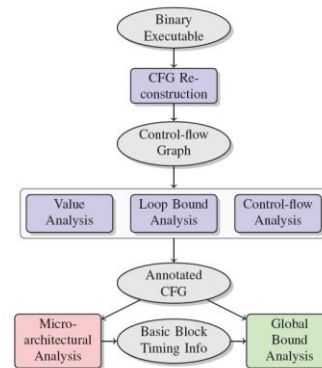


Integrating Symbolic Pipeline Analysis with Abstract Caches

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Static WCET Analysis

- **Computes:** Safe upper bounds on WCET of a task
 - safety critical applications, hard WCET
- **Requires:** Full coverage of
 - feasible program paths
 - inputs
 - hardware states (pipeline, caches)
- **Method:** AI + ILP [aiT WCET analyzer]
 - abstract interpretation
 - integer linear programming
 - several interacting analyses



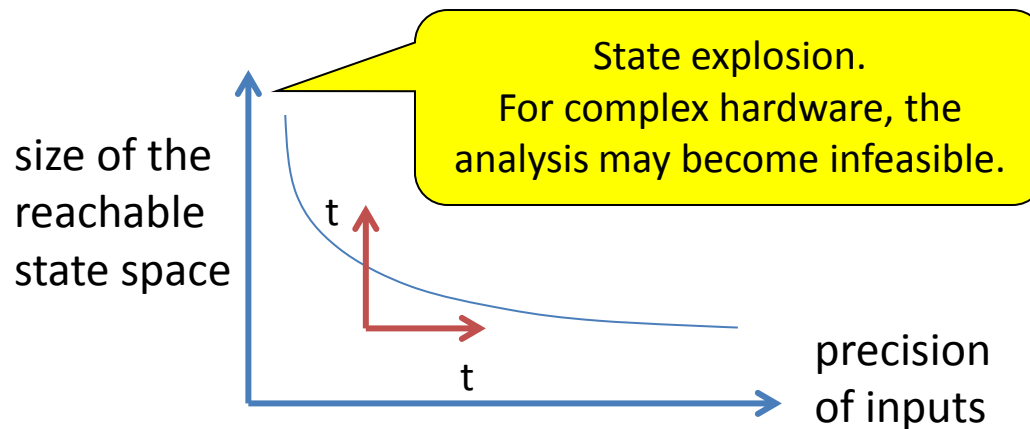
Micro-architectural Analysis

- **Pipeline analysis**
 - abstract model
 - pipelining, speculation, branch prediction
 - computes an over-approximates reaching states
 - counts cycles on the basic block level
- **Cache analysis**
 - computes an over-approximation of possible cache contents
 - predicts sure hit or miss, or don't know
- **Analysis inputs**
 - from control flow and value analysis
 - branch and call targets, possible register contents
 - all inputs may be *imprecise*



Imprecise Inputs

- **Intervals and sets** instead of precise values for
 - possible register contents
 - call and branch targets
- **Effect** on hardware-level analysis
 - must consider more possibilities



Efficient State Space Coverage

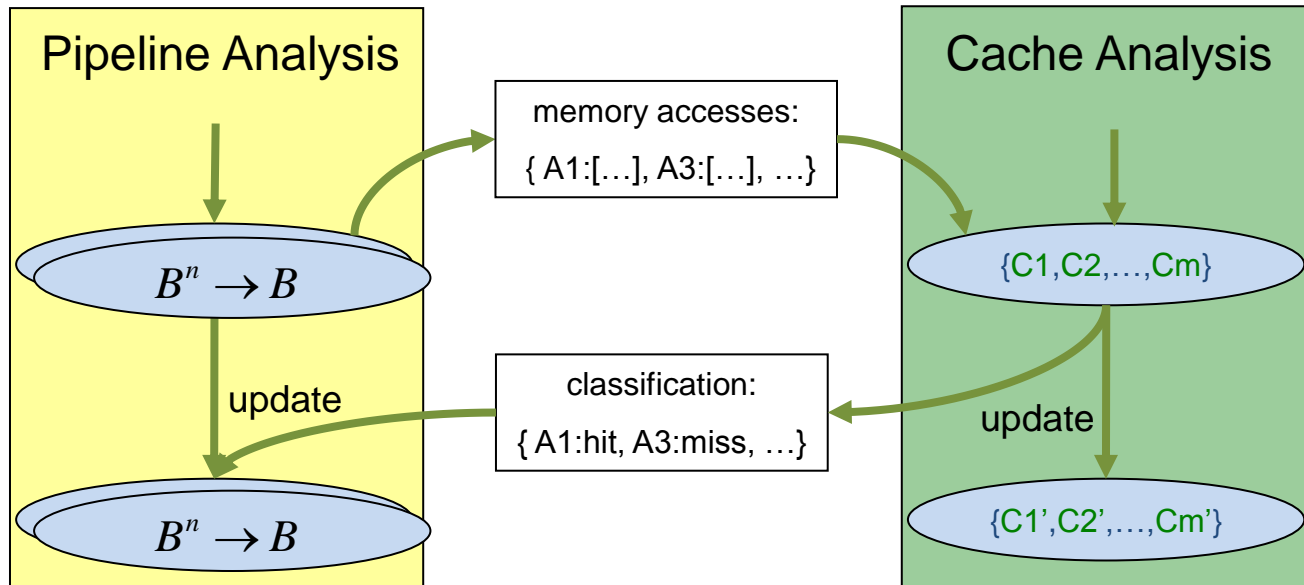
Cache analysis

- Abstract interpretation
- Efficient join operator
 - safe over-approximation
 - loses precision
- Efficient **abstract** representation

Pipeline analysis

- Enumerates all states
 - Timing anomalies
 - Domino effects
- Symbolic representation
 - Implicit, Boolean functions
 - Binary decision diagrams
 - Exploits redundancies
- Efficient **implicit** representation

Pipeline-Cache-Interaction

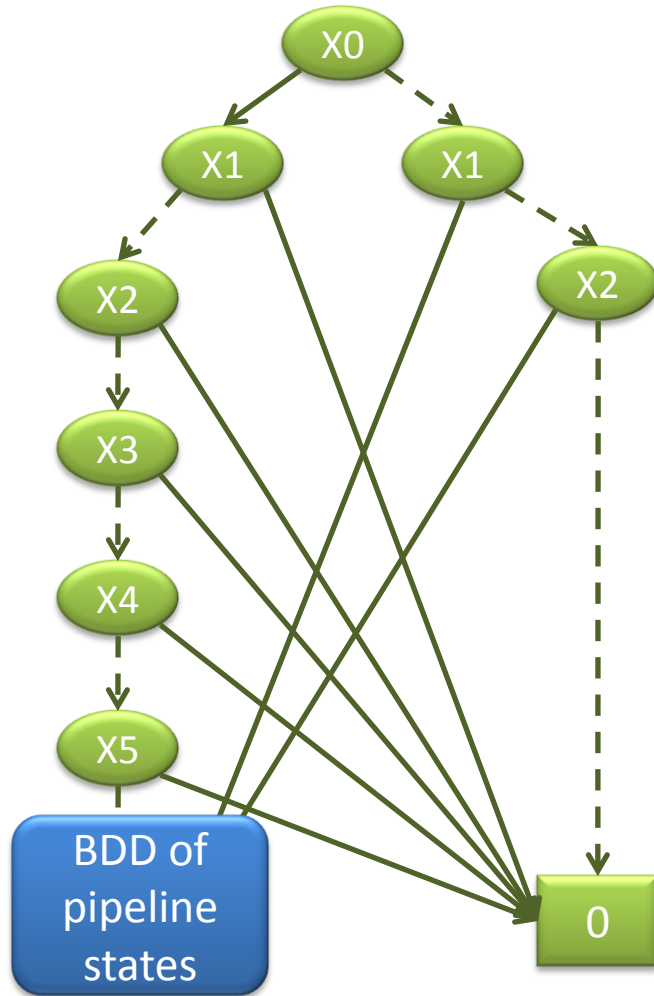


- Exchange information only between pipeline and cache states which cover the same execution history.
 - Requires a combination of pipeline and cache states.
- Explicit-state pipeline analysis uses tuples $(P_k, C_k), (P_{k'}, C_{k'}) \dots$
- What to do for symbolic pipeline analysis?

Proposed Solution

- Combine **many** pipeline states with **one** cache state
 - semi-symbolic domain
 - 1 domain element = set of tuples: $(B^n \rightarrow B, \hat{C})$
- **Maintain property**
 - all pipeline states in the same tuple access the same cached memory interval
- **Requirement**
 - memory is addressed via the first variables in the BDD representation of pipeline states

Memory Access Interval

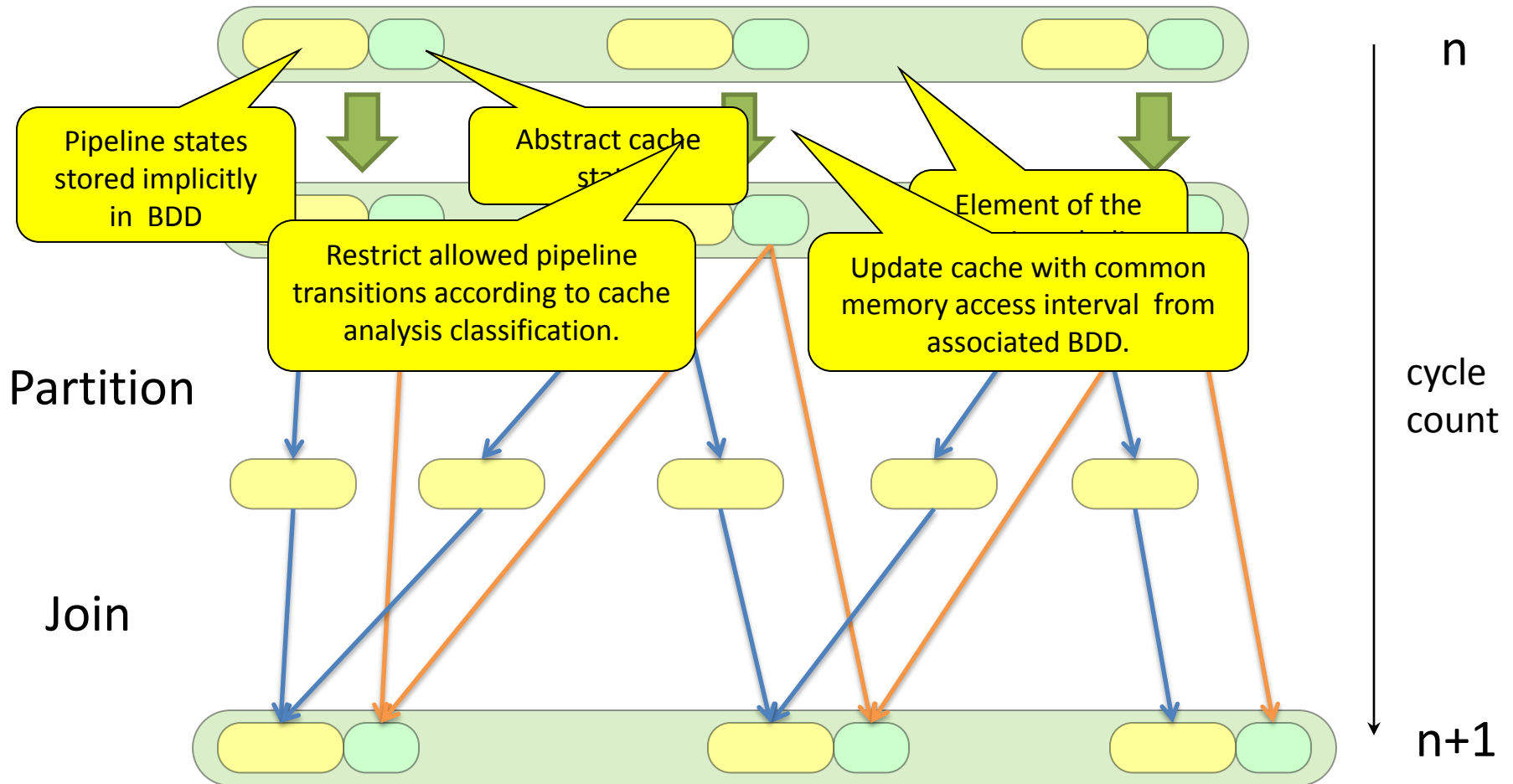


X0	X1	X2	X3	X4	X5
1	0	0	0	0	0
0	1	-	-	-	-
0	0	1	-	-	-

X0	X1	X2	X3	X4	X5	ub
1	0	0	0	0	0	32
0	1	1	1	1	1	31
0	0	1	1	1	1	15

X0	X1	X2	X3	X4	X5	lb
1	0	0	0	0	0	32
0	1	0	0	0	0	16
0	0	1	0	0	0	8

Update + Balancing



Experiments

- **Goal:** asses the efficiency of the proposed domain
 - efficient if combining **many** pipeline states with single cache state
- **Instrumented** the (explicit-state) **MPC755 pipeline** model to report for each access into cached memory:
 - accessed memory address or interval
 - type of access: instruction or data
 - cycle count since start of the current basic block
- **6 tasks** of a typical industrial software project
 - fully unrolled & annotated (avoid serious state explosion)
 - still reaches up to a few thousand states per cycle

Results

Explicit-state experiment

- cycle count since start of the current basic block
- # states with same cycle count and memory access
- # different memory accesses with same cycle count

Semi-symbolic domain

- state space exploration layer
- # pstates per BDD
 - Average : 27 / 20
 - Maximum : 8544 / 8115
- # tuples per domain element
 - Average : 2 / 1
 - Maximum : 42 / 6

Conclusion

- Covering the state space of complex hardware can become a problem in static WCET analysis.
- **Solutions** for caches and pipelines **already exist**.
 - abstract interpretation of caches, symbolic pipeline analysis
- Proposed an approach for integration.
 - **Efficient if combining many pipeline states with a single abstract cache state.**
 - Experimental data indicates that favorable combinations can be expected.
 - May lose precision due to additional joins of cache states.

Q & A