

*Scheduling energy consumption with local
renewable micro-generation and dynamic
electricity prices*

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GREEMBED'10
April 12, 2010



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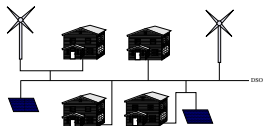
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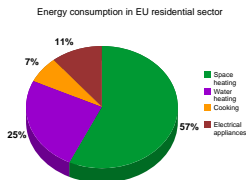
▶ Smart grid

- ▶ Transformation of electricity generation: distributed generation
- ▶ Transformation of electricity trading: real-time pricing, short-term contracting



▶ Smart home

- ▶ Transformation of electricity consumption: peak demand response, balancing power, load adjustment



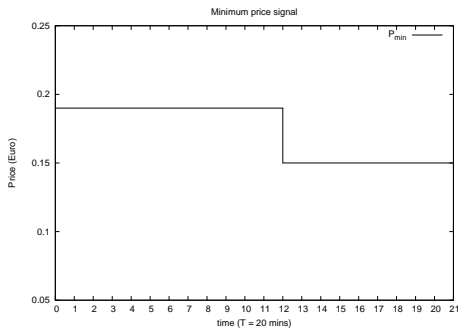


System model

Price signals

- ▶ a set P of price signals
- ▶ assumed to be predicted or provided for some time

$$p_{min}(t) = \min\{p_i(t)\}$$

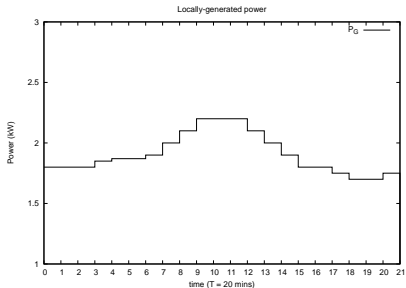




Locally-generated power

- ▶ a set G of local power micro-generators such as photovoltaics and wind mills
- ▶ $P_{G_i}(t)$ depends on weather, location
- ▶ assumed to be predicted
- ▶ assumed to be costless

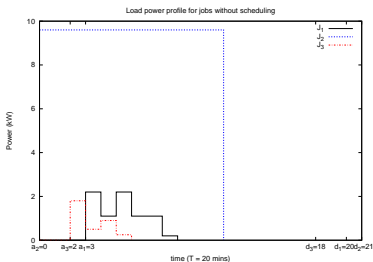
$$P_G(t) = \sum P_{G_i}(t)$$





Flexible tasks

- ▶ a set J of flexible tasks
- ▶ $J_i = (a_i, d_i, pr_i, L_i)$
 - ▶ a_i : earliest start time
 - ▶ d_i : deadline
 - ▶ pr_i : preemptability
 - ▶ L_i : load power profile





Problem statement

The scheduling problem

Given a task set J , a price signal set P , a locally-generated power P_G , and maximum allowed consumable power at any instant as P_{max} ; determine a schedule of the tasks such that the total cost for their execution is minimized.



Discretization of the problem

We assume piecewise constant functions with interval T

- ▶ $p_{min}(t), P_G(t)$ in interval $[\min(a_i), \max(d_i)]$

becomes

- ▶ $p_{min}[n], P_G[n]$ of length

$$N = \frac{(\max(d_i) - \min(a_i))}{T}$$

$L_i(t)$ becomes $L_i[n]$ with length

$$N_{L_i} = \frac{\text{length}(L_i)}{T}$$



Cost function

- ▶ Schedule for task J_i , $s_i[n]$: sequence of 0s and 1s
- ▶ Power consumed by J_i

$$P_i[j] = \begin{cases} L_i[\sum_{k=1}^j s_i[k]] & \text{if } s_i[j] = 1 \\ 0 & \text{otherwise} \end{cases}$$

- ▶ Total power consumed by all tasks

$$P_{tot} = \sum_i P_i$$

- ▶ Power to be billed (negative values are zeroed in P_{billed})

$$P_{billed} = P_{tot} - P_G$$

- ▶ Cost of the energy

$$C = P_{billed} \cdot p_{min} \cdot T$$



Constraints

- ▶ Tasks are scheduled to start after their earliest starting time:

$$\forall J_i : \quad a_i \leq T \cdot \min\{k : s_i[k] = 1\}$$

- ▶ Tasks are scheduled to finish before their deadlines:

$$\forall J_i : \quad d_i - T \geq T \cdot \max\{k : s_i[k] = 1\}$$

- ▶ Task J_i is scheduled as many times as the length of its load power profile:

$$\sum_{k=1}^N s_i[k] = N_{L_i}$$

- ▶ If task J_i is not preemptable, then it should be scheduled to run all at once:

$$pr_i = 0 \Rightarrow s_i(l) = 1 \\ \text{for } l \in [\min\{k : s_i(k) = 1\}, \max\{k : s_i(k) = 1\}]$$

- ▶ At no time, the total power withdrawn by all tasks exceeds the allowed maximum, P_{max} .

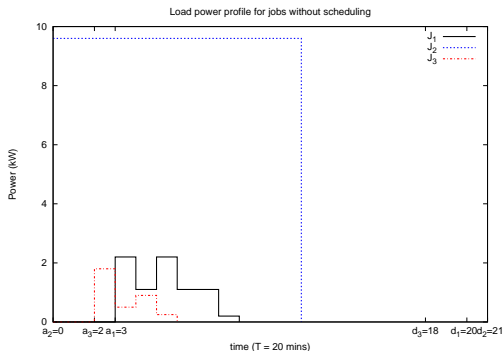
$$P_{tot}[k] \leq P_{max} \quad \text{for all } k$$



Case study

| Task (J_i) | earliest start | deadline | total duration | preemptable |
|-----------------|----------------|----------|----------------|-------------|
| Clothes washing | 1:00 | 6:40 | 2h | no |
| Car recharge | 0:00 | 7:00 | 4h | yes |
| Dish washing | 0:40 | 6:00 | 1h20' | no |

$$P_{max} = 15kW$$



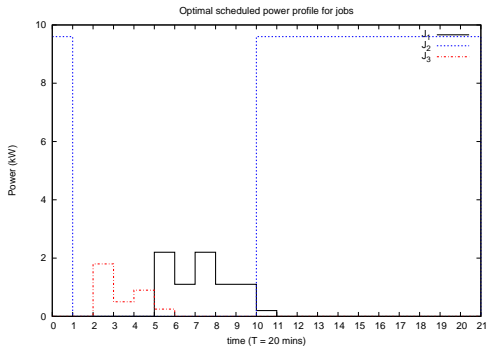


Case study

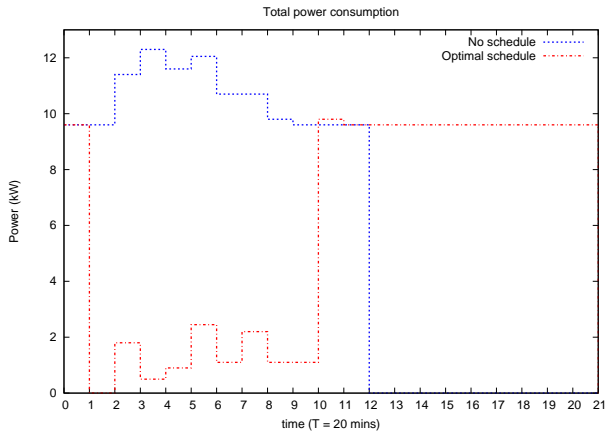
- ▶ huge search space

$$\binom{21}{12} \cdot 12 \cdot 11 = 38,798,760 \text{ valid schedules}$$

- ▶ took 35 minutes on a 1.8 GHz Intel Pentium Dual Core computer with 2GB of RAM
- ▶ 23% cost reduction from €6.5 to €5.0



- ▶ better use of P_G





Scheduling algorithm

- ▶ worst case complexity is $O(2^{MN})$
 M : number of tasks, N : number of time slots
- ▶ the scheduler should run in a reasonable time when new tasks arrive or predictions change
- ▶ need for fast admittance tests

$$P_{max} \cdot N \geq \sum_i \sum_j L_i[j]$$

ICT requirements

- ▶ A controller device
 - ▶ reads price signals,
 - ▶ makes contracts for short-terms with different DSOs.
 - ▶ communicates with home appliances (start, pause, resume and task information)
 - ▶ runs the scheduling algorithm
 - ▶ can be integrated with a smart metering device
 - ▶ may communicate with other controllers nearby
- ▶ need for standards for interoperable devices and seamless integration



Use cases

- ▶ Home level
- ▶ Community level
 - ▶ better trading power;
 - ▶ less communication/computation requirements on the infrastructure;
 - ▶ less cost of the ICT infrastructure per home due to sharing;
 - ▶ more predictable consumption at the community level;
 - ▶ ability to impose peak demand response and balancing power policies at the community level.
 - ▶ privacy concerns due to making household tasks transparent to a shared controller;
 - ▶ a community-level scheduling might provide less optimal results than home-level scheduling from the stand point of single users.

Management of locally-generated energy

- ▶ store, sell or waste



Conclusion

- ▶ A scheduling problem has been proposed to save money for household tasks based on the current trends in electricity markets, smart grids and smart homes.
- ▶ Finding the optimal schedule through exhaustive search is not feasible.
- ▶ We need efficient heuristics that would work for large number of tasks and time slots; and run on embedded systems.

Future work

- ▶ develop heuristics, assess their performance
- ▶ evaluate optimization performance in presence of prediction errors in p_{min} and P_G
- ▶ investigate negotiation in buying and selling of the energy
- ▶ investigate scheduling policies for demand side management



Thank you!

Questions?