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# WCET analysis of a parallel 3D multigrid solver executed on the MERASA multi-core

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# Motivation

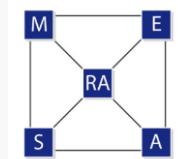
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- **Hardware trends**
  - common requirements
    - low energy consumption
    - high integration / low cost
    - high performance
  - architectural solution: **multicore processors**
- **New issues**
  - time-predictable multicore architectures
    - projects: **MERASA**, PREDATOR
  - parallel programming of embedded applications
  - **WCET analysis of parallel programs**
    - reporting experiments on a data parallel industrial application

# The MERASA project



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## □ Objectives:

- develop **multi-core processor designs** (from 2 to 16 cores) and **system-level software** support for mixed-critical workloads
- provide **WCET analysis tools**
  - static analysis: OTAWA
  - hybrid measurement-based: RapiTime
- evaluate with **pilot studies** submitted by industrial partners

## □ Partners:

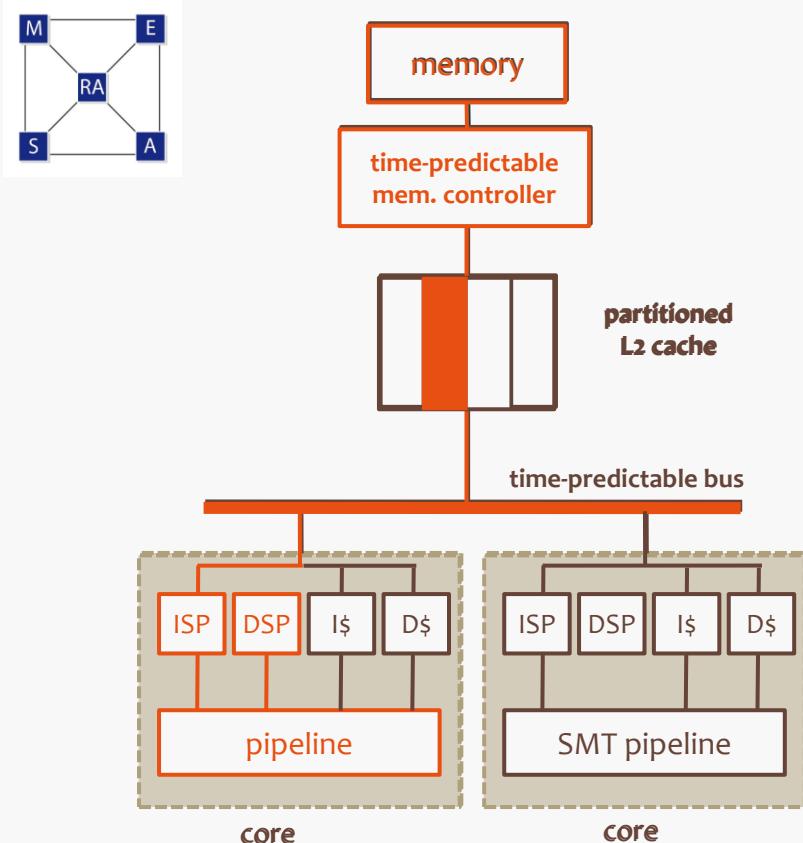
- **Participants:** Univ. Augsburg, Barcelona Supercomputing Center, Univ. Toulouse, Rapita Systems Ltd., Honeywell
- **Industrial Advisory Board:** Airbus, BAUER Maschinen GmbH, European Space Agency, Infineon , NXP

# The MERASA multicore and system software

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## □ Isolation of critical threads

- ▣ specific scheduling in pipeline
- ▣ time-predictable bus
- ▣ partitioned L2 cache
- ▣ time-predictable memory controller

## □ System software

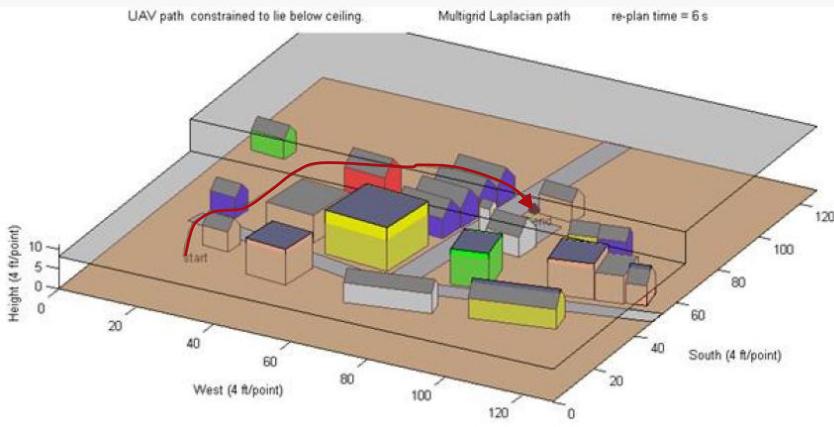
- ▣ time-predictable memory management
- ▣ POSIX-compliant time-predictable synchronizations
- ▣ two-level thread scheduling

# The 3D multigrid solver

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## □ Part of application for airbone collision avoidance (Honeywell)

- path planning using Laplace's equation
- numerical solution:
  - domain partitioning → voxels
  - set potential of each point to average of neighbors
  - several iterations
- multigrid technique:
  - several phases, varying grid size



## □ Version considered in this study

- three phases -- each phase = **interpolation** step + **iteration** step

```
for (x=0; x<NX; x++)  
for (y=0; y<NY; y++)  
for (z=0; z<NZ; z++)  
v[x][y][z] = compInterpolate(old_v);
```

```
for (i=0; i<NUM_ITE; i++)  
for (x=0; x<NX; x++)  
for (y=0; y<NY; y++)  
for (z=0; z<NZ; z++)  
v[x][y][z] = compIterate(v);
```

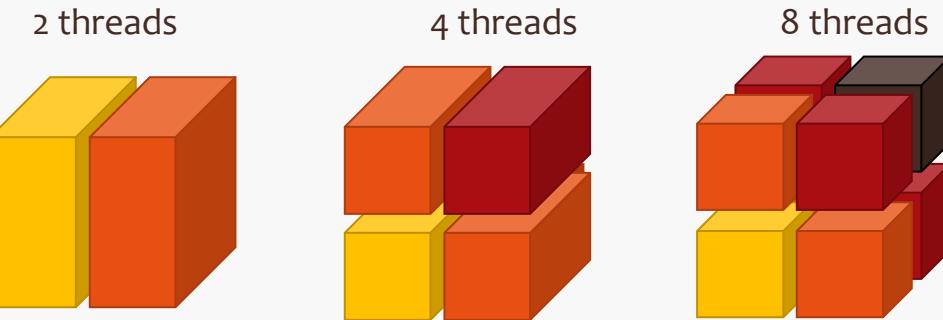
# The parallelized 3D multigrid solver

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## □ Domain splitting



## □ Orchestration

main thread

```
readImage();  
initiate P1S1(interpolation)  
  
wait for child threads  
initiate P1S2(iteration)  
  
wait for child threads  
... // same for P2 and P3
```

child thread #1

```
wait for main thread  
process P1S1(interpolation)  
  
wait for main thread  
process P1S2(iteration)  
  
... // same for P2 and P3
```

child thread #2

```
wait for main thread  
process P1S1(interpolation)  
  
wait for main thread  
process P1S2(iteration)  
  
... // same for P2 and P3
```

# Breakdown of the WCET

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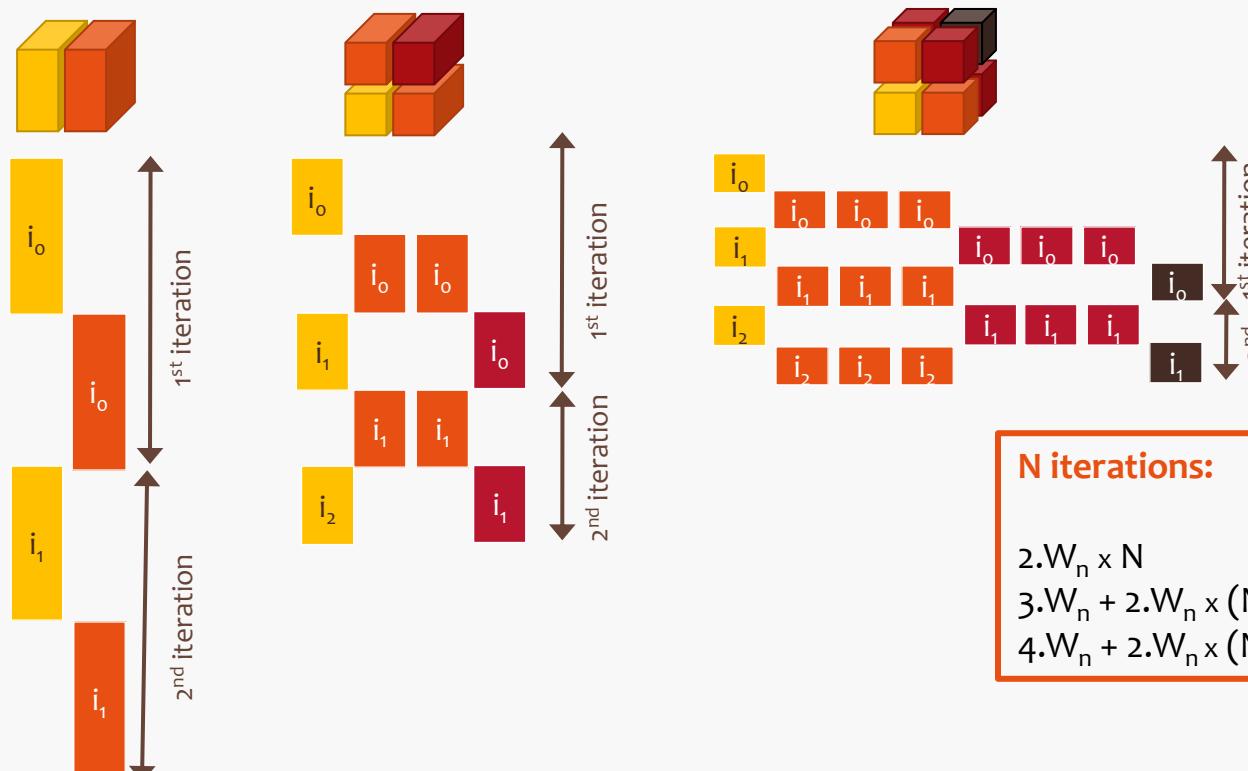


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## □ First level



## □ Second level



# Analysis of synchronizations

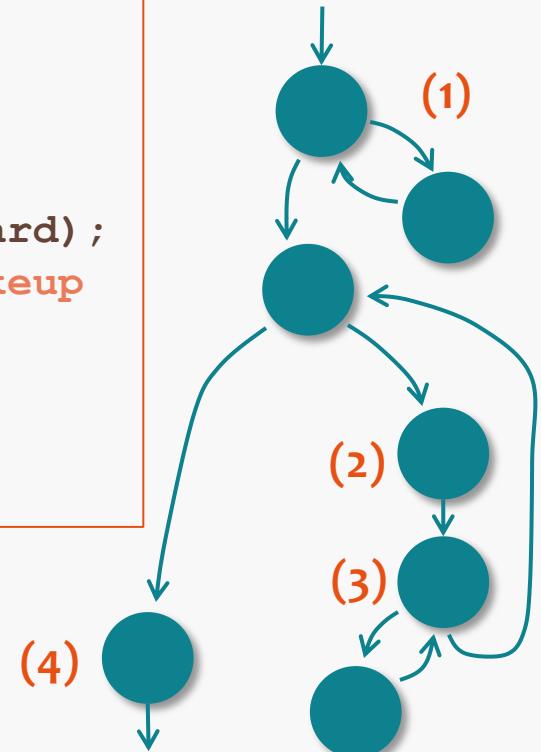
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## □ Example: mutex lock



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```
int mutex_lock(mutex_t mutex) {
(1)    sl_lock(&mutex->guard) ;
...
    while (trylock (&mutex->thelock) {
        ... // insert thread into queue
        sl_unlock_and_set_suspended (&mutex->guard) ;
(2)    spinlock_lock (&mutex->guard) ; // on wakeup
(3)
}
...
(4)    sl_unlock (&mutex->guard) ;
}
```



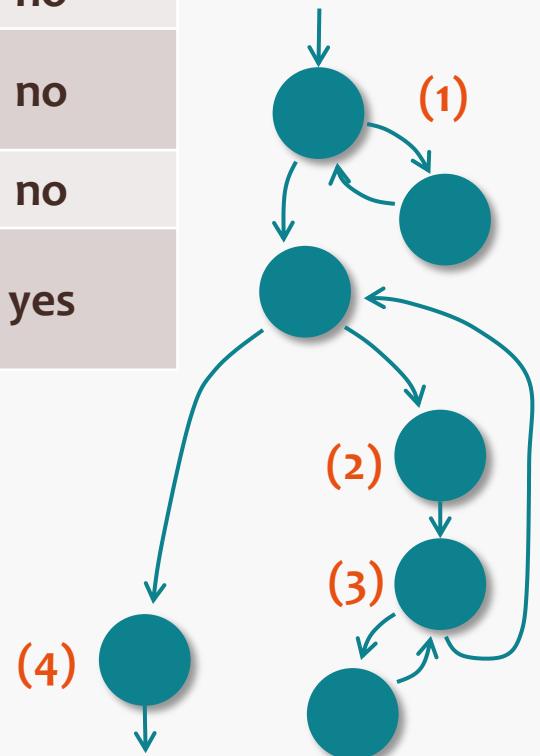
# Analysis of synchronizations (cont'd)

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Term	depends on ...	
	# threads	application
$T_e$	Execution time when the variable is free	no
$T_{w1}$	Overhead execution time when the thread has to wait for the variable	no
$T_{w2}$	Waiting time related to system-level variables	yes
$T_{w3}$	Waiting time related to application-level variables	yes



# Experimental results

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## □ Estimated WCETs (OTAWA – perfect ISP – no L2 cache)

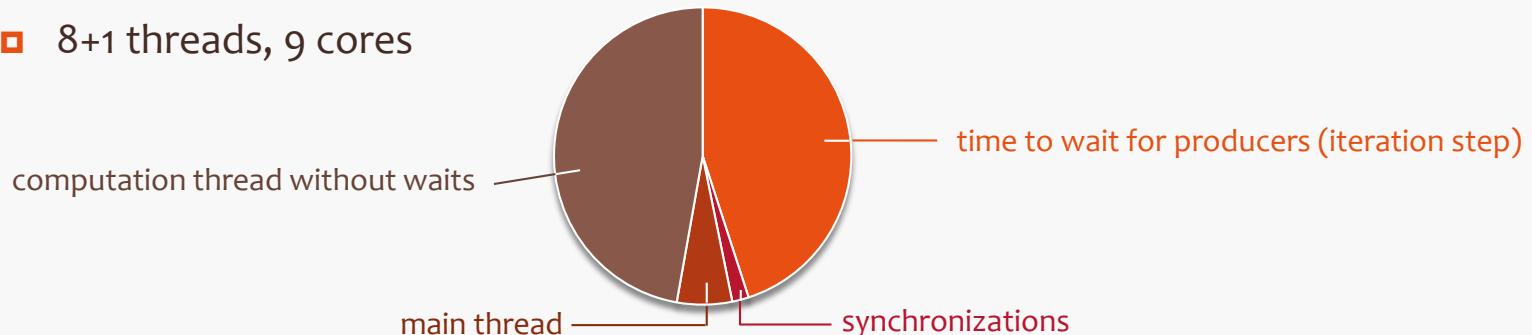
	# threads			
	1	2+1	4+1	8+1
1 core	53,990,765	-	-	-
3 cores	58,297,375	66,849,369	-	-
5 cores	62,603,985	73,372,109	40,389,428	-
9 cores	71,217,205	86,417,589	47,678,968	28,105,761

Annotations on the table:

- A teal circle with  $\times 1.3$  points to the value for 3 cores.
- An orange circle with  $\div 1.9$  points to the value for 5 cores.
- A teal circle with  $\div 2.5$  points to the value for 9 cores.

## □ WCET breakdown

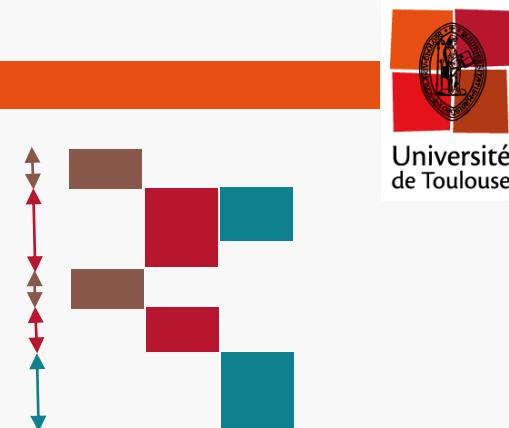
- 8+1 threads, 9 cores



# Lessons learnt and future work

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- WCET analysis of a parallel application requires:
  - understanding the structure of the application
    - dependency analysis
  - analyzing the synchronizations
    - where?
    - number of concurrent threads?
    - paths on which locks are hold
- Automatic WCET analysis of parallel code:
  - retrieve required information (structure + synchronizations)
    - source/object code analysis
    - user annotations
  - orchestrate the WCET computation
- Parallelisation guidelines
  - to facilitate WCET analysis



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Questions?