Mapping Software-Defined Radio Applications onto MPSoCs

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Where innovation starts

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Radio Baseband Processing

- > The three main stages are filters, modem and codec.
- These stages are customarily implemented as hardware blocks.





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MPSoC architectures for SDR

 combine homogeneous and heterogeneous multiprocessing, including GPPs, vector processors and weakly programmable accelerators.





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- high workload: e.g. in smartphones, a digital workload of 100GOPS within 1W power budget.
- > One of our solutions is variation-aware dataflow-based design flow.



Variation-aware dataflow-based design flow

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Modeling



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 - scheduling techniques for scarce resources such as power and memory



Modeling

Mapping/scheduling



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Variation-aware dataflow-based design flow

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 - modeling techniques for radios, storage, arbitration, etc
 - scheduling techniques for scarce resources such as power and memory
 - analysis techniques to compute buffer-sizing, latency and throughput.



Modeling

Mapping/scheduling

Analysis **TU/e** Technische Universiteit Eindnoven

What Synchronous Dataflow Scenarios are and how to analyse them

How to model radios using scenarios dynamism in Long Term Evolution (LTE) and how scenarios capture dynamism in LTE



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- consists of actors that communicate tokens through FIFO channels.
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Iteration

- ▶ is a set of actor firings, as specified by the repetition vector.
- an iteration is marked by the production times of initial tokens, that is recoreded in a time-stamp vector γ.

Scenario-aware Dataflow (SADF)

 A scenario - a single mode of operation of an application.



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Synchronous Dataflow Scenarios

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- ► For a scenario, behavior is mostly invariable and modeled with SDF.



scenario 2



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Synchronous Dataflow Scenarios

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FSM **TU/e** Technische Universiteit Eindhoven

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FSM-SADF

is a tuple F = (S, f), consisting of a set of scenarios S and a FSM f.







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$$\mathbf{M} = \begin{bmatrix} 2 & 3 & 2 \\ 5 & 6 & 5 \\ 7 & 8 & 7 \end{bmatrix}, \ i_{xy} = \text{distance b/n tokens } y \& x \text{ tokens}$$

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• Period = 7 time-units & hence, throughput = $\frac{1}{7}$ iters. per time-unit.

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Exact methods*

 the inverse of maximum cycle mean (MCM) of the state space.



* M.Geilen and S. Stuijk, CODES+ISSS, 2010



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Conservative methods[†]

linear upper-bounds to γ_k



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- linear upper-bounds to γ_k
- find a vector τ
 ⁻ and a scalar number λ such that γ_k ≤ τ
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- Scenario graph: $\max_{s \in S} \tau_s + \lambda_s$

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- Scenario graph: $\max_{s \in S} \tau_s + \lambda_s$
- Reference schedule: MCM of the FSM where the weight of each node is τ_s + λ_s.

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LTE frame structure for FDD



 allocation of the grid to channels varies between sub-frames.

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FSM-SADF model of LTE



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Worst-case throughput computation

 WCT computation of the FSM-SADF model (×10⁻⁴ sub-frames/time-unit)

Method	1	2	3	4	5
Name	Static SDFG	Scenario graph	Reference schedule	State- space	MaxP∣us
WCT	2.6	5.2	6.6	8.9	8.9

Scenario-based techniques improve the static SDF result by 2 to 3.4 times more.



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- Methods 2 and 3 trade accuracy for lower run-time useful for iterative DSE algorithms.
- Methods 4 and 5 give the exact WCT, at the cost of run-time.
- Method 5 has a run-time in the order of tens of second.



Conclusions

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Thank you! Questions?



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