

# ***A generic framework for blackbox components in WCET computation***

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# ***PLAN***

***introduction***

examples of partial analyses  
generalization of the approach  
experimentation with OTAWA  
conclusion

# ***WCET\* computation***

WCET computation by static analysis  
program control flow analysis  
architecture effects analysis  
WCET computation → IPET

IPET (Implicit Path Enumeration Technique)

WCET = maximization of  $t_i \times x_i$

under constraints of an ILP system

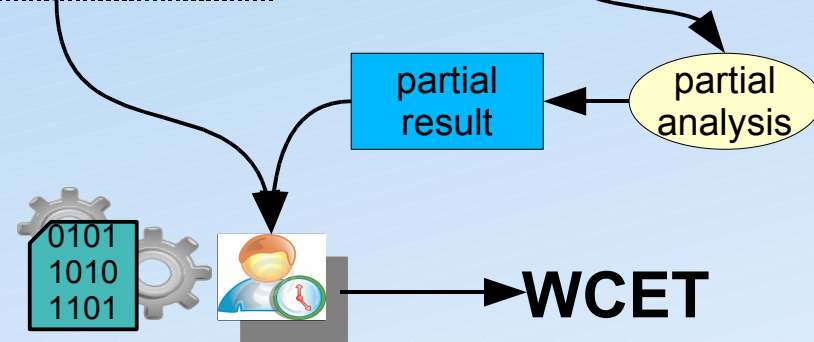
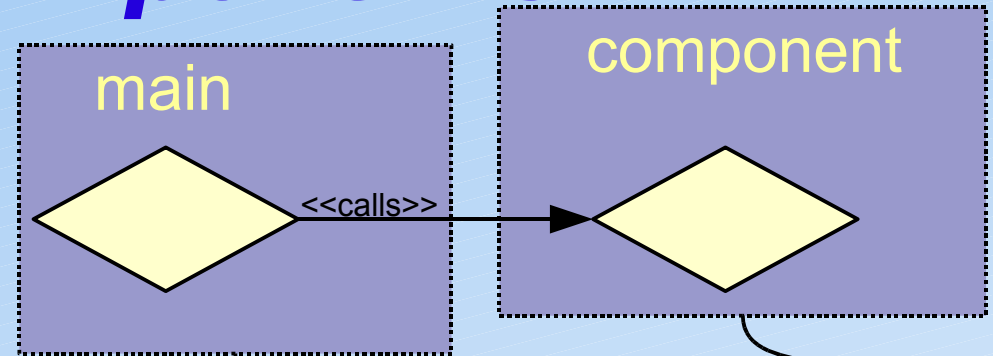
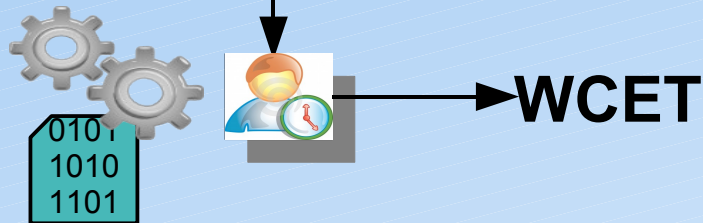
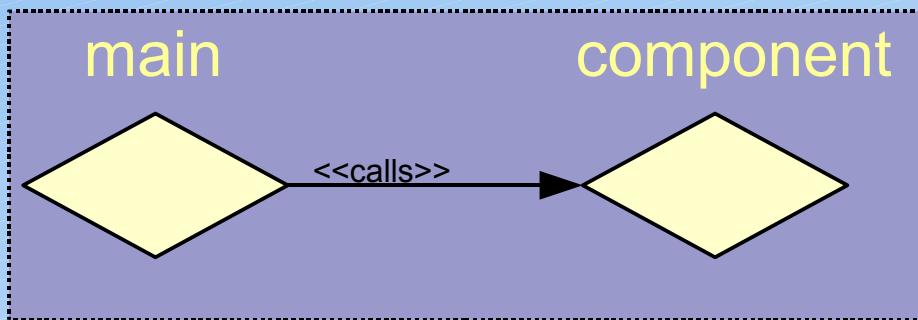
program control flow

hardware effects

ILP solver

\* Worst Case Execution Time

# *blackbox components*



WCET computation done on a whole program  
program = multiple components  
lack of informations on blackbox components  
→ no computation  
how to handle blackbox component ?  
partial analysis

\* Worst Case Execution Time

# ***PLAN***

introduction

***examples of partial analyses***

generalization of the approach

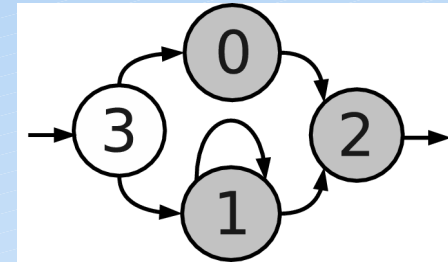
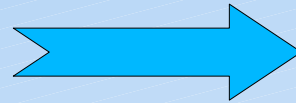
experimentation with OTAWA

conclusion

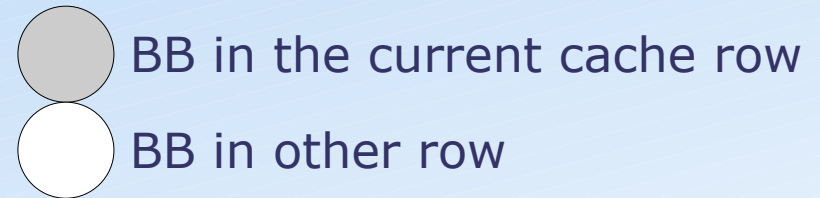
# *the instruction cache analysis*

```
#define ERR -1
#define OK 0
int status = OK
int fact (int y) {
    int result = 1;
    int i;
    if (y >= 0) {
        for (i = 1; i <= y; i++)
            result *= i;
        status = OK;
    } else status = ERR;
    return result;
}
```

**example program**



**CFG**



transfer function

$state_{\text{after}} = transfer(state_{\text{before}})$

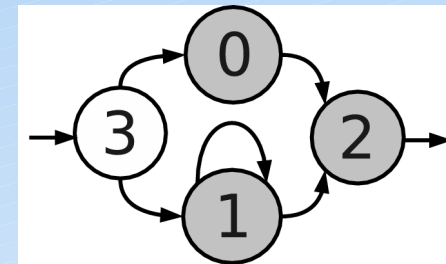
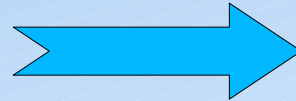
summary function

categories =  $summary(state_{\text{before}})$

# *the loop bounds estimation*

```
#define ERR -1
#define OK 0
int status = OK
int fact (int y) {
    int result = 1;
    int i;
    if (y >= 0) {
        for (i = 1; i <= y; i++)
            result *= i;
        status = OK;
    } else status = ERR;
    return result;
}
```

**example program**



**CFG**

transfer function

$state_{\text{after}} = transfer(state_{\text{before}})$

summary function

$bounds = summary(state_{\text{before}})$

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# *content of the partial result*

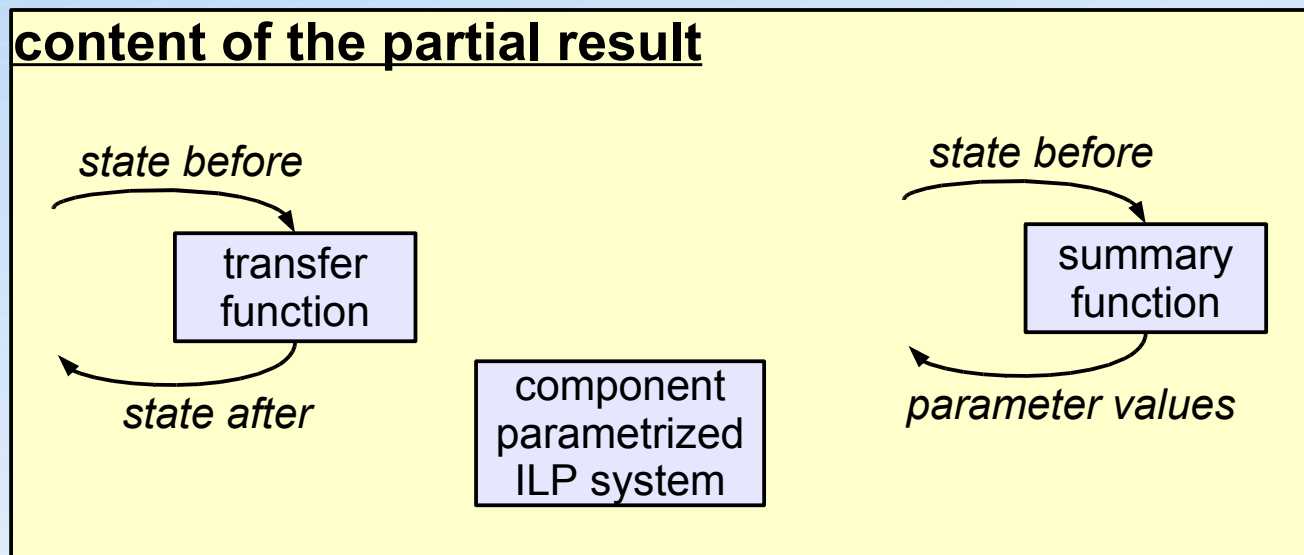
## **for each analysis**

effect : component  $\rightarrow$  main program =  
transfer function for each

effect : calling context  $\rightarrow$  ILP parameters =  
summary function

## **for the component**

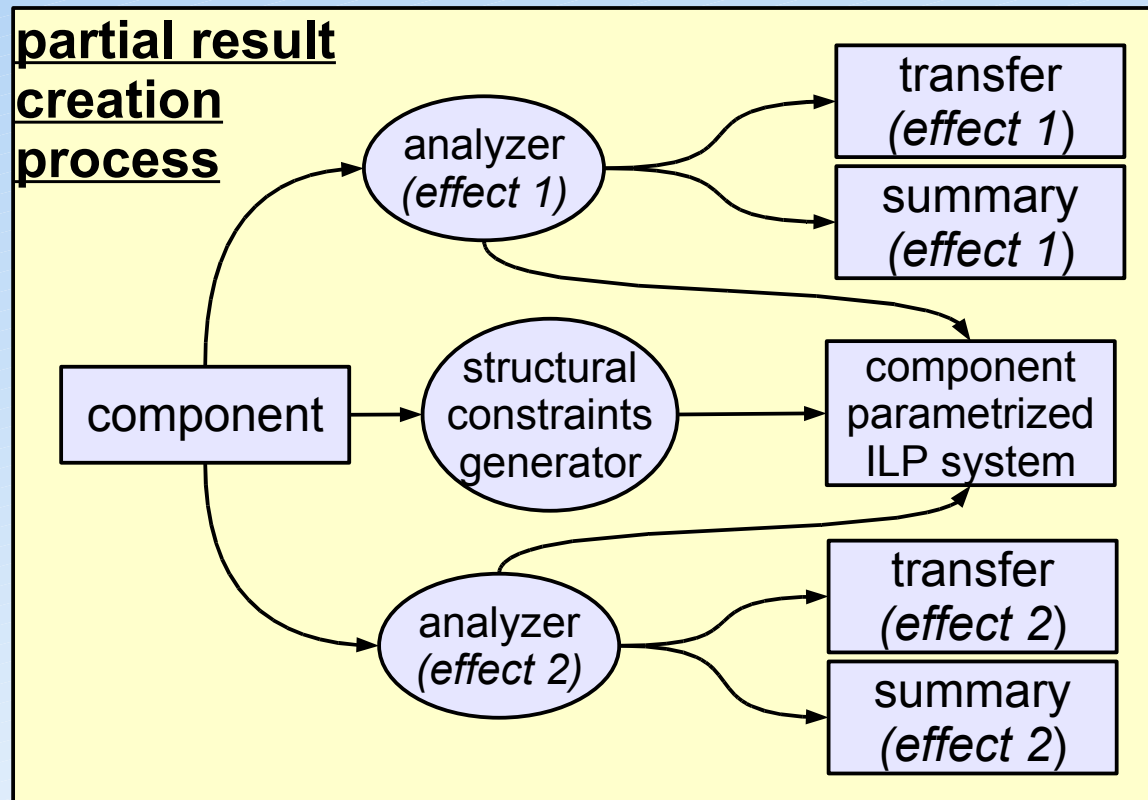
contribution of component = parametrized ILP system



# ***(1) creation of the partial result***

each analyzer → transfer  
→ summary

all analyzers → one ILP system

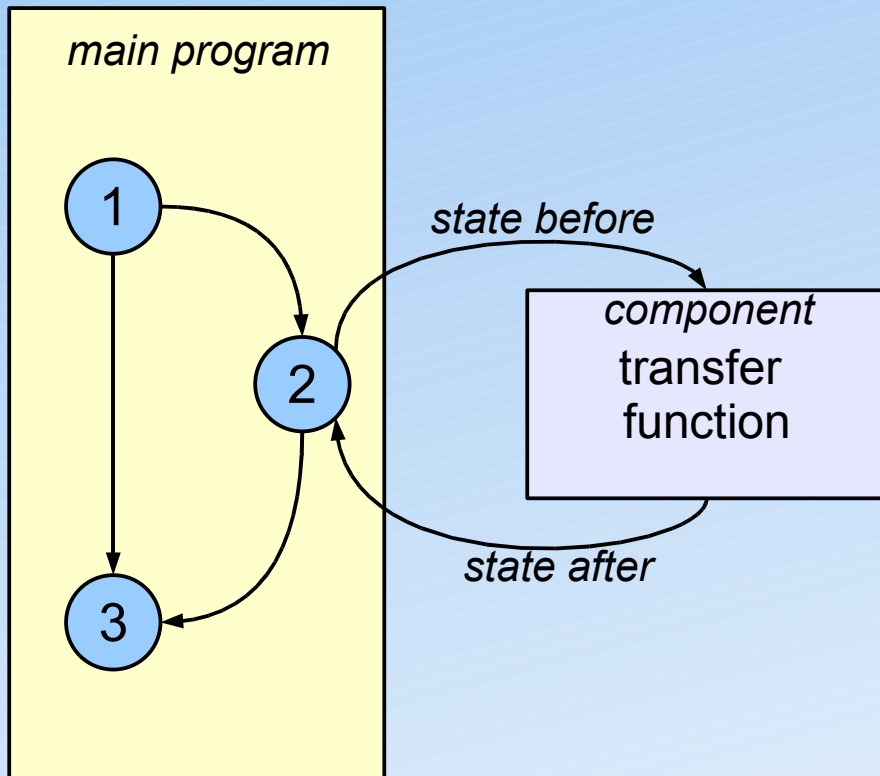


## ***(2) usage of the partial result***

with the *transfer* function

do each analysis on the  
main program

get the component entry  
state for each analysis



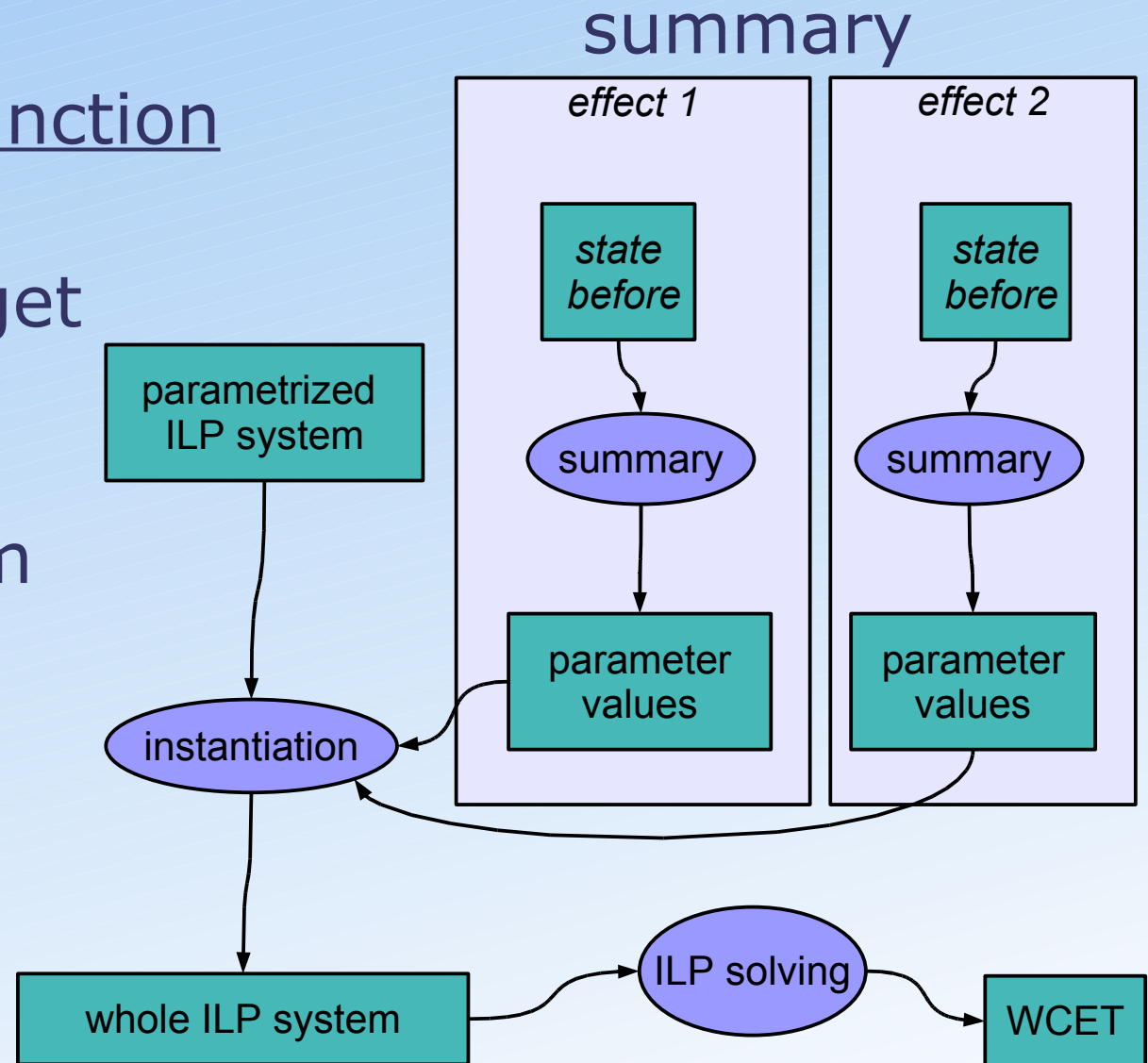
## *(2) usage of the partial result*

with the summary function

apply summaries to get  
parameter values

instantiate ILP system

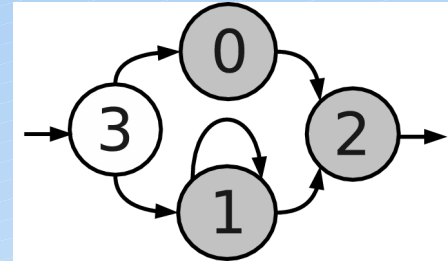
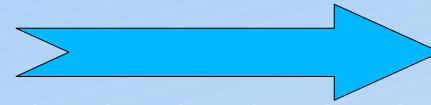
solve ILP system



# partial result example

```
#define ERR -1
#define OK 0
int status = OK
int fact (int y) {
  int result = 1;
  int i;
  if (y >= 0) {
    for (i = 1; i <= y; i++)
      result *= i;
    status = OK;
  } else status = ERR;
  return result;
}
```

## example program



## CFG

### parametrized ILP system

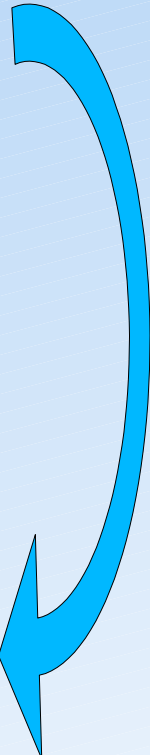
if ( $p_0 = \text{Always\_Hit}$ )  $x_0^{\text{miss}} = 0$   
if ( $p_1 = \text{Always\_Hit}$ )  $x_1^{\text{miss}} = 0$   
if ( $p_2 = \text{Always\_Hit}$ )  $x_2^{\text{miss}} = 0$   
 $e_{1,1} = p_3 \cdot e_{3,1}$

### transfer

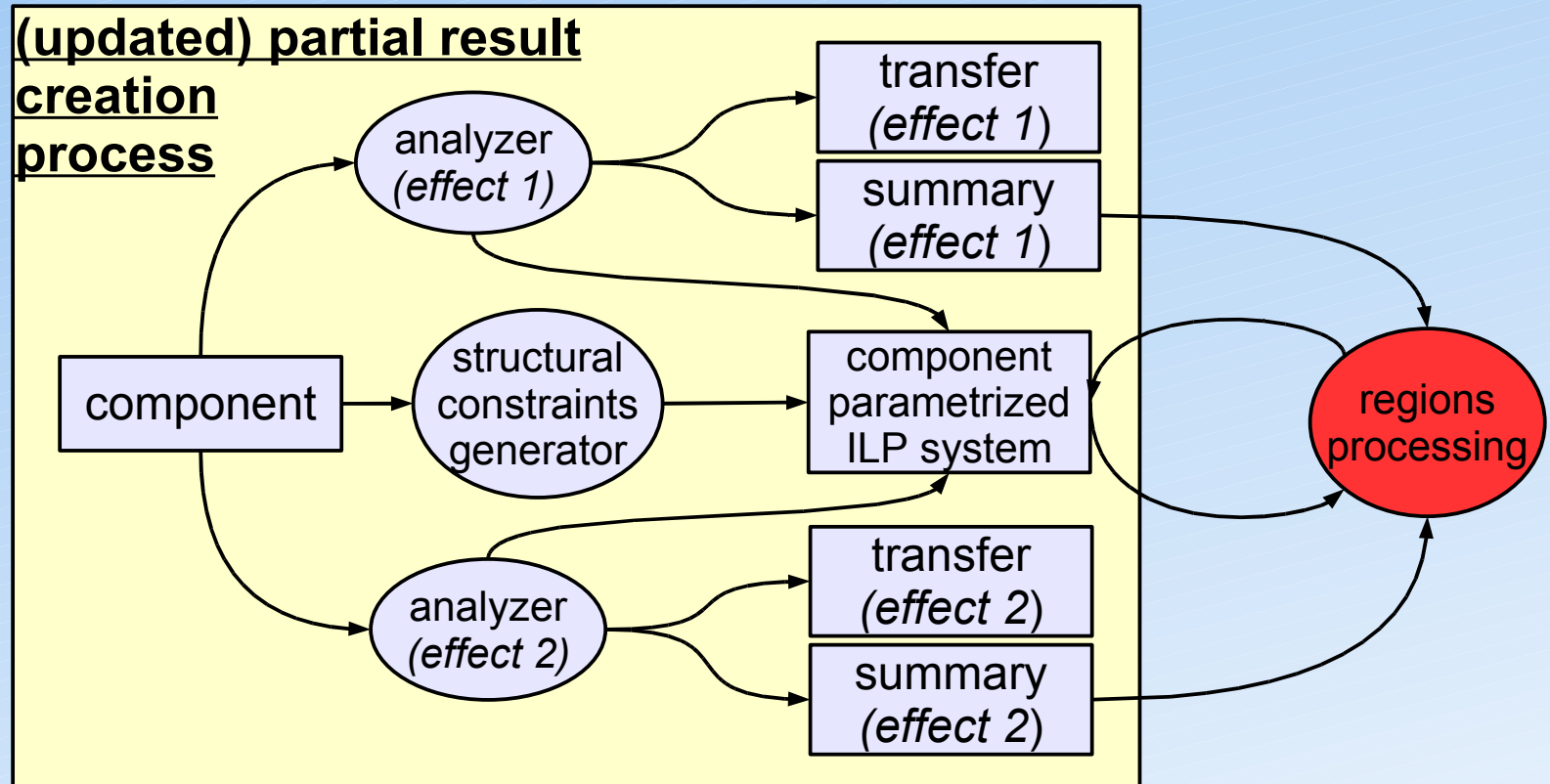
**cache transfer:**  
ageing (must)  
block 0  $\Rightarrow$  aged 2 times  
block 1  $\Rightarrow$  aged 2 times  
inserted blocks (must)  
block 2  $\Rightarrow$  inserted at age 0  
**loop bounds transfer:**  
{ $x = x + 10$ }

### summary

**cache summary:**  
CacheBlock 0, parameter  $p_0 =$   
if  $\text{must}(0) < 4 \Rightarrow$  Always Hit  
else Not Classified  
CacheBlock 1, parameter  $p_1 =$   
if  $\text{must}(1) < 4 \Rightarrow$  Always Hit  
else Not Classified  
CacheBlock 2, parameter  $p_2 =$   
if  $\text{must}(2) < 3 \Rightarrow$  Always Hit  
else Not Classified  
**loop bounds summary:**  
Loop 1, parameter  $p_3 = 3y$



# Precomputing independent regions



general idea:

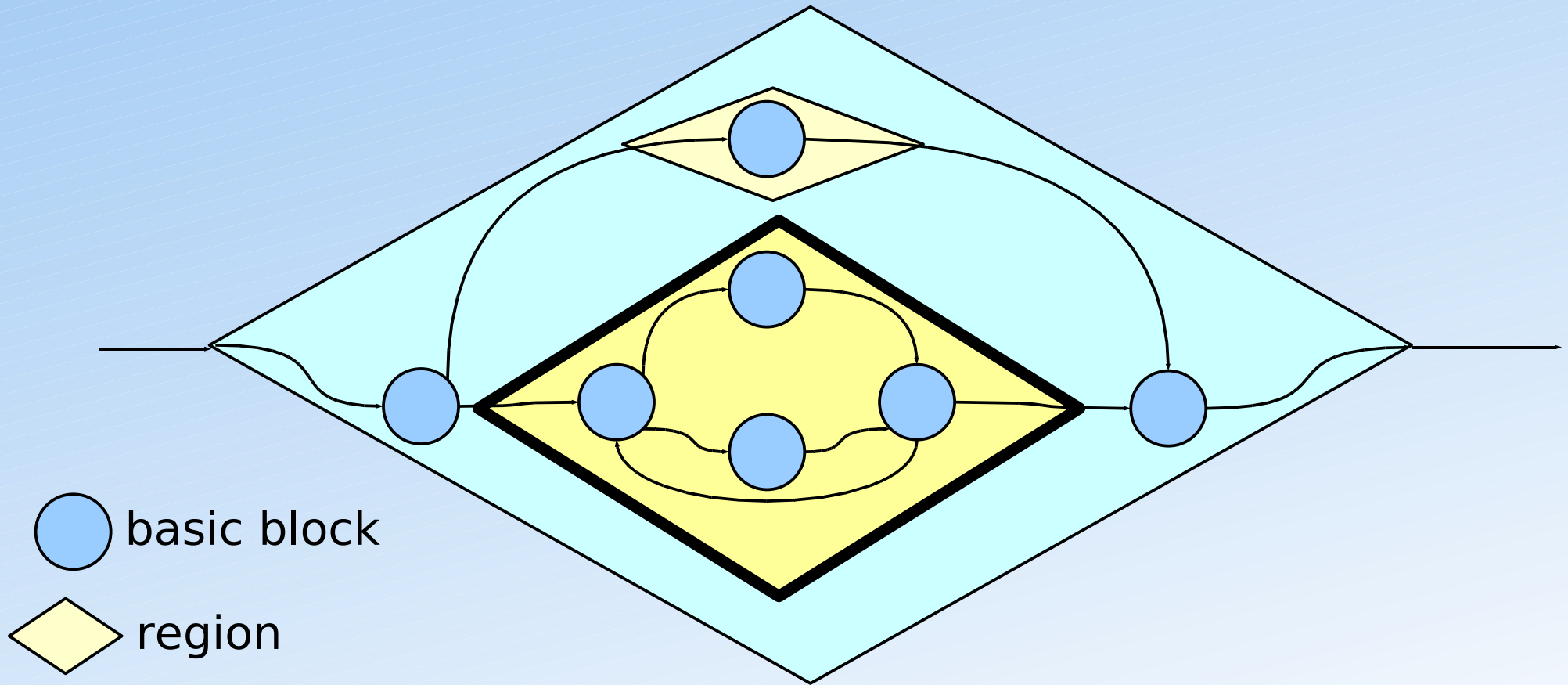
pre-compute the WCET of context-independent parts of the component to produce a smaller ILP system.

*(we need the summary functions to know which parts are context-independent)*

# *Precomputing independent regions*

Single-Entry Single-Exit regions

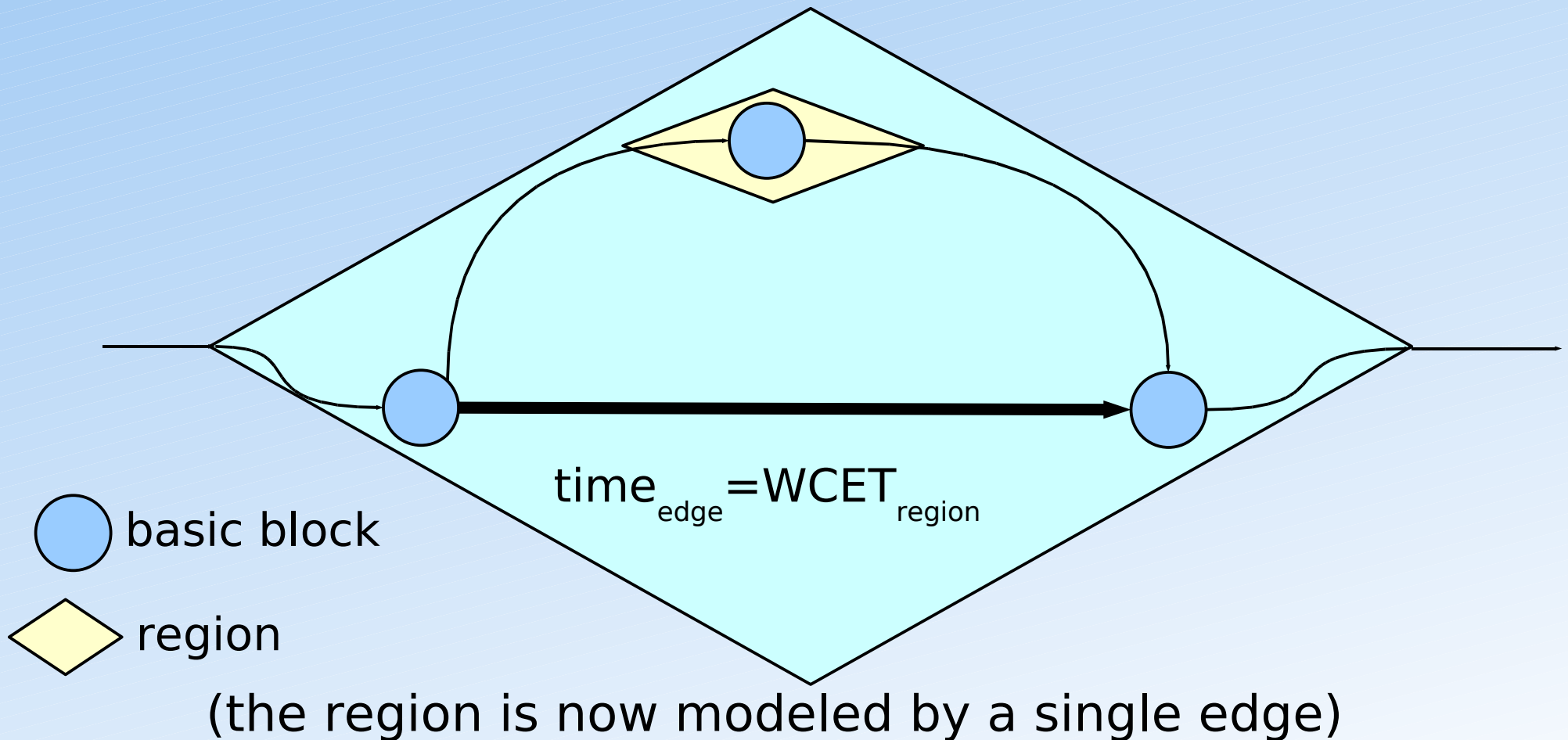
Constant-WCET regions can be pre-computed



# *Precomputing independent regions*

Single-Entry Single-Exit regions

Constant-WCET regions can be pre-computed





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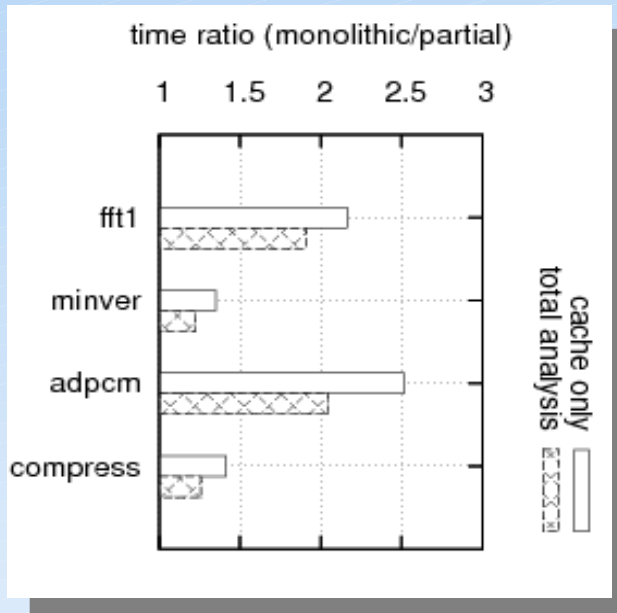
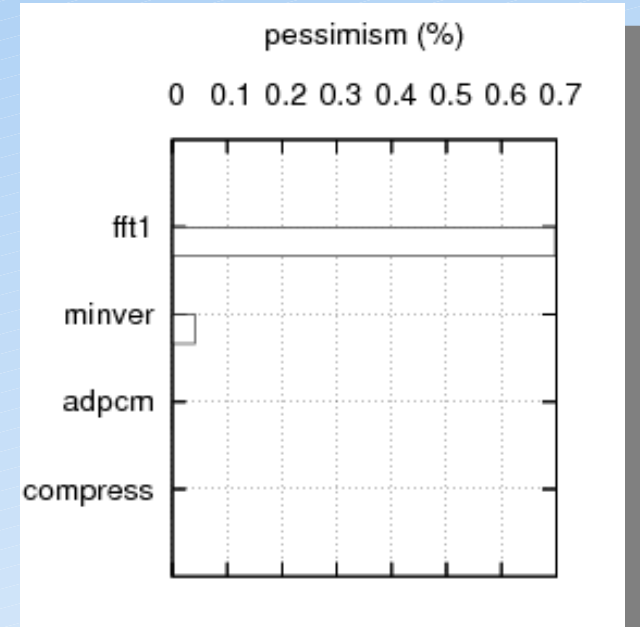
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***experimentation with OTAWA***

conclusion

# experimentation with OTAWA

pessimism increase measurement  
average WCET increase: 0.18%



analysis time measurement  
average gain: 1.86 **times**

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# ***Conclusion and future works***

blackbox components analysis  
handle COTS correctly  
speed up WCET computation

future works

finishing branch predictor partial analysis  
automate summary/transfer function  
creation