

Logical Time at Work:

Capturing Data Dependencies and Platform Constraints

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Semantics as constrained logical time

(Modeling) languages are defined by:

- syntax
- (behavioral) semantics

Semantics should be explicit

Clock Constraint Specification Language (CCSL)

- Define the semantics of syntactical models
- Makes the syntactical models executable



Outline

Semantics of data-flow formalisms: Synchronous Data-Flow (SDF), Multidimensional-SDF

- The description of all the possible system execution schedules
- Functional description: internal constraints = data dependencies
- Time refinement enforce the constraints
- Environment/execution platform constraints



Outline

Semantics of data-flow formalisms: Synchronous Data-Flow (SDF), Multidimensional-SDF

- The description of all the possible system execution schedules
- Functional description: internal constraints = data dependencies
- Time refinement enforce the constraints
- Environment/execution platform constraints
- Define explicit semantics for SDF models
- Translate data-dependencies to execution dependencies
- Multi-dimensional semantics (MDSDF)
- External constraints



Clock Constraint Specification Language

Modeling and Analysis of Real-Time and Embedded systems (MARTE)

- Companion of the Time Package
- Chronological relations between events
- Clocks = possibly infinite and possibly dense totally ordered sets of instants

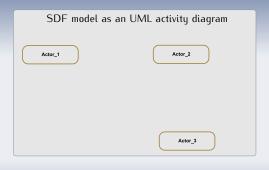
ccsl relations:

- precedence (≺)
- coincidence (≡)
- exclusion (#)

CCSL clock expressions:

- filteredBy (▼) by a binary periodic word
- delay (\$) by an integer value

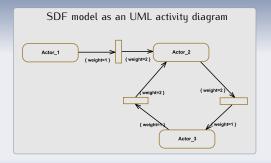




Syntax:

- computational elements
 - actors

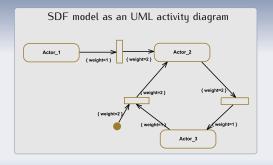




Syntax:

- computational elementsactors
- FIFO channels arcs

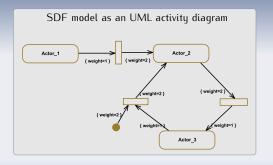




Syntax:

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Syntax:

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Synchronous data-flow semantics:

- fixed amount of data elements produces/consumed at each firing
- o local producer/consumer rules defined by each arc



General representation of a SDF arc

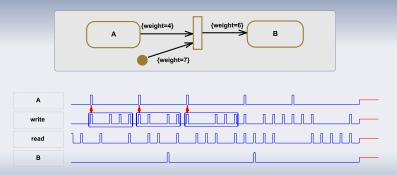


Relevant events in the system:

- Actor firings
- Element-wise write and read events on arcs



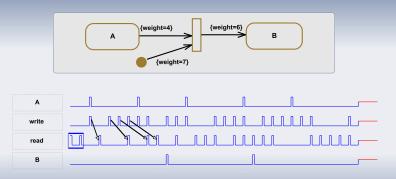
General representation of a SDF arc



 Each actor firing is followed by write_{weight} write events on each outgoing arc



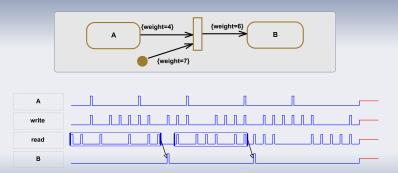
General representation of a SDF arc



 On each arc, read events (delayed by the initial value) must follow write events



General representation of a SDF arc



 read_{weight} read events on each of its ingoing arc precede an actor firing



General representation of a SDF arc



② write (read \$ delay)



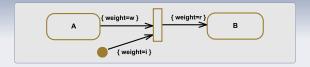
 \bigcirc (read \blacktriangledown (0^{weight-1}.1)^ω) \frown actor





SDF: execution dependency semantics

General representation of a SDF arc



Direct precedence (computed) between the two clocks



SDF: execution dependency semantics

Example of SDF arc



Direct precedence (computed) between the two clocks



 $(Clock_{master} \lor (011)^{\omega}) \lor (Clock_{slave} \$ 1)$



Direct precedence computation algorithm

Iterative algorithm to compute the parameters of the general execution precedence relation:

$$(Clock_{master} \lor M) \lor ((Clock_{slave} \$ d) \lor S)$$



Direct precedence computation algorithm

Iterative algorithm to compute the parameters of the general execution precedence relation:

$$(Clock_{master} \lor M) \subset ((Clock_{slave} \$ d) \lor S)$$

$$d = \lfloor delay_{weight} / read_{weight} \rfloor$$
 $initial = delay_{weight} \mod read_{weight}$



Direct precedence computation algorithm

Iterative algorithm to compute the parameters of the general execution precedence relation:

$$(Clock_{master} \lor M) \lor ((Clock_{slave} \$ d) \lor S)$$

$$d = \lfloor delay_{
m weight}/read_{
m weight}
floor = \lfloor 7/6
floor = 1$$
 $initial = delay_{
m weight} \mod read_{
m weight} = 7 \mod 6 = 1$

	initial		+4	+4	-6+4	-6+4
tokens	1		5	9	7	5
			< 6	> 6	> 6	done
binary		(0	1	1)



Local rules – global functionality

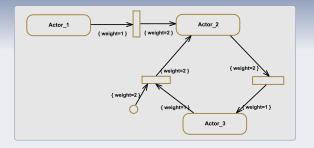


$$(Actor_A \lor (011)^{\omega}) \lt (Actor_B \$ 1)$$





Local rules – global functionality







Encoding MDSDF in CCSL

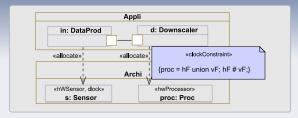


- straightforward multi-D extension of 1-D SDF
- quasi-independent relations producer/consumer by dimension

$$in_1 \leftarrow (hF_1 \lor (1.0^2)^\omega)$$
 $hF_1 \leftarrow (vF_1 \lor (1.0^2)^\omega)$
 $in_2 \leftarrow hF_2$ $(hF_2 \lor (0^8.1)^\omega) \leftarrow vF_2$



External constraints



multidimensional order – environment constraints

$$in_1 = (s \lor 1. (0)^{\omega})$$
 $in_2 = s$
 $hF_1 = (hF \lor 1^3. (0)^{\omega})$ $hF_2 = (hF \lor (0^2.1)^{\omega})$
 $vF_1 = (vF \lor 1^9. (0)^{\omega})$ $vF_2 = (vF \lor (1.0^8)^{\omega})$

execution platform constraints

$$proc = hF + vF$$
 $hF \# vF$



System refinement

Formal specification encoding the entire set of schedules corresponding to a correct execution

Enforce the constraints:

- internal (data dependencies)
- external (environment or execution platform)
- buffer capacities reverse constraints
- physical time durations: execution, communication, ...



Repetitive Structure Modeling

Application – data dependencies Appli-rsm ar: AR br: BR cr: CR «Shaped» «Shaped» «Shaped» shape = {6.2.*} shape = {3,*} shape = {2,*} a: A c: C «Shaped» «Shaped» dr: DR «Shaped» «Shaped» shape = {4.* shape = {2,2,3,* shape = {6.6.*} shape = {4.6.*} «Shaped» shape = {2,3,*} «Shaped» shape = {6,*}



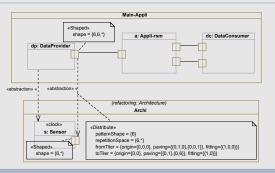
Repetitive Structure Modeling

Architecture — execution platform | Architecture — execution platform | Architecture | Architec



Repetitive Structure Modeling

Multi-D structures - projection in time and space





Internal constraints

Execution dependencies: precedence (ccsl from rsm::Order Behavior::Order) Order {{finishAction Sensor ((finishBehavior Aby 2 strictly ((finishBehavior B strictly precedes startBehavior C ;}} by 3 strictly precedes precedes startBehavior B :)) startBehavior A;}} «timedProce. «timedProce.. «timedProce... «timedProces Call A_flatten Call B_Flatten Action Sensor Call C Flatten ((startBehavior A {{startBehavior B ((startBehavior C nonStrictly precedes nonStrictly precedes nonStrictly precedes finishBehavior A:33 finishBehavior B;}} finishBehavior C:33 ((finishBehavior A «timedProce... strictly precedes Call D Flatten startBehavior D:33 {{startBehavior_D nonStrictly precedes finishBehavior_D;}}

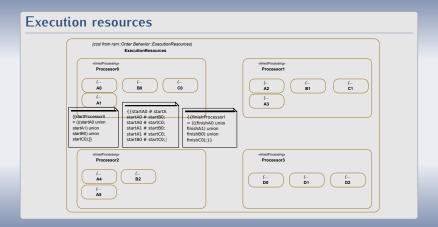


Internal constraints

Task parallelism – unbounded resources (ccsl from rsm::Order Behavior::Flatten) Flatten B_Flatten C_Flatten A_flatten {{startA0 nonStrictly precedes finishA0:33 CO A0 A1 A2 D_Flatten А3 {{startBehavior A {{(sup(finishA0, nonStrictly finishA1, finishA2, precedes (inf finishA3, finishA4, (startA0, startA1, finishA5)) = A5 finishBehavior A:1 startA4, startA5));33 D2



External constraints





Time Model & CCSL (MARTE)

Logical Time:

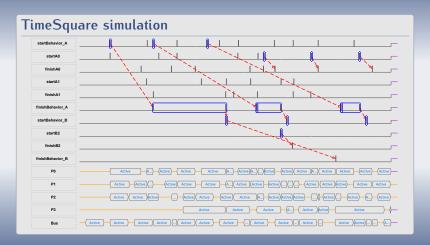
- Dependencies (data, execution)
- Parallelism
- Concurrence (resources: execution, communication, storage)

Projection into Physical Time

- Physical clocks
- Durations
- Delays
- o ...



Schedule constraints





Conclusion

Results:

- Define explicit semantics
- For analysis, static scheduling, state space exploration
- Simulation in TimeSquare¹

More complex languages:

- Array-OL: no predefined order
- Uniform loop nests: complex dependence analysis

Heterogeneous systems

http://www-sop.inria.fr/aoste/dev/time_square