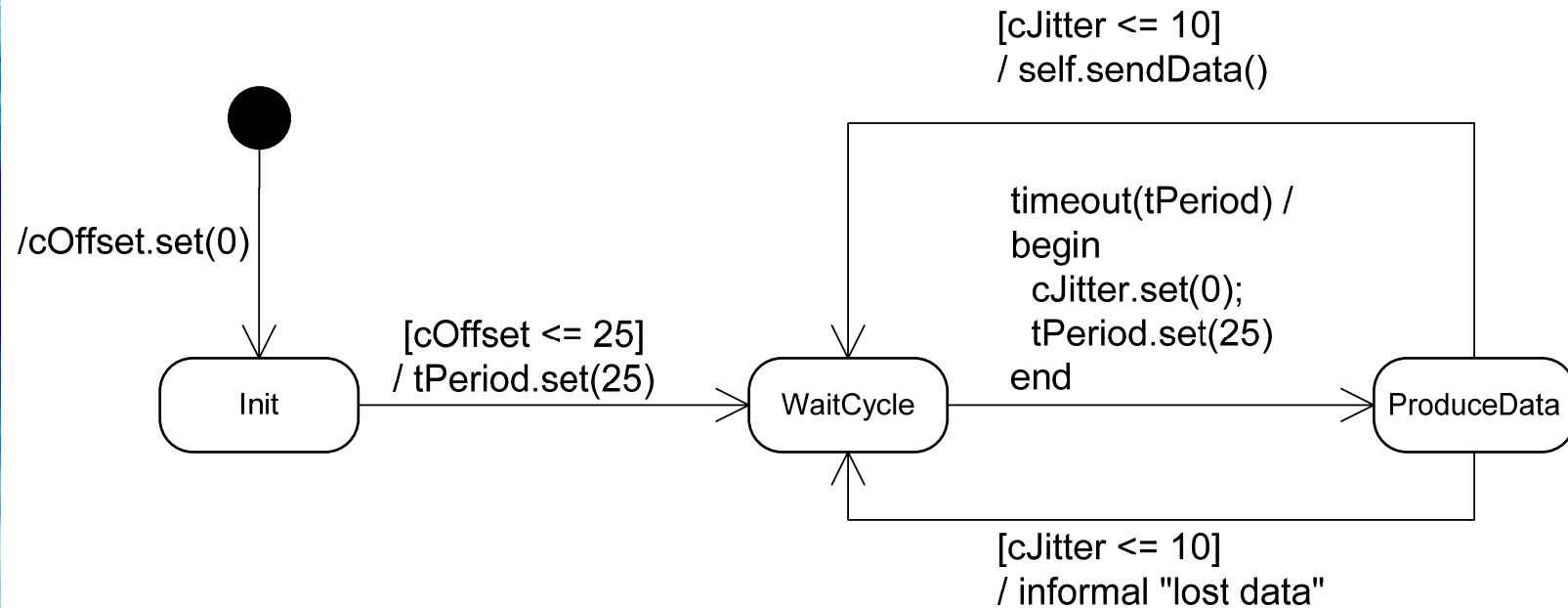
Data bus manager: timing requirements



UML and tools

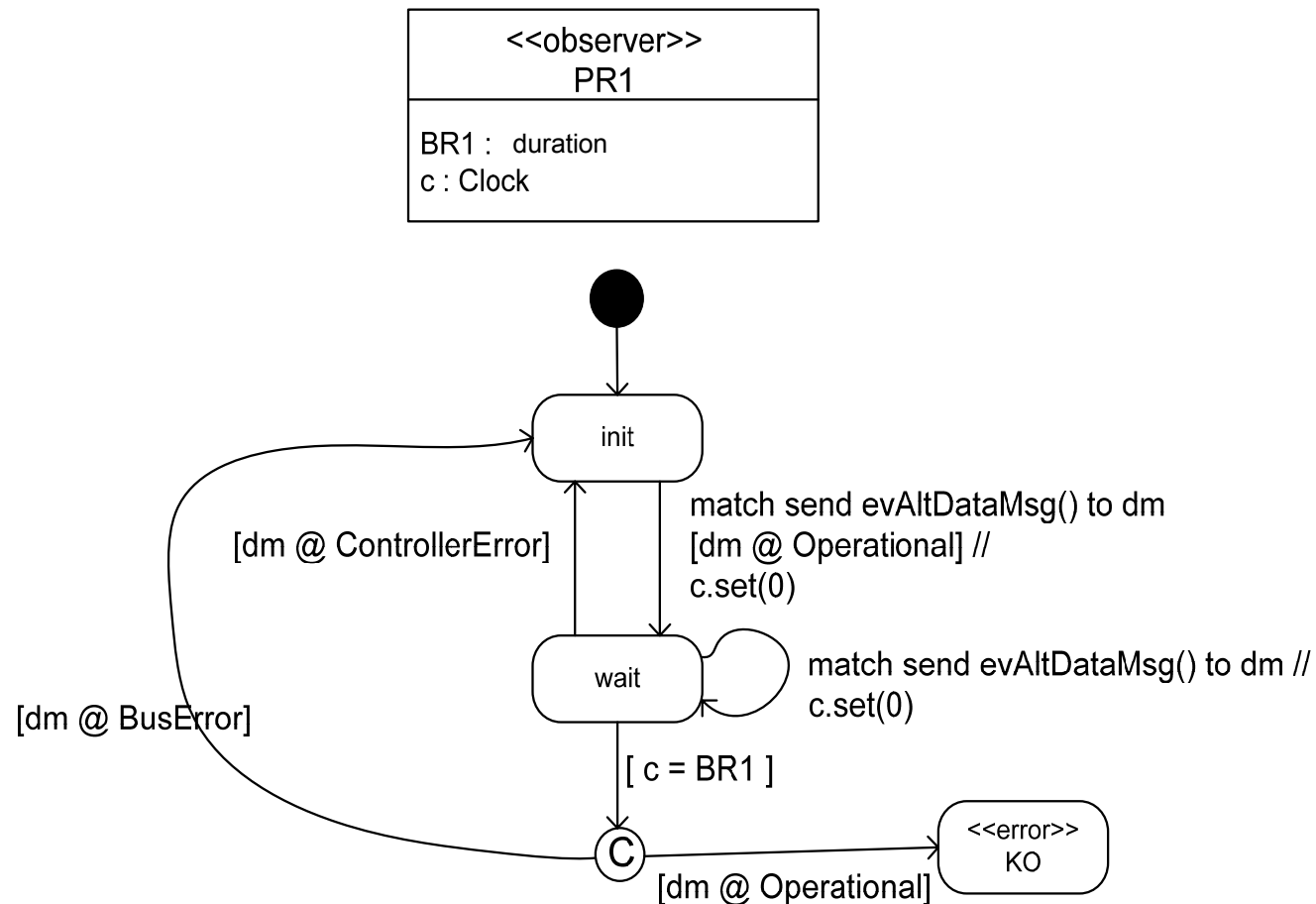
- OMEGA UML profile [\[http://www-omega.imag.fr/\]](http://www-omega.imag.fr/)
(executable semantics)
 - active components, state machines
 - asynchronous signals & synchronous calls
 - precise timing (TA w/ urgency), dynamic priority
 - safety properties : observers
- IFx toolset [\[http://www-if.imag.fr/\]](http://www-if.imag.fr/)
 - simulation
 - model checking
 - test generation,...

Environment model : data sources

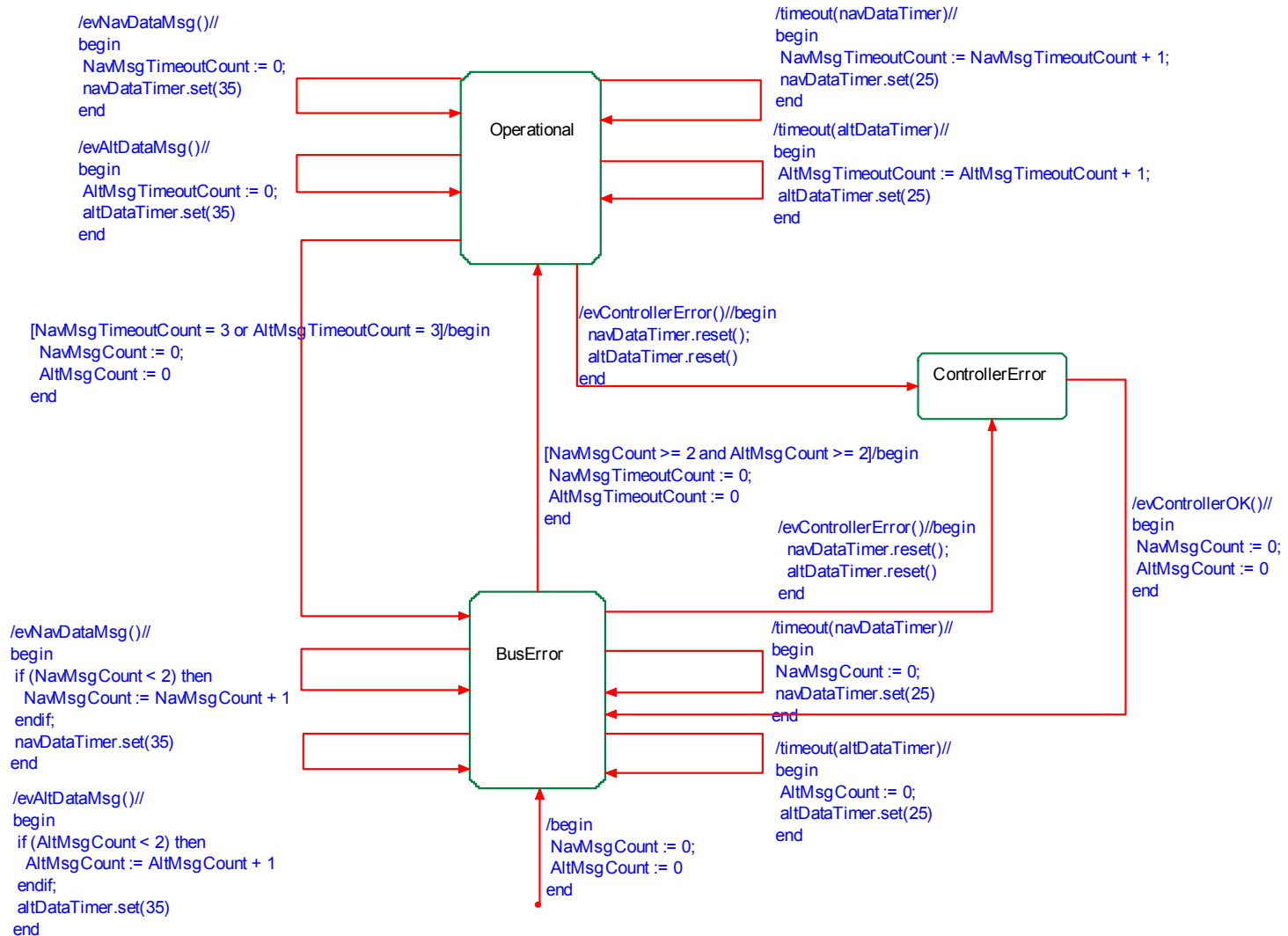


property : measuring reactivity R1

Absence of transmission is discovered within delay BR1

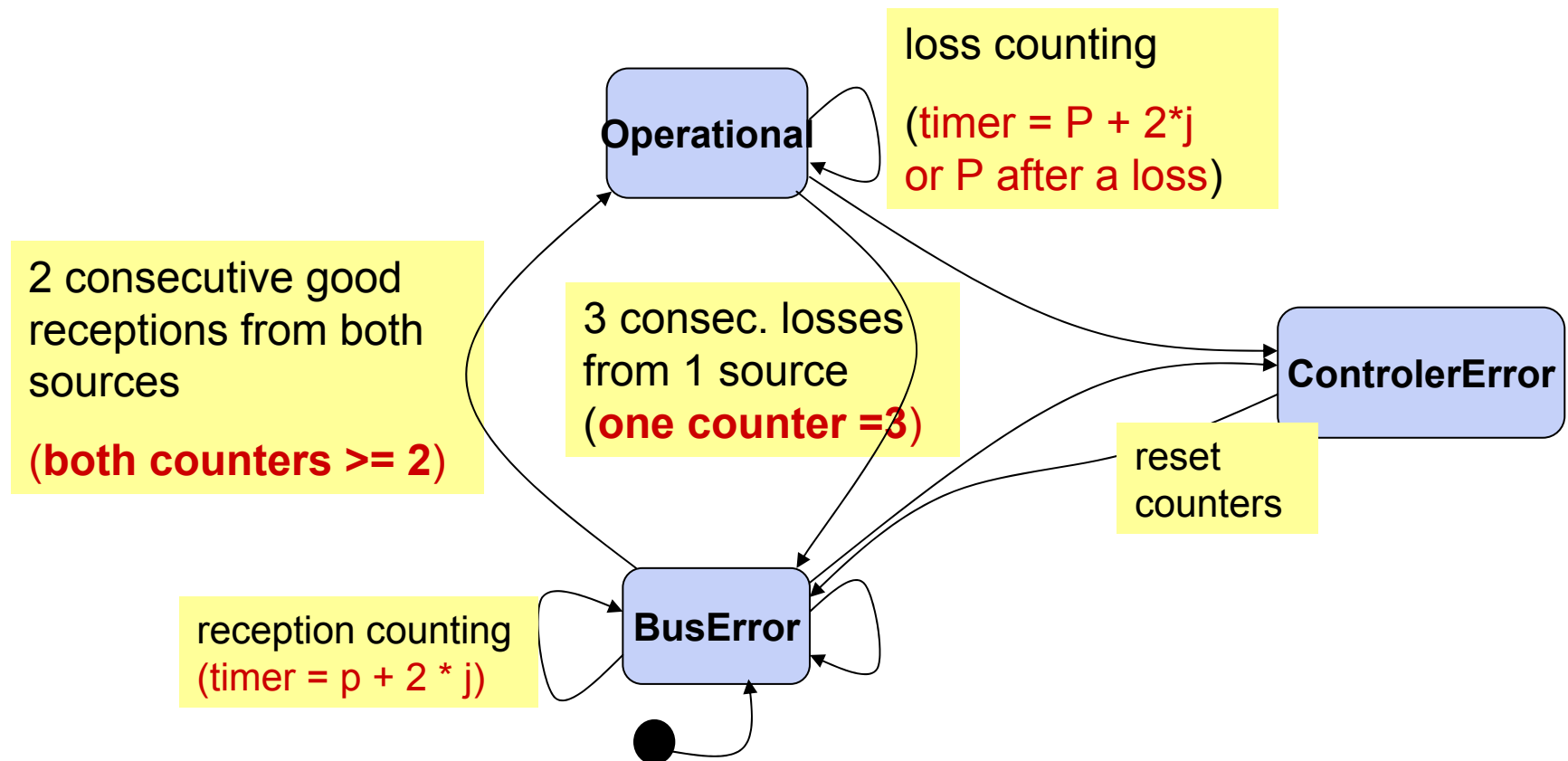


Model of the Databus Manager



Abstract model, Version 1: for every sender

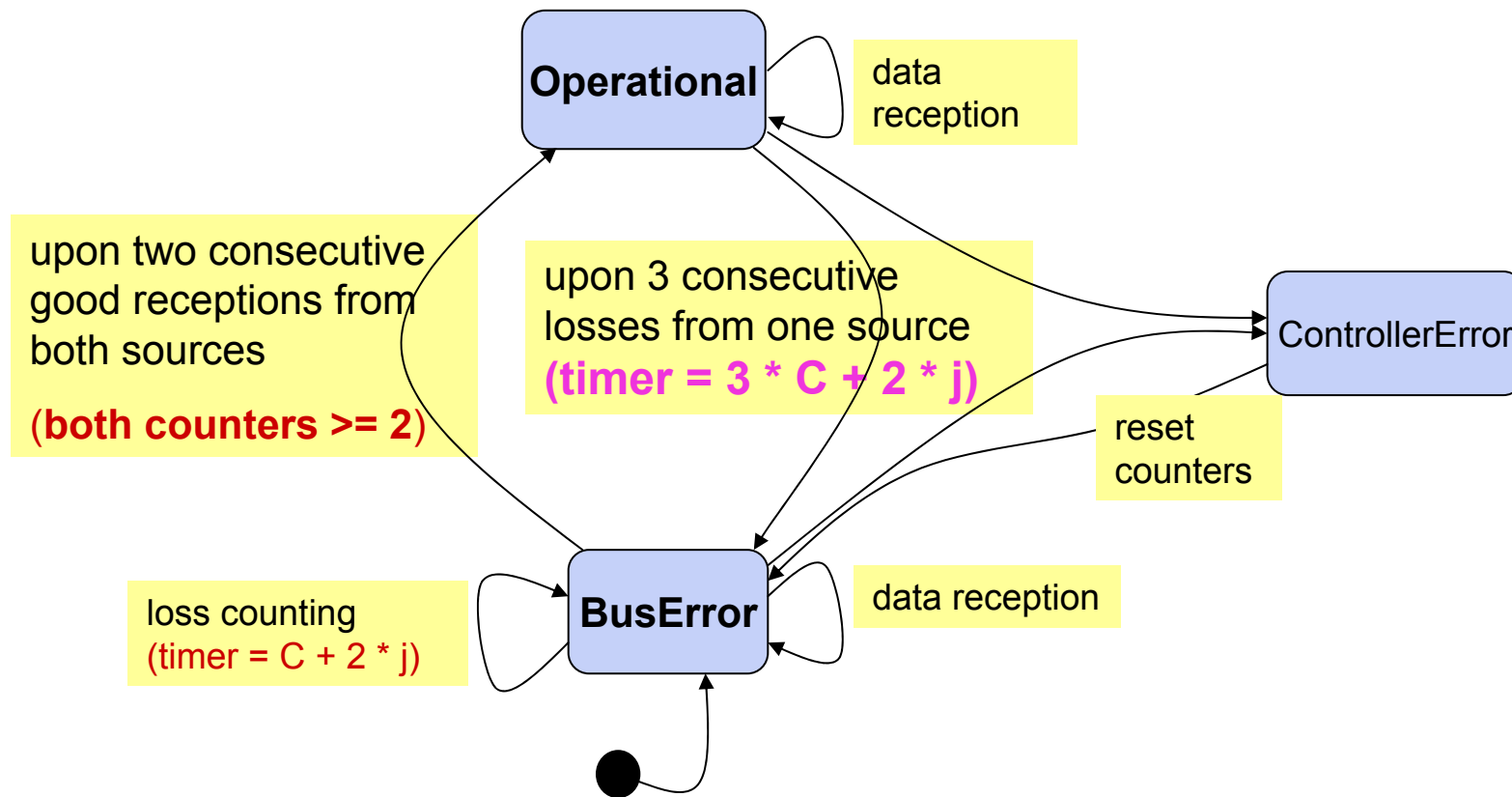
- Determine for every period “signal ok” or “signal loss” and count them
 - in state operational: count consecutive losses
 - in state BusError: count consecutive receptions
- Depending on counter values decide if status switch is needed



Verification shows: reactivity property holds

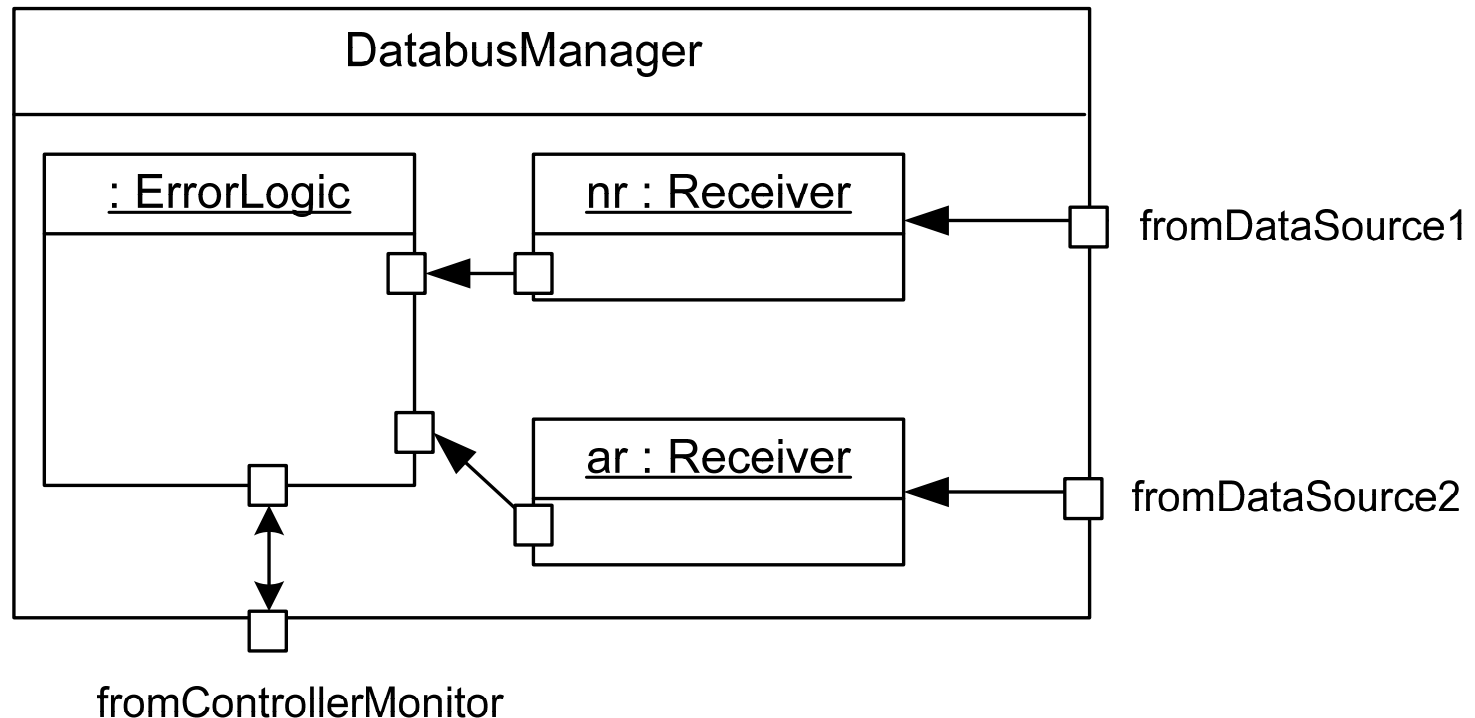
Version 2: Asymmetric

- In state **operational**, look for a “long timeout” ($3 \cdot C + 2 \cdot j$) – no need for counter
- in state **BusError**, no change



Verification shows: this version is less reactive

compositional re-design ⇒ better abstraction results



- Error logic parametric in the number of receivers
- Project property to a single receiver
- Consider 1 concrete receiver and abstract all others

Verification parameters and results

Configuration	Number of states	Number of transitions	User time
Initial model with only one source (no <i>CM</i> polling) (<i>non-conservative</i>)	1084	1420	< 1s
Initial model with two synchronized sources (no <i>CM</i> polling) (<i>non-conservative</i>)	99355	151926	36s
Initial model with two de-synchronized sources (no <i>CM</i> polling) (<i>conservative</i> – does not finish)	> 1136768	> 1676126	> 9m30s
Abstract model, 10ms <i>CM</i> polling (<i>conservative</i> – does not finish)	> 1494864	> 701120	> 8m12
Abstract model with no <i>CM</i> polling (<i>non-conservative</i>)	118690	174871	45s
Abstract model with lazy <i>CM</i> polling (<i>conservative</i>)	155166	263368	1m21s

Conclusions

- Push button verification for whole systems –
not there yet !
- Still useful for analysis of "hard points"
- Essential success factors :
 - "user-friendly" language and tools
 - use of abstractions
 - easy ones (~ automatic) : timing relaxation, ...
 - harder ones : design refactoring ← "design for verification" skills