The AADL Behavioral annex : an example of use

Frank Singhoff, LISyC/UBO UML-AADL workshop'2010

Verification with the

real-time scheduling theory (1/5)

□ Real-time systems:

- 1. Functions of real-time systems may have timing constraints.
- 2. Example of timing constraints : deadlines
- 3. Functions must meet timing constraints: how to check them at design time ?
- Timing analysis with real-time scheduling theory (also called "Rate Monotonic Analysis"):
 - 1. Modeling functions: simplified models of task = processor demand + deadline (e.g. periodic task model).
 - 2. Use of well known task schedulers (e.g. Rate Monotonic).
 - 3. Verification: either by simulation or by feasibility tests.

Verification with the real-time scheduling theory (2/5)

Example of feasibility test to perform verification : worst case response time of a set of periodic tasks, sharing resources, scheduled by Rate Monotonic, (Joseph & Pandia 1986):

$$R_{i} = C_{i} + B_{i} + \sum_{j \in hp(i)} \left[\frac{R_{i}}{P_{j}} \right] \cdot C_{j} \leq deadline_{i}$$

Verification method:

- □ Compute Ri, the worst case response time of task i as a sum of:
 - 1. Ci, worst case execution time of task i (or task capacity).
 - 2. Bi, worst case waiting time of task i for shared resources.
 - 3. Waiting time for the processor due to scheduling of the other tasks.
- Compare Ri with the task i deadline.
- Weakness of this approach : low accuracy of Ci and Bi may lead to compute a too pessimistic Ri.

Verification with the real-time scheduling theory (3/5)

```
PROCESS IMPLEMENTATION a_process.I
SUBCOMPONENTS
thread1 : THREAD a_thread.I;
thread2 : THREAD a_thread.I;
resource1 : DATA a_data.I {
Concurrency_Control_Protocol=>PCP};
END a_process.I;
```

```
THREAD IMPLEMENTATION a_thread.I

PROPERTIES

Dispatch_Protocol => Periodic;

Deadline => 10 ms;

Period => 10 ms;

Compute_Execution_Time=> 1 ms ..7 ms;

Bound_On_Data_Blocking_Time => 7 ms;

END a_thread.I;
```

□AADL Model without Behavioral annex specification:

□Verification is possible.

But : no thread behavior specification => we assume that threads use the shared resource during all their execution time, which leads to a pessimistic Bi.

□May imply a pessimistic worst case response time (Ri)

Behavioral annex provides a mean to increase accuracy of the model.

Verification with the real-time scheduling theory (4/5)

```
THREAD IMPLEMENTATION a thread.I
                                        PROPERTIES
PROCESS IMPLEMENTATION a process.I
                                         Deadline => 10 \text{ ms};
SUBCOMPONENTS
                                         Compute_Execution_Time => 1 ms .. 7 ms;
thread1 : THREAD a thread;
                                         Bound On Data Blocking Time => 7 ms;
thread2 : THREAD a thread;
                                         ANNEX Behavior Specification {**
resource1 : DATA a data {
                                          states s : initial complete final state;
Concurrency_Control_Protocol=>PCP};
                                          transitions t : s -[on dispatch]-> s
END a process.I;
                                           {computation(4 ms); resource1!<;</pre>
                                            computation(3 ms); resource1!> };
                                        **};
                                        END a thread. I;
```

□AADL model with Behavioral annex specification:

- □Specification of the thread behavior: specify when threads use the shared resource (allocate and release time of the resource).
- □Less pessimistic Bi approximation (Bi=3 and not 7), and then more accurate worst case response time (Ri).

Bi

Verification with the real-time scheduling theory (5/5)



- AADL and its annexes are standardized, and make possible tool interoperability.
- Example of an AADL Tool chain (Ellidiss & Univ. Of Brest):
 - 1. STOOD or ADELE AADL editor (both core AADL & annex specifications)
 - 2. Static analysis of Behavioral annex specifications. Produce enriched Cheddar XML with shared resource accesses (Bi).
 - 3. Analysis with feasibility tests by Cheddar.

Conclusion of the panel

□ AADL allows architecture designers to:

- □ Model, verify and generate embedded real-time systems architectures.
- □ Assume an iterative development process, and allows model refinements.

□ AADL core language and annexes:

- □ Core language defines system architectures.
- Annexes give different view points of system architectures (e.g. verification, code generation, ...).
- Examples of annex: ARINC 653, behavioral, data modeling, error, ...
- □ AADL has been used as a pivot language in numerous research projects: IST-ASSERT, TOPCASED, SPICES, ANR Flex-eWare, MOSIC, SAVI AVSI, ...
- Many AADL interoperable tools: OSATE, Furness, TINA, CPN/AMI, Lotos, BIP, GreatSPN, STOOD, ADELE, Ocarina, Cheddar, …
- □ Further readings: <u>http://www.aadl.info</u>