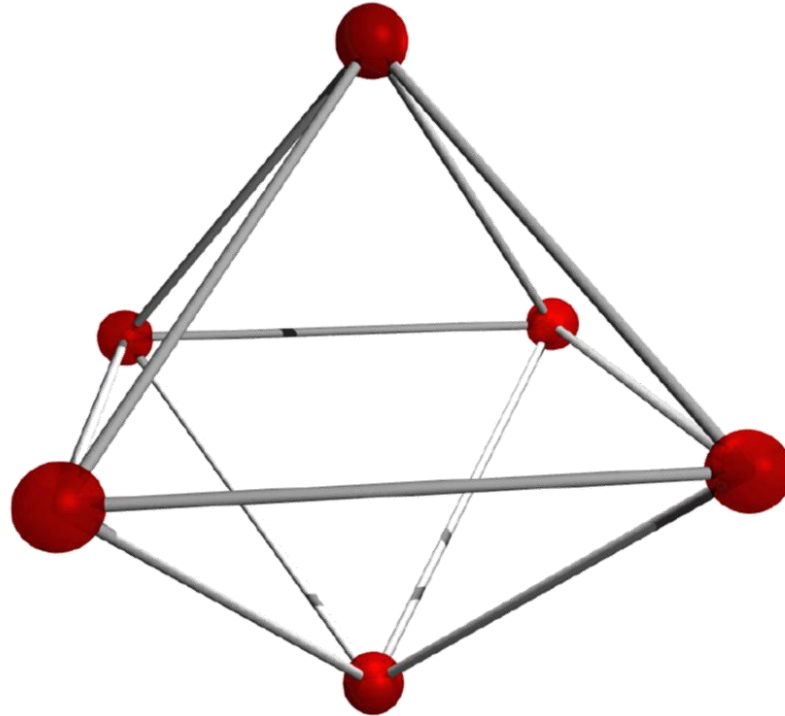
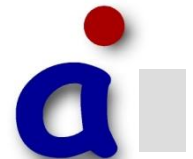


Worst-Case Timing Estimation and Architecture Exploration in Early Design Phases



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Early Estimation Problem



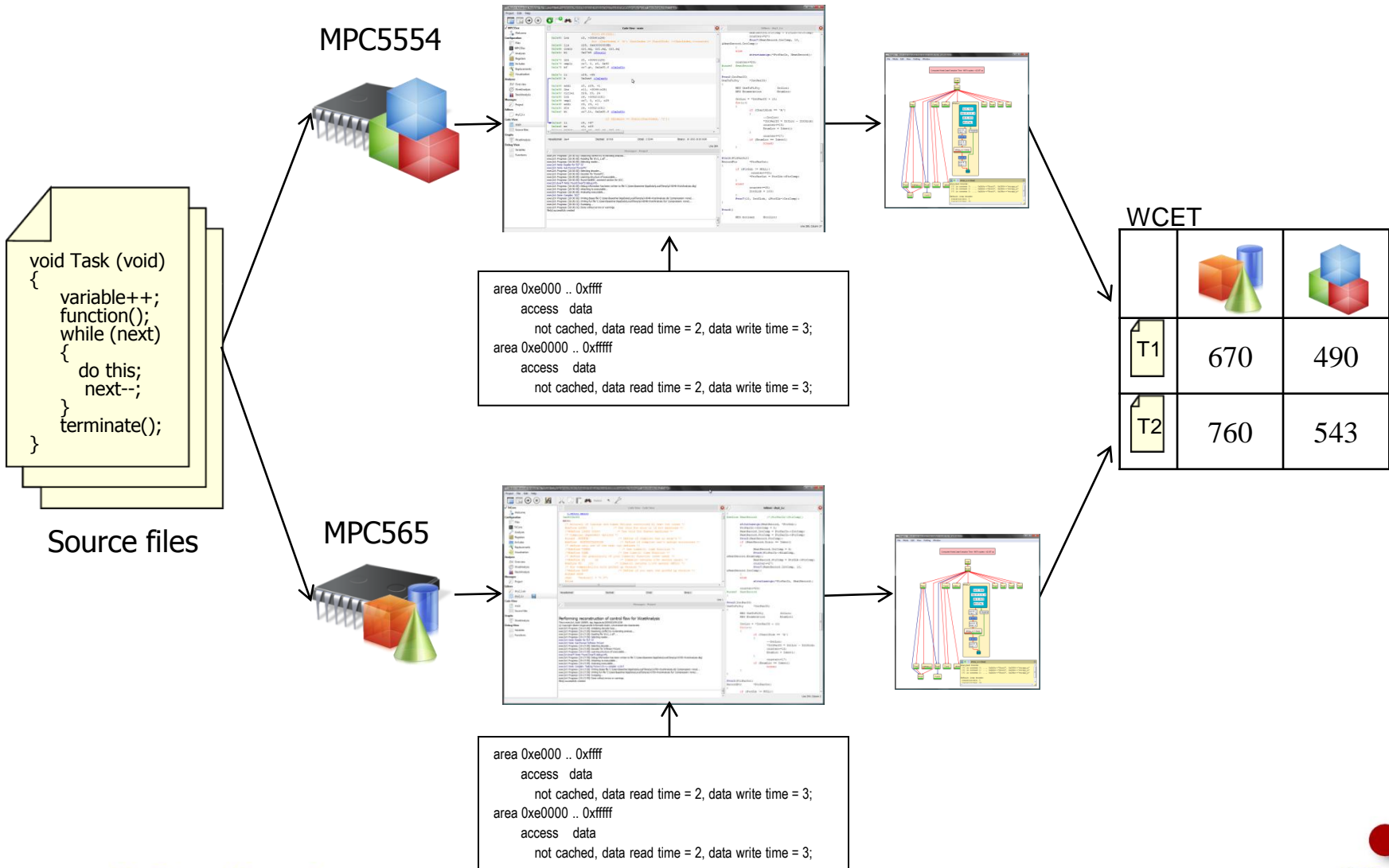
```
void Task (void)
{
    variable++;
    function();
    while (next)
    {
        do this;
        next--;
    }
    terminate();
}
```

Source code or
models

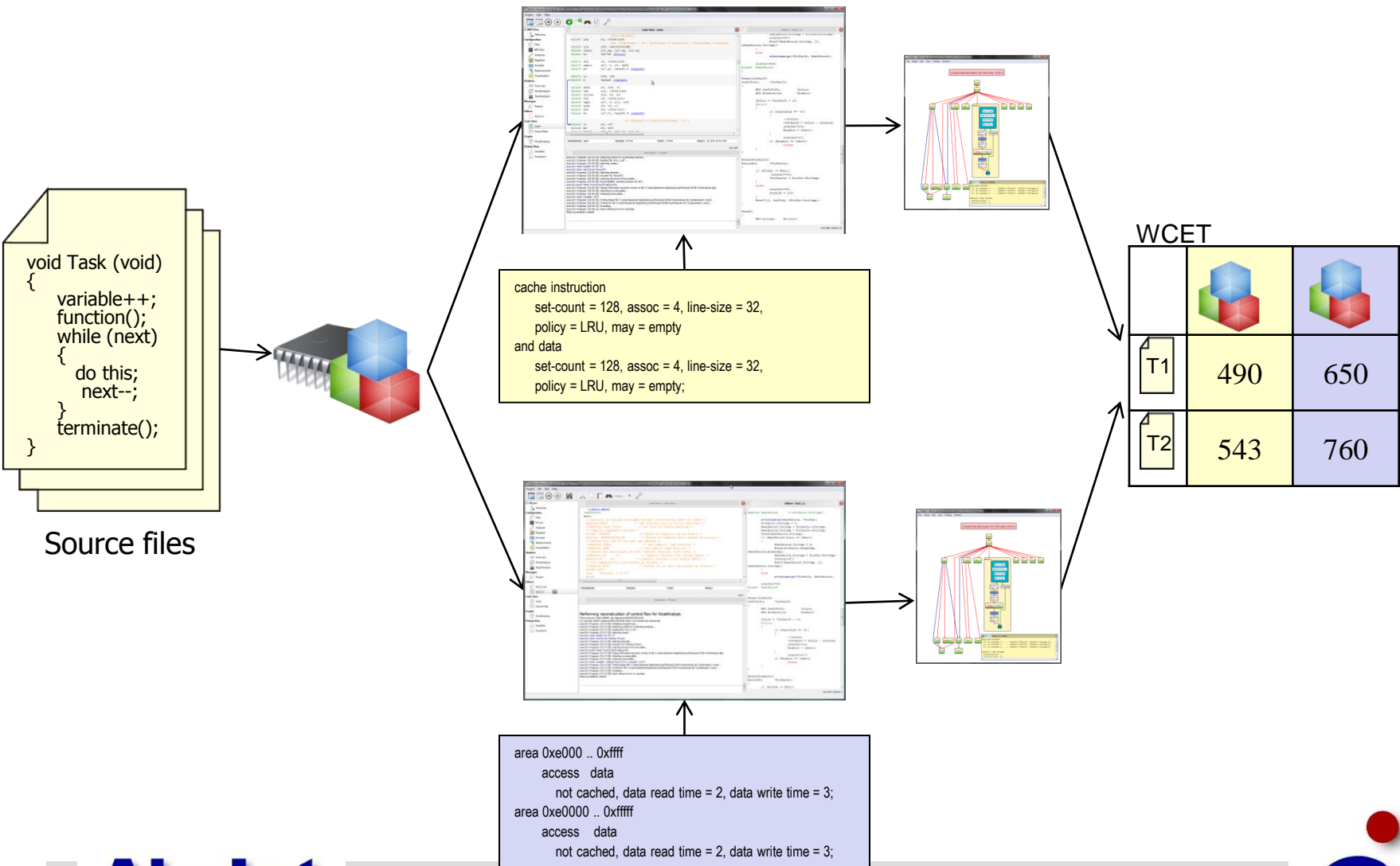
TimingExplorer

- Goal: assist in the exploration of alternative system configurations
- Design:
 - based on aiT
 - parameterizable core
- Requirements: (representative) source code

Core choice



Architecture configuration



From Source to Executable

- Compilation and linking are user's responsibility
- Compiler effects
 - Production compiler
 - Standard compiler

aiT

Application Code

```

void Task (void)
{
  variable++;
  function();
  next++;
  if (next)
    do this;
  terminate();
}

```

Software Characteristics

```

loop "_codebook" + 1 loop exactly 16 end;
recursion "_fac" max 6;
SNIPPET "printf" IS NOT ANALYZED
AND TAKES MAX 333 CYCLES;

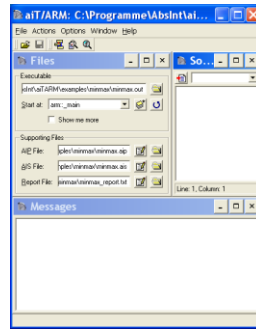
```

Compiler Linker

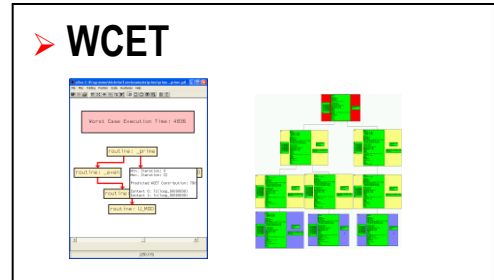
Executable (*.elf / *.out)



aiT



Entry Point 



Architecture Specification

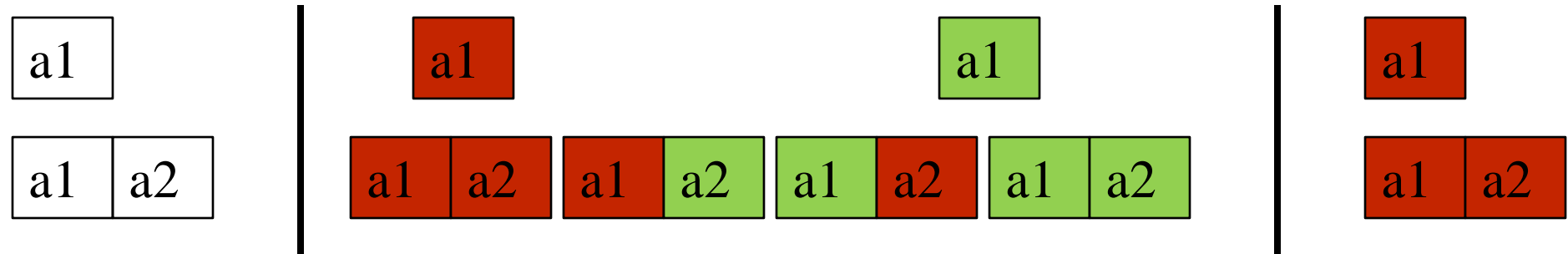
➤ Cache

- Size, line size, replacement policy, associativity
- Cache access behaviour:
 - Default strategy for unknown accesses – hit/miss
 - Turn off cache analysis and treat all accesses as hit/miss/unknown

➤ Completely parameterizable memory map

Pipeline Analysis

➤ Use local worst case



➤ Goal

- Reduce computation resource usage
- Speed up the analysis

Source-code Analyses

- Goal: reduce the need for manual annotation
- May introduce imprecision:

```
for (int i=0; i<100; i++)
```

...

- Analyses: function pointer resolution, loop bounds
 - Framework: SATIrE, LLNL-ROSE, PAG

Future work

- Case study
- Integration in a system-level analysis
 - XML Timing Cookies

Source-code Analyses

- Goal: reduce the need for manual annotation
- May introduce imprecision:

```
for (int i=0; i<100; i++)
```

...

- Analyses: function pointer resolution, loop bounds
 - Framework: SATIrE, LLNL-ROSE, PAG

Function Pointer Resolution

- Compute a points-to set for each function pointer variable
- Handles complex structures
- Output: annotations

Test benchmark

Program	Version	Lines of code	Indirect calls
diction	0.7	2 037	3
grep	2.0	12 417	3
gzip	1.2.4	8 163	4
sim6890	0.1	3 290	97

Loop-bound Analysis

- Express the upper loop bound as a parametric formula
- One symbolic formula per loop
 - value analysis \Rightarrow concrete bounds

Loop-bound analysis results

- Detects and assigns a bound to all loops
- Precision (Mälardalen benchmark suite; provided precise results of the value analysis)
 - $>$ SWEET
 - \geq oRange
- Automatically generated code
- Effective for typical counter loops