The AADL Behavioural annex¹

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Introduction

Architecture languages \rightsquigarrow verification of dynamic aspects:

- composition, temporal properties,
- quantitative analysis: sizing, performance, energy,
- reliability.

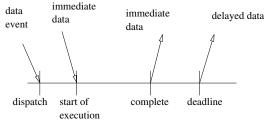
The behavioural annex

Why a behavioural annex for AADL?

- AADL relies on the analysis of source text.
- AADL calls and flows do not depend on data.
- AADL behaviour is basically described by:
 - the AADL execution model
 - quantitative aspects like wcet.

The behavioural annex allows a refinement of these aspects.

The AADL data port protocol



The AADL behavioural annex:

- describes how the dispatch is triggered.
- gives access to the received and sent data
- describes what happens when control is gained
- says when control is relinquished

The behavioural annex

- Describe the internal behavior of component implementations as a state transition system with guards and actions.
- Extend the default run-time execution semantics that is specified by the core of the standard, such as thread dispatch protocols.
- Provide behavioral refinement for mode transitions.
- Introduce subprogram calls synchronization protocols.

These extensions are introduced through properties and annexes.

```
Introduction
The behavioural annex
Conclusion
```

```
behavior_annex ::=
  [ variables behavior_variable + ]
  [ states behavior_state + ]
  [ transitions behavior_transition + ]
behavior_state_kind ::=
  [ initial ] [ complete ] [ final ]
execution_behavior_transition ::=
  [ behavior_transition_label : ]
  source_state_identifier , source_state_identifier *
  -[ [ behavior condition ] ]->
  destination state identifier [ behavior actions ] :
```

thread merger
features
p1 : in event data port Basic_types::integer;
p2 : in event data port Basic_types::integer;
m : out event data port Basic_types::integer;
end merger;

```
thread implementation merger.twopersistentstates
annex behavior_specification {**
variables
    x1 : data Basic_types::integer;
    x2 : data Basic_types::integer;
    states
    s0 : initial complete state;
```

comp : state;

next1, next2 : complete state;

transitions

s0 -[on dispatch p1]-> next2 { x1 := p1 }; s0 -[on dispatch p2]-> next1 { x2 := p2 }; next1 -[on dispatch p1]-> comp { x1 := p1 }; next2 -[on dispatch p2]-> comp { x2 := p2 }; comp -[x1 < x2]-> next1 { m!(x1) }; comp -[x2 <= x1]-> next2 { m!(x2) }; ***};

Synchronization protocols

- subprogram calls in AADL is synchronous.
- alternate synchronization protocols (HRT-HOOD):

Server_Call_Protocol: type enumeration (ASER,HSER,LSER)
applies to (provides subprogram access);

- ASER : the caller is never blocked.
- LSER : the caller waits for the acceptation of the request.
- HSER : the caller waits for the completion of the request and gets results if any.

Conclusion

- The behavioural annex is now part of the AADL standard.
- Supported by the OSATE-TOPCASED project.
- Verification support through the FIACRE language : pivot verification language of the TOPCASED project.
- Used for case studies in industrial research projects.