

TuBound – A Conceptually New Tool for Worst-Case Execution Time Analysis¹

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Requirements for WCET analysis tools

WCET analysis

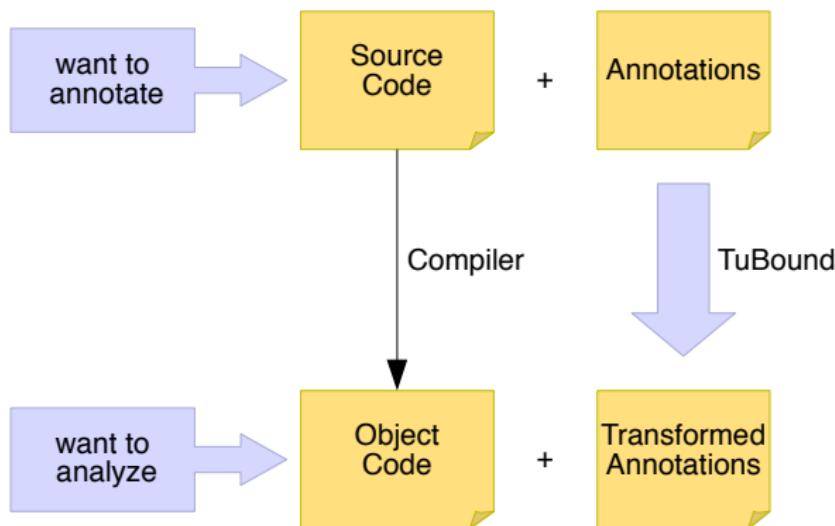
Depends on manual annotations

- Results from static analysis may not be sufficient
- User may have additional *domain-specific knowledge*

WCET analysis

Must be performed on the final low-level machine code

How to bring both worlds together



- ➡ Bridge the gap between desired annotation level and analysis level

TuBound is a WCET analysis tool developed at TU Vienna.

New Concepts

- Source-based Annotations
- Source-based Optimization
- FlowTrans to transform flow constraints alongside the compiler

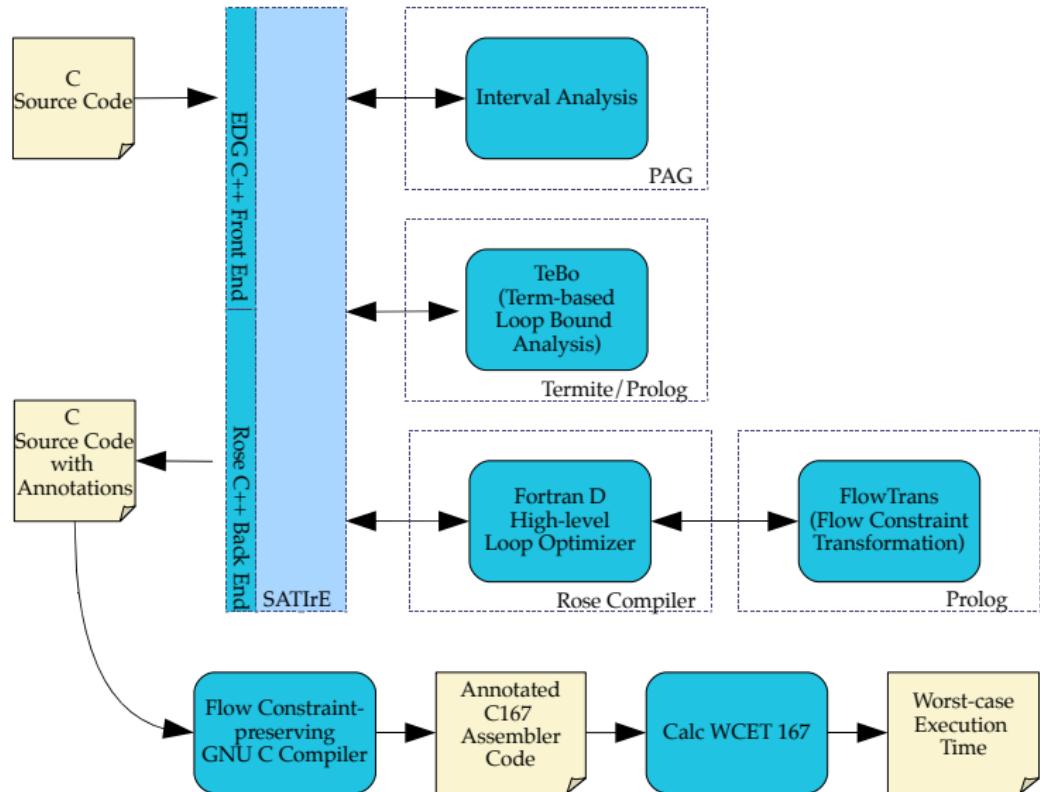
Architecture Overview

Three Phases

Start-up and Annotation

Program Optimization and WCET Annot. Transformation

Compilation and WCET calculation



I. Start-up and Annotation 1/2

1. Parsing

- EDG C++ front end
- Output: Rose abstract syntax tree (AST)

2. Interval Analysis

- SATIrE ICFG
- Program analyzer generator (PAG)
- Interprocedural interval analysis
- Output: Variable value ranges attached to the AST

I. Start-up and Annotation 2/2

3. Loop Bound Analysis (TeBo)

- Based on the SATIrE term representation of the AST
 - Implemented in Prolog
 - Solves equation system
 - Constraint logic programming
- ➡ Output: Flow constraint annotations (as #pragma nodes) in the AST

II. Program Optimization and WCET Annotation Transformation

1. Fortran D loop optimizer

- Part of the Rose compiler
 - Source-to-source loop optimizations
- Output: Optimized program and optimization trace

2. FlowTrans

- Flow constraint transformation rules
- Output: Annotated, optimized program

III. Compilation and WCET Calculation

1. Compilation to Assembler Code

Flow constraint-aware Compiler:

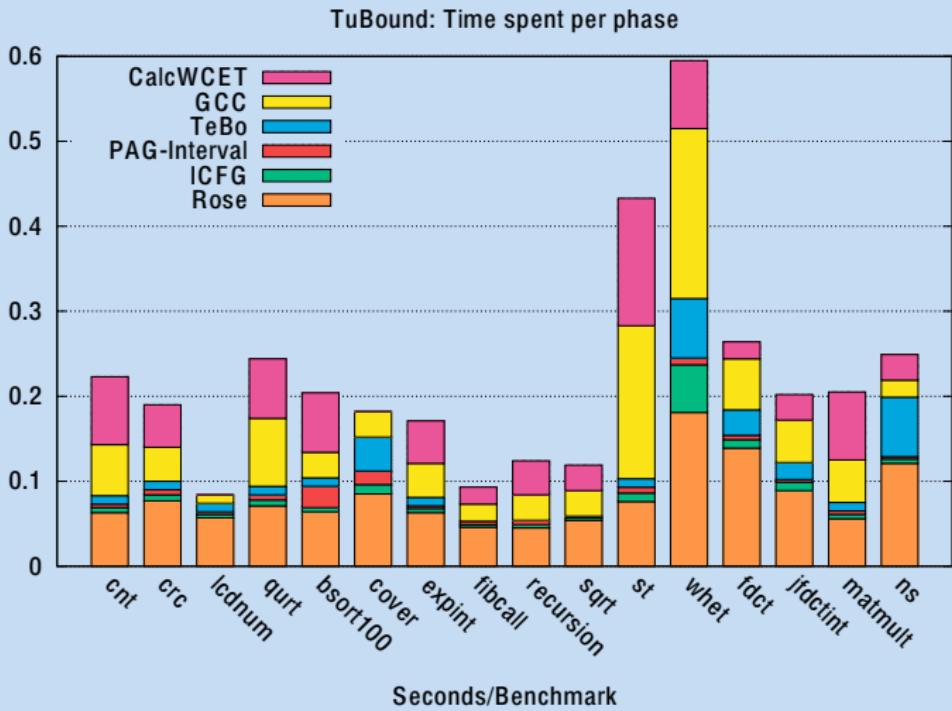
- Customized GCC 2.7 for the Infineon C167 microcontroller
 - Input: annotated C sources
 - Preserves the annotations
- Output: Annotated C167 assembler code

2. WCET Calculation (CalcWCET_{167})

- Instruction table lookup to yield WCET of basic blocks
- Combined equation system:
 - WCET of basic blocks (exec. times)
 - ICFG (connections)
 - Flow constraints (relative frequencies)
- Integer linear programming (ILP) to find maximal solution

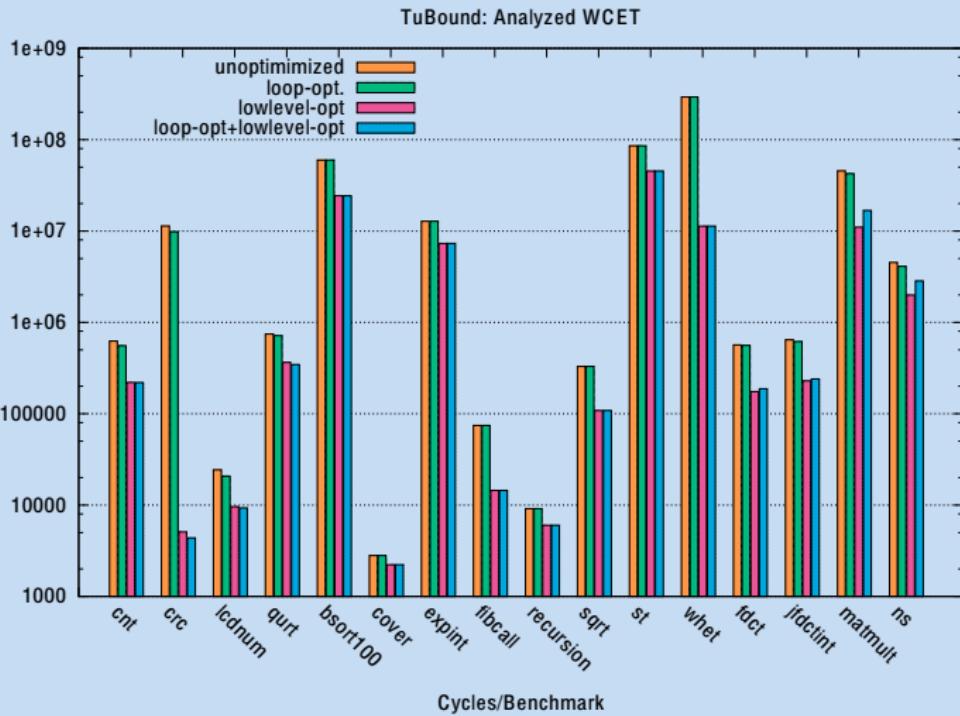
Measurements

Analysis Runtime



Measurements

Automatically calculated WCET



Conclusion

Source-based annotations

- Portable
- Easy to use
- Flexible integration of back ends
- Allow for perpetual refinement

Measurements

- Acceptable performance
- Optimizations improve WCET

➡ WCET Tool Challenge 2008

Thank You!



The CoSTA Team

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Example 1/3

Example

Original program

```
1 int sort(int* a, int n) {  
2     int i,j;  
3     for (i = 0; i < n; i++) {  
4         for (j = 0; j < i; j++) {  
5             if (a[i] > a[j]) SWAP(a[i],a[j]);  
6         }  
7     }  
8 }  
9  
10 int main()  
11 {  
12     int a[100];  
13     sort(a, 100);  
14 }
```

Example 2/3

Example

After loop bounding

```
1 int sort(int* a, int n) {  
2     int i, j;  
3     for (i = 0; i < n; i++) {  
4 #pragma wcet_loopbound(100)          13  
5         for (j = 0; j < i; j++) {  
6 #pragma wcet_loopbound(99)          14     int main()  
7             if (a[i] > a[j])           15     {  
8                 SWAP(a[i],a[j]);    16         int a[100];  
9             }                      17         sort(a, 100);  
10        }                         18     }  
11    return 0;  
12 }
```

Example 3/3

Example

After loop constraint analysis

```
1 int sort(int* a, int n) {  
2 #pragma wcet_marker(m1)  
3     int i, j;  
4     for (i = 0; i < n; i++) {  
5 #pragma wcet_marker(m2)  
6 #pragma wcet_loopedbound(100)  
7 #pragma wcet_constraint(m2=<m1*100)      18  
8     for (j = 0; j < i; j++) {  
9 #pragma wcet_marker(m3)  
10    #pragma wcet_loopedbound(99)  
11    #pragma wcet_constraint(m3=<m1*4950)      19    int main()  
12        if (a[i] > a[j])  
13            SWAP(a[i],a[j]);  
14    }  
15 }  
16 return 0;  
17 }
```

```
19    int main()  
20    {  
21        int a[100];  
22        sort(a, 100);  
23    }
```